Dublin City University

Faculty of Engineering and Computing School of Electronic Engineering

Quality-Oriented Adaptation Scheme for Multimedia Streaming in Local Broadband Multi-Service IP Networks

Submitted for the fulfilment of the requirements for the degree of

Doctor in Philosophy (Ph.D.)

Gabriel-Miro Muntean

Supervisors: Dr. John Murphy Dr. Liam Murphy

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DECLARATION

I hereby certify that this material, which I now submit for assessment on the programme of study leading to the award of Doctor of Philosophy is entirely my own work and has not been taken from the work of others save an to the extent that such work has been cited and acknowledged within the text of my work.

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Signed:

Date: 23/001/03



Life is an ocean, love is a boat

In troubled waters that keeps us affoat

When we started the voyage, there was just me and you

Now gathered round us we have our own crew.

Dahon, "The Voyage"

Se lasa seara, Molly Malone, atarna luna de franjurii cetii, zeii din ceruri motaie-n tron in pub-uri canta petrecaretii.

Si infierbanta ispitele mute cupe de whisky si anason Si focuri sacre se-aprind nevazute-Se lasa seara, Molly Malone.

Se lasa seara, Molly Malone si toate-s parca o reverie, Iar din sageata lui Cupidon picura-ntr-una stropi de magie.

Insa cand haul cu neagra-i cange pe nesimtite de tine se-agata, toti zeii lumii-s mute falange, tacuti ca pestii tai din Piata.

Si lumea-i toata un Babilon. Ce neagra-i noaptea, Molly Malone!

Ivo Muncian, "Noapte la Dublin"

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II

Abstract

The research reported in this thesis proposes, designs and tests the Quality-Oriented Adaptation Scheme (QOAS), an application-level adaptive scheme that offers high quality multimedia services to home residences and business premises via local broadband IP-networks in the presence of other traffic of different types. QOAS uses a novel client-located grading scheme that maps some network-related parameters' values, variations and variation patterns (e.g. delay, jitter, loss rate) to application-level scores that describe the quality of delivery. This grading scheme also involves an objective metric that estimates the end-user perceived quality, increasing its effectiveness. A server-located arbiter takes content and rate adaptation decisions based on these quality scores, which is the only information sent via feedback by the clients.

QOAS has been modelled, implemented and tested through simulations and an instantiation of it has been realized in a prototype system. The performance was assessed in terms of estimated end-user perceived quality, network utilisation, loss rate and number of customers served by a fixed infrastructure. The influence of variations in the parameters used by QOAS and of the network-related characteristics was studied. The scheme's adaptive reaction was tested with background traffic of different type, size and variation patterns and in the presence of concurrent multimedia streaming processes subject to user-interactions. The results show that the performance of QOAS was very close to that of an ideal adaptive scheme. In comparison with other adaptive schemes QOAS allows for a significant increase in the number of simultaneous users while maintaining a good end-user perceived quality. These results are verified by a set of subjective tests that have been performed on viewers using a prototype system.

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Chapter I Introduction

Abstract

As an introductory chapter of this thesis, the first chapter presents the current situation in the market of multimedia presentations that tends to hugely develop and expand by changing the very manner these services are delivered to customers. It seems that this significant change is related to providing high quality, interactive and/or on-demand multimedia-related services to home residences and business premises. The chapter starts with a presentation of this tendency, the associated problems and the challenges that come with this development. Different existing solutions are then mentioned looking at network-related technologies, provided services, technical solutions for multimedia distribution and the consequent provided quality and necessary efforts. Next the chapter describes the motivations of the work that stands behind this thesis and states the problem and the goal of the research. The proposed solution is then presented and the significant contributions of the thesis are listed. A short outline of the thesis ends this chapter.

1.1 Multimedia Presentations

Multimedia presentations have taken at least four major divergent directions: i) shows in cinema theatres, presentation halls etc., ii) programs delivered via broadcast TV, radio, cable TV, etc., iii) movies and documentaries played from tapes and DVDs rented and/or bought and iv) multimedia streaming over different types of networks, including the Internet. Each of these directions has significant advantages and important disadvantages that make it more or less popular than the others. The cinema spectators for example may appreciate the high quality of the shows and the opportunity to socialise, but this involves physical presence of many people in hub like halls with inherent problems such as booking, traffic, parking, etc. The home comfort as opposed to the latter made the delivery of multimedia programs to homes via TV or radio very popular. The latest enhancements such as cable TV and digital TV provide the viewers with a wider choice, offer

interactivity, and introduce new services like for example tele-voting, T-mailing and T-gaming. Apart from this, a large number of viewers have found that the rented/bought video tapes are very convenient since they allow for choosing both the show and the time of presentation. This has shown that the one-fits-many approach, which is economically beneficial, is not what the customers desire, but rather one-fits-one solutions that allow for choice flexibility. The DVDs and the latest home theatre systems, that have added high quality to multimedia presentations, brought them closer to the cinema experience while offering to the users the TV-related conveniences and the possibility of both content and showing time selection.

The computer-based multimedia streaming, a very different type of multimedia presentation, has become increasingly popular lately, especially over the Internet, attracting millions of users. The exponential increase in computer users, in Internet-connected computers and in quantity of information, including multimedia data, available and exchanged via the Internet that have exceeded a linear increase in available resources, made very likely congestion to appear. The congestion and the consequent losses that affect the multimedia viewers in their perceived quality are the greatest disadvantage of multimedia streaming. Another disadvantage is the need for some basic training in order to allow for using computer-based services. Also today the quality of these multimedia presentations is much lower than that of the other presentations previously mentioned. A definite advantage is the large variety of available services offered (for example in the same class of multimedia streaming-based services there are radio, Web-TV, pre-recorded and live multimedia transmissions, educational presentation, etc.). Another advantage is the convenience of using these services in conjunction with other ones, Internet-related or computer-based.

Currently there is a trend that very likely will cause a major change in the way information and entertainment are delivered to consumers (a significant part of them in the form of multimedia presentations) [1, 2, 3]. It seems that the existing parallel directions of multimedia presentations are going to merge in the form of on-demand access to rich media and full-motion high quality multimedia to home residences or business premises, as part of a large set of personalised high-quality services. This will take advantage of some of the benefits and will minimise some of the disadvantages related to different types of multimedia presentations. The success or failure of this trend depends on widespread market acceptance, which, in turn, relies heavily on the technical solutions involved, on the popularity and the quality of services provided, and on the price the enduser must pay.

Briefly, for the on-demand access to high-quality multimedia presentations from homes to be successful, there is a need for:

- A delivery network that can support increased resource requirements related to highquality multimedia streaming and the delivery of other heterogeneous services
- A wide range of attractive personalised on-demand high-quality services that can determine the customers to choose paying for the new solution
- A delivery solution that offers high quality services that will both attract the customers and will allow for the service providers to make profits.

Next possible solutions for each of the above-mentioned problems are briefly presented.

1.1.1 Delivery Networks

The problem of choosing a delivery network for high-bitrate multimedia traffic to the homes with tightly imposed cost constraints is not simple. This becomes even more complex when, due to economic pressures, other types of services are required to use the same infrastructure in order to reach the customers. This is unlike what happened in the past when service providers and network operators have built separate networks for different services provided (e.g. telephony, cable TV, etc.).

The technologies that allow access to residential users could be either wireless or wireline. Wireless distribution options include fixed terrestrial wireless, wireless local area networks (WLAN), mobile wireless and satellite systems. Wireline solutions include the telephone network, the cable TV, the power line network and the separate distribution infrastructure built by so-called over-builders. More details about these solutions are given in the second chapter that presents the related works. It is worth to mention that the emerging wireline broadband IP networks constitute an important solution for distributing these high bit-rate multimedia-based services to the viewers. However, for their success, other services have to be offered as well, and solutions for their distribution have to be proposed in order to make them more appealing to the customers.

1.1.2 Offered Services

Some of the most important services that could be offered via broadband networks are digital and interactive TV, digital and interactive audio, high-speed data transmission, and other

services such as gaming, betting, voting, banking, shopping etc. More details about these services are presented in the second chapter. However it is important also to use a distribution solution for these services in order to ensure their high quality and good utilisation of the existing infrastructure.

1.1.3 Distribution Solutions

The new services associated with broadband connectivity can become successful and attract a large number of customers if their quality is high, their price is low and they bring benefits to both network operators and service providers. The quality is assessed depending on the service provided, varies with the technical solution chosen for the delivery of the service and is subject to subjective considerations. The price paid by the customers and the benefits for the network operators and service providers are influenced by the overall performance of the delivery solution. Significant components of the service distribution performance are the infrastructure utilisation, the number of customers simultaneously served with a certain service or group of services, and the quality of these services.

In the next chapter the term "quality of service" (QoS) is defined and its meaning in relation with the quality of broadband services is explained in detail. Then, different solutions for providing desired QoS, their advantages and disadvantages are presented along with different options for assessing the quality of the provided services, mainly multimedia-based. Among the best-known solutions for providing QoS are bandwidth over-provisioning, traffic engineering, QoS architectures and application-level adaptive solutions. The application-level adaptive schemes, which take the distribution networks as they are, provide the least complex and the most flexible mechanisms for providing certain QoS, although with no guarantees. These are the main reasons, for focusing the research presented in this thesis on an approach based on application-level adaptive schemes.

1.2 Research Motivation

For 2003 and the near future, in spite of the global economic slowdown, IDC¹ estimates a sustained growth in the number of broadband connections to residential users (e.g. broadband connections will surpass 20 million in Europe alone), while the equipment and product markets will continue to grow in volumes (i.e. the expansion drive will be the differentiated product offerings

¹ IDC, http://www.idc.com

and an increase in the availability of broadband specific content and applications). However a further deterioration of price levels will affect the revenues of service providers and network operators (e.g. a 9% drop is expected for 2003) [4]. Therefore the trend towards multi-service IP-networks that would allow the use of popular IP-based applications and low cost hardware predicted in [1, 2, 3] may be accelerated. At the same time a GartnerG2 study [5] concluded that the consumers are prepared to pay a premium for broadband connectivity only in conjunction with a "must have" application that may convince them they need broadband (e.g. fewer than 10 percent of Internet households think broadband alone currently provides good value). Related to possible services to be attracted by, a 2002 study² found that the broadband services the most US households would pay for are those that have multimedia components, especially entertainment services (44% of the subjects), communications-based services (42%), and education-related services (39%). All of these have both high bandwidth requirements and timing constraints that may put significant pressure on the network providers' delivery infrastructure. They also suggest that the service providers have to offer a wide range of services with rich content in order to become attractive for the residential customers.

In consequence, as previously mentioned, the networks used for delivery, the attractiveness, range and quality of the provided services and the technical solutions for distributing these services to their receivers are of a paramount importance for a successful wide-scale deployment of these high-quality services. Different possible solutions have already been discussed, and their advantages for the network operators, the service providers and the customers have been assessed. In this context the motivations for this work are presented briefly as follows.

Need to Support High Diversity of Services

The service providers, the network operators and the customers look forward at providing, respectively having access at highly diverse services such as VoD, VoIP (IP telephony), high rate data transfers, etc. However these services have different types and therefore various requirements that have to be accommodated, while being delivered by the same multi-service broadband IP-based infrastructure. In this context there is a need for multimedia-based services that influence or are influenced in a minimal manner by traffic produced by other type of services (e.g. data transfer).

² Michael Pastore, "Broadband Lacks a European Audience", CyberAtlas, Feb. 5, 2002, http://cyberatlas.internet.com

Increased Network Infrastructure Utilisation

Service providers and network operators have to take full advantage of the existing network infrastructure and make incremental investments to support revenue-producing services in order to increase service penetration and improve infrastructure utilisation. Increasing the number of simultaneously served customers and the network utilisation decreases the quality of service in general. Thus there is a need to balance the goals of providing high-quality rich content services of diverse types, and of reducing the network infrastructure necessary for the provision of these services.

Personalised Services to Heterogeneous Customers

The scalability issue may have another dimension apart from number of viewers: heterogeneity of customers. In order to be considered acceptable, any novel multimedia-based solution has to be able to satisfy customers with different expectations. Therefore there is a need for the "one-fits-many" approach to be replaced by "one-fits-one", providing personalised, interactive services to customers that may be connected via heterogeneous links.

Trade-off Between Performance and Quality

QoS solutions in generally involve many trade-offs. For example in multimedia streaming in order to reduce the quantity of data to be sent across the network, compression algorithms are being used that remove streams' redundancies, but leave the streams vulnerable to transmission errors. To further reduce the quantity of data lossy multimedia encoding techniques purposely leave aside some information, reducing the quality of the streams. As results, the higher the compression rate is and therefore the narrower bandwidth necessary for transmission, the lower the streams' quality and the lower their resilience to potential transmission errors. Similarly for time-sensitive applications, smaller size buffers help reducing streaming delays, but cannot accommodate highly bursty traffic causing losses that more severely affect the quality of the remotely transmitted streams. In consequence there is a need for very good trade-off between the performance and quality, especially in the presence of different types of traffic.

1.3 Problem and Goal

Broadband multi-service IP networks are either being deployed by over-builders or through transformation of existing cable TV networks, and many popular IP applications are ready to be provided as services. However, a technical solution to the provision of these services is still required.

The problem this thesis addresses consists of delivering multimedia-based services via local broadband multi-service IP networks while balancing:

- customers' need for high quality service
- service providers' and network operators' goal of increased infrastructure utilisation and more customers served.

Since apart from being time-sensitive, these services have very high bandwidth requirements that make their support expensive, the latter is achieved by building an *inexpensive* application-layer adaptive mechanism that would adjust the transmitted quality level to the delivery conditions only when congestion is building up and may severely affect the quality of the service provided. This mechanism should allow for serving a higher number of customers from limited resources, which would constitute the main benefit of the proposed solution. However, due to the routine-like daily and weekly schedule for the customers with periods of very high and very low usage of the multimedia-based services that are the highest bandwidth consumers, it is expected that such adjustments to be only temporary matching current peak times of the cable TV service audience (i.e. mainly evening). This adaptive mechanism should maintain good end-user perceived quality for the delivered services in order to meet the customers' quality expectations.

The goal of this work is to propose, design and test an application-level end-to-end adaptive mechanism for streaming multimedia that offers high quality of services to home residences via local broadband IP-networks subject to very high traffic of different type, with various size and variation pattern. The scheme should not interfere with the provided services' interactivity and personalisation characteristics and should find a solution for the adaptation that has the least effect on the end-user perceived quality.

1.4 Solution and Contributions

The work's goal was achieved by proposing the Quality-Oriented Adaptation Scheme (QOAS) for adaptive multimedia deliveries via local broadband multi-service IP networks. This scheme relies on estimates of the end-user perceived quality as the best assessment of the degree of QoS provided made at the client. For this estimation to be accurate, some network-based parameter values, variations and variation patterns were mapped into application level QoS grades that reflect the quality of delivery. These grades are then send via feedback to the server and used to trigger adaptive adjustments of the streaming process according to the reported delivery conditions in order to provide best quality given the situation.

The proposed adaptive scheme has been designed, modeled, implemented and tested through simulations and a real prototype system in order to both verify and validate the scheme's performance. Also the end-user perceived performance as estimated by an objective metric has been measured and the scheme's behavior has been assessed according to the test results. These tests have first checked the scheme's adaptive reaction to sudden changes in network's traffic and have included traffic of different type, size and variation patterns. Then the effect some variations in the parameters used by the adaptive scheme have on its performance has been tested, along with the influence of some network-related parameters characteristics on its functionality. It was also very important to test the benefits brought by this adaptive scheme in terms of estimated end-user perceived quality, network utilization, loss rate and number of customers served by a fixed infrastructure in comparison with an ideal scheme that would use all the available bandwidth, would achieve 100% utilization and 0% loss. The QOAS performance was also compared with other proposed mechanisms for delivering multimedia.

Since there is not a generally accepted metric for the objective assessment of the quality of video streams, subjective tests have been performed on real viewers. For this a prototype system has been built that makes use of the proposed adaptive scheme and a test bed that would allow for different other traffic to interfere with the multimedia traffic generated by the prototype system was used in order to test the scheme's performance. The subjective test results have verified and confirmed the good results obtained from the simulated tests.

Next the contributions of the QOAS - the proposed application-level adaptive solution for high quality multimedia streaming in local broadband multi-service IP-networks - are highlighted:

 QOAS uses a novel client-located grading scheme that maps some network-related parameters' values, variation and variation patterns on application-level QoS scores that describe the network's traffic conditions. The QOAS adaptation is based only on transmitting these client-computed QoS scores that estimate the current quality of delivery to the server

- The end-user perceived quality as estimated by an objective metric is actively considered during the adaptation process, increasing the effectiveness of the adaptation and the estimated end-user perceived quality
- The scheme's behaviour is very close to one of an ideal adaptive scheme in terms of estimated end-user perceived quality, loss rate and link utilisation when used for multimedia streaming in the presence of traffic of different types, size and variation pattern. During testing the end-user perceived quality while using QOAS was within 1% from the one if the ideal adaptive scheme was used, the loss rate was almost 0% and the link utilisation was more than 99.6% in the large majority of cases.
- The scheme allows for a significant increase in the number of customers that can be simultaneous served while maintaining a good end-user perceived quality, even in comparison with other existing solutions for delivering multimedia. The results of the tests performed show that 23% more customers could be served by using QOAS than by using TFRCP [6], 33% more clients than by using LDA+ [7], and 39% more users than by using a non-adaptive solution.

1.5 Short Outline of the Thesis

This thesis is organised in eight chapters that present the subject of the research performed, the related works, the proposed solution and its testing and conclusions drawn.

This first chapter has mainly presented the motivation for the research, the problem to be solved, the goal, the solution and the contributions. The second chapter presents different related works, whereas the third describes the context of the solution. The forth chapter focuses on the detailed presentation of the proposed application-layer adaptive scheme - QOAS and includes the architecture of the multimedia delivery system that implements it. The fifth chapter aims at presenting both the simulation model and the prototype system that have implemented QOAS and were used for testing it, whereas the sixth chapter presents the tests performed and their results. The seventh chapter draws some conclusions and highlights possible future work directions. The list of references and the appendixes end the thesis.

1.6 Summary

The first chapter starts with a presentation of existing divergent directions in multimedia presentations including cinema shows, TV programs, tape and/or DVD movies and multimedia streaming. It then presents their merging tendency in form of on-demand-based access to rich media and full-motion high quality multimedia to home residences as part of a large set of personalised, high-quality services. In order for this to be successful a delivery network that would support the increased requirements in resources related to high-quality multimedia streaming and the delivery of other heterogeneous services is needed as well as a wide range of attractive, personalised, on-demand high-quality services that would make the customers pay for them and a delivery solution that would offer high quality for the services at a low cost. Next this chapter mentions existing solutions related to these three issues. In this context the chapter also presents the motivation for the research, the problem to be solved and the research's goal and it ends with a description of the proposed solution and of the significant contributions made.

Chapter II Related Works

Abstract

The second chapter of this thesis presents significant works related to the proposed Quality Oriented Adaptation Scheme (QOAS). These works were classified in directions of interest and span from different solutions for achieving success in providing high-quality multimedia-based services to remote viewers to compression techniques, adaptive delivery solutions and multimedia streams user-perceived quality assessment. Also solutions for improving the performance of multimedia deliveries are explored looking from a broad point of view and including error control, delivery techniques, protocols and delivery architectures. Comments are made and conclusions are drawn in relation to the applicability of the presented works to QOAS in broadband IP-networks.

2.1 Overview

When multimedia data is transmitted over an IP network, including a broadband multiservice IP network, among the very few assumptions about the capabilities of the network that could be made is that it is able of delivering packets to a destination. However, there are no guarantees that all packets will be delivered and there is not any mechanism to inform if a packet does not reach its destination. If a sequence of packets is sent to the same destination, the host computer must not assume that the network will maintain packet order, and also it must not assume that the network will maintain the relative timing of the packets. Also the source of data cannot assume that there is any particular throughput rate, bandwidth, or end-to-end delay.

In this context extensive research has tried to find solutions in order to provide desired Quality of Service (QoS) for applications with different requirements, mainly time sensitive or resource intensive. This chapter defines QoS and then presents in detail proposals for providing QoS and directions for QoS assessment.

Among the least complex and most flexible solutions for providing desired levels of QoS, adaptive applications for multimedia streaming take the networks as they are and employ different mechanisms that complement the IP network's basic functionality. For example adaptive control schemes could inform the applications about the loss rate, throughput or other network-related parameters or about the estimated perceived quality at the end-users or could take adjustment measures (e.g. modification of the transmission rate) in order to improve the quality of delivery if decided to be necessary. Although these adaptive control solutions cannot guarantee the provision of certain QoS, they will increase the quality of delivery over high loaded networks, trying to avoid congestion. In this chapter, different research directions related to the **adaptive control schemes** are presented and some of the most significant solutions.

Complementing the effort of these adaptive control schemes aimed at distributing high-quality multimedia streams with high bandwidth requirements and tight timing constraints, other proposed solutions can be used in conjunction, in order to provide increased quality with little effort. Since bandwidth is a limited and expensive resource, among the mostly used solutions are **compression techniques** that reduce the quantity of data to be sent across the networks, while maintaining a good quality for multimedia streams subject to compression. Measuring the end-user perceived quality is also significant in the effort to provide the adaptive control schemes with accurate information about the effect the network conditions have on the quality of delivery. This is also important during the development stage when the solutions have to be tested. Therefore there is a need for **the end-user perceived quality assessment**.

Other solutions are used in conjunction in order to provide increased performance of multimedia deliveries in terms of quality, bandwidth requirements and cost. Among them the error control mechanisms have the capability to detect and correct errors, (i.e mainly transmission errors), minimising their effect on the quality of the service provided and increasing therefore the expected end-user perceived quality. Different network approaches for the localisation of information to be accessed were also taken into account, including caching, proxy servers and content distribution networks that were devised and used to bring the data closer to the customers in order to minimise its transport paths over the loaded sections of the networks. Different delivery solutions, including broadcast, multicast and unicast, were proposed to deliver the same content to one or a group of receivers in a one-to-one or one-to-many approach, balancing the need for reducing delivery effort with the increase in personalisation of provided services.

In this chapter these directions are also explored, proposed solutions are mentioned, their performances are compared and the most suitable for the usage as part of the proposed application-layer adaptive scheme (QOAS) are indicated.

2.2 High-Quality On-Demand Multimedia Presentations

As previously mentioned, for the success of the on-demand high-quality multimedia presentations services to home residences and business premises, there is a need for: a delivery network that can support resource-intensive heterogeneous services, a wide range of attractive personalised on-demand high-quality services and a delivery solution that offers high quality services while best utilising the infrastructure. Possible solutions for each of these problems are presented in detail next.

2.2.1 Delivery Networks

The technologies that allow access to residential users could be either wireless or wireline. Wireless distribution solutions include fixed terrestrial wireless, wireless local area networks (WLAN), mobile wireless and satellite systems. Wireline solutions include the telephone network, the cable TV, the power line network and the separate distribution infrastructure built by so-called over-builders.

2.2.1.1 Wireless Solutions

Wireless solutions are cheaper than wireline ones, having the advantage of low deployment cost (no wires), although there is a licence fee for the spectrum. Among them, fixed terrestrial wireless services provide connectivity from a base station to a stationary point (e.g. home). First-generation of commercially proprietary systems could provide data rates of 1 - 10 Mb/s making use of certain spectrum bands. For this purpose local multipoint distribution service (LMDS) and multipoint multi-channel distribution service (MMDS) were defined. In reality operators like Sprint³ and MCI WorldCom⁴ have provided only 1 Mb/s. Although second-generation LMDS and MMDS have been assessed for standardisation by bodies like IEEE in its 802.16 Working Group⁵, the

³ Sprint, http://www.sprint.com

⁴ MCI WorldCom, http://www.mci.com

⁵ IEEE 802.16 Working Group on Broadband Wireless Access Standards, http://grouper.ieee.org/groups/802/16

bandwidth provided is not enough for delivering very high-quality multimedia. Providing real broadband will be possible by addressing, WLAN technologies or by using the future IEEE 802.11 and its variants⁶ and the European Telecommunications Standards Institute (ETSI) HIPERLAN/2⁷ standards that are meant to provide speeds of up to 50 Mb/s. Unfortunately the coverage area of WLAN-based solutions is limited to microcells with a typical radius of less than several hundred feet that does not eliminate totally the need for broadband wired access. Mobile wireless has a completely different approach, targeting portable and mobile communication and computing devices, instead of broadband. However even speeds of 2 Mb/s theoretically achieved by the International Telecommunication Union (ITU) 3G standard IMT-2000⁸ (limited in practice to hundreds of kilobits per second) are not enough for real broadband local networks [8]. Satellite-based solutions already provide broadband services mainly for broadcasting programs, but have performance limitations that make them unsuitable especially for interactive and bi-directional communications. These limitations refer to high latencies and uplink-related problems.

2.2.1.2 Wireline Solutions

Wireline solutions, although more expensive due to the cost of wiring, could be more effective if already existing infrastructures are used. Telephone networks, mainly twisted-pair copper-based, with slight upgrades, have already been used to provide near-broadband connectivity using Digital Subscriber Line (DSL)-based solutions. Different DSL flavours such as asymmetric DSL (ADSL), G.lite DSL and very high-data rate DSL (VDSL) were proposed to either reduce cost by providing a very narrow upstream channel, to operate concurrently with the telephone service or to provide speeds of tens of megabits per second. The latter operates only if fiber is introduced into the network to reduce the copper lines' lengths. This is because the biggest problem with DSL is that it does not work over wires longer than a certain distance (18,000 feet for ADSL) [8]. The DSL transmissions could be also affected by interference from signals from adjacent lines. These problems limit the usage of the telephone networks as support for broadband connectivity. Power lines could be used for data communications and although the technology allows for a 10 Mb/s connectivity⁹, there is a strong concern about the interference with well-established wireless

⁶ IEEE 802.11 Working Group for WLAN Standards, http://grouper.ieee.org/groups/802/11

⁷ The European Telecommunications Standards Institute, HIPERLAN/2, http://www.etsi.org/technicalactiv/hiperlan/hiperlan2.htm

⁸ ITU - International Mobile Telecommunications-2000 (IMT-2000), http://www.itu.int/home/imt.html

⁹ LEA, http://www.lanergy.com

applications such as amateur radio, emergency broadcast services, etc. Using them on large scale is not a viable solution until these reasons of concern are dealt with. **Cable TV networks**, the existing hybrid-fiber-coax (HFC) systems that feed the conventional TV sets or set-top boxes and **over-builders' cable infrastructures** (e.g. RCN¹⁰), the newly deployed structures with fiber at their core offer services that could be easily upgraded to serve the distribution of rich content media and other services, following the principle "pay-as-you-grow". In this thesis cable networks and cable operators terms refer to both of these cases.

2.2.1.3 Cable-based Solutions versus Satellite Broadcast

The cable operators are currently in a difficult position, after being challenged successfully by satellite broadcasting companies that are offering more channels at a lower price, with a significant impact on their subscriber base. Therefore many of them [1] have already started to consider their competitive advantages and are trying to shift fundamentally their business approach.

- First they can easily move from broadcast to unicast, offering more personalised content delivery, according to subscribers' needs. Briefly they can offer what the customers want and when they want it ("on-demand").
- Second, the possibility of using the return channel and the low latencies involved make the introduction of a new set of interactive services possible.
- Third, upgrading their infrastructure by introducing fiber, the bandwidth offered becomes very competitive.
- Fourth, by using Gigabit Ethernet and switching to IP-architectures many other services
 could be offered to the subscribers based on numerous and very popular IP-based
 applications, including high quality multimedia streaming, by taking advantage of the
 broadband availability.

Therefore, many cable operators in different parts of the globe have already upgraded their networks by introducing fiber into their systems offering broadband connectivity to residential users. For example if in 2001 the percentage of households with broadband connections was very low in Europe (1.93% in Germany, France and Britain), and moderate in America (13%)² and parts

¹⁰ RCN, http://www.rcn.com

of Asia (17% in Korea)¹¹, in 2002 the market for broadband services experienced a significant increase (it doubled in Europe, reaching more than 12.6 million homes¹²). This is while the service providers are continually introducing Fiber-to-the-home (FTTH) and Fiber-to-the-curb (FTTC) systems on a large scale¹³ in order to offer even higher bandwidth. In USA, for example, besides BellSouth¹⁴, Sprint³ and Verizon¹⁵, which already use FTTH and FTTC systems to service more than 300,000 households, several new broadband service providers are building large scale FTTC or FTTH systems in California, Tennessee and Texas [9] and municipalities or other public authorities have launched FTTH projects in various towns. At the same time the estimates show¹³ that the FTTH systems in the USA will reach 2.65 million homes by 2006 and FTTC systems - another 1.9 million in an on-going expansive process.

2.2.1.4 Broadband Multi-service IP Networks

Meanwhile, apart from cable operators' interest, there is also an industry push for the development of all-IP-multi-service distribution systems from key players such as the International Engineering Council (IEC)¹⁶ and Cisco [10], the market leader in IP routing. This push becomes increasingly consistent with time and is accompanied by efforts driven by the Multi-service Switching Forum (MSF) - founded in 1998 by Cisco¹⁷, Bellcore/Telcordia¹⁸ and MCI WorldCom⁴ - that aim at developing standards and architectures for an open, standards-based network that supports multiple services on a common network infrastructure [11]. Part of this effort is the multi-service IP network seen as the next generation network that would support both the delivery of high quality multimedia streams and different other IP-based services to both home residences and businesses [2]. For instance, since 2001 there are companies on the market like GoldTV¹⁹, a provider of broadband on-demand multimedia-related services based in Milan, Italy, and

^{11 &}quot;Korea Leads Broadband Internet Service Market", KoreaNow, Nov. 30, 2002, http://kn.koreaherald.co.kr

¹² J. H. Bakkers, "European Broadband Market Predictions And Preliminary Analysis", Jan. 2003, http://www.idc.com

¹³ KMI Corporation, "Fiber-To-The-Home To Reach 2.65 Million Homes By 2006", Press Release, 2001, http://www.kmicorp.com/press/011015.htm

¹⁴ BellSouth, http://www.bellsouth.com

¹⁵ Verizon, http://www.verizon.com

¹⁶ International Engineering Consortium (IEC), http://www.iec.org

¹⁷ Cisco Systems, http://www.cisco.com

¹⁸ Bellcore/Telcordia Technologies, http://www.telcordia.com

¹⁹ GoldTV Italia, http://www.tvgold.it

NovaMedia²⁰, a digital broadcast and on-demand multimedia provider based in Reykjavik, Iceland that really do deliver rich content services over such a multi-service all-IP based infrastructure.

This significant evolution towards broadband connections and a large interest in providing diverse services from a low-cost infrastructure will offer a large market opportunity for the IP-based services. In the near future, a global evolution from the existing Hybrid Fibre Coax (HFC) networks towards all-IP architectures that would allow an almost universal use of popular IP-based applications and low cost hardware is predicted [1, 2, 3]. In consequence new services that make use of this infrastructure, including on demand delivery of multimedia ("video-on-demand" - VoD), that have already been launched, are waiting for very large-scale deployment. For this to happen, other services have to be offered as well, and solutions for their distribution have to be proposed in order to make them more appealing to the customers.

2.2.2 Offered Services

Broadband subscribers can be grouped in two basic classes: "lean forward" and "lean back" users [12]. The first ones are typical PC users of high-speed Internet services and most of their activity is very interactive in nature, "leaning" forward to access the service. The latter ones are non-interactive by nature, passive, "leaning" back on the chair and enjoying the experience. Currently the first category constitutes the base for broadband subscribers and the corresponding market is expanding slowly. The second type of potential subscribers represents a very large market opportunity for the network operators and service providers, which could be transformed into revenues by offering among other services ones that would make them interact. Shortly, the customers will be able to watch selected movies on request, to send messages, to shop, to learn, to explore websites with rich content, to watch live programs, to record them, to listen to high quality audio, to select and listen to radio stations, to download quickly data, to take part in interactive gaming and debates, etc. This will improve their experience of studying, communicating, shopping and mainly of being entertained. Introducing these services through their TV sets would make the transition easier to a world in which the future-TV and the computer will be synonyms.

Some of the services that could be offered via broadband connections are presented next.

²⁰ NovaMedia Iceland, http://www.media.is/pdf/interactivetv.pdf

2.2.2.1 Digital and Interactive TV

Digital and Interactive TV includes Digital Video Broadcast, Pay-Per-View, Personal Video Recording, Video-On-Demand, Near-Video-On-Demand, Videoconferencing, Electronic Program Guide, T-Commerce, etc.

Digital Video Broadcast or Digital TV, provides very high image quality, better resolution and colour while using transmission related facilities more efficiently than the analog TV. It uses broadcast or multicast to reach the customers and it is seen only as a first stage in delivering rich content video broadband services, as it does not provide customer personalisation and flexibility. Pay-Per-View (PPV) services represent Digital TV programs that are transmitted unicast or multicast only to users that have paid to view them. These services are mainly used for live events. Personal Video Recording (PVR) or Time-shifted TV (TsTV) allows for VCR-like controls for the live transmitted programs: recording, pause, skip, rewind, etc. Video-on-Demand (VOD) and Subscription Video-on-Demand (SVOD) are the ideal applications for broadband IP networks. They provide entertainment-on-demand by taking advantage of the networks' two-way communication capabilities. They are like video or DVD rental and in general the specified content is unicast streamed only to those users that have paid for the service. They provide full VCR-like controls. SVoD refers to the subscription to an entire series. A VoD service is termed Near-VoD (NVoD), if the subscribers who order a particular movie to start within a specific time window are grouped together [13] in order to save bandwidth. The major disadvantage of NVoD is the lower flexibility offered to the customers. Videoconferencing allows for two or more people at different locations to see and hear each other at the same time, sometimes even to share computer applications for collaboration. This offers new possibilities for schools, libraries, businesses, including formal instruction, connections with guest speakers and experts, multi-party project collaboration, professional activities, and community events. Interactive Program Guide (IPG) or Electronic Program Guide (EPG) is an on-screen listing of the available programs, which can be organised by channel, time, genre or personal interest in a user-friendly manner. The desired program is regularly selected using the remote control. T-Commerce refers to online commerce done through the TV environment (choose and buy using the remote control and the TV set).

2.2.2.2 Digital and Interactive Audio

Digital and Interactive Audio services include Digital Radio Broadcast, Audio-On-Demand services, Voice over IP (capable of providing IP telephony, etc.).

Digital Radio Broadcast or Digital Audio Broadcast²¹ provides reliable interference free reception - especially in-car - and near CD quality sound, and could be complemented by additional data services, such as listing song titles, news and sports updates, next program to be broadcasted etc.. Audio-On-Demand (AoD) services are similar to VoD and provide music, news, and audio-related entertainment in general at request. Normally the specified content is unicast transmitted to customers, allowing for VCR-like capabilities. Voice over IP (VoIP) describes the use of the Internet Protocol (IP) [14] to transfer speech/voice between two or more sites. In general, this means that the voice signal is sampled, compressed and encapsulated into data packets and then transferred across an IP-based network along with all other data packets. VoIP is mainly used for IP Telephony, significantly reducing costs in comparison with classic circuit switching-based solutions.

2.2.2.3 High Speed Data Transmission

High Speed Data Transmission refers to the transmission of different content data packets with a significant higher speed in comparison with dial-up networks, for instance. The most important applicability is for downloading data files for later usage, transferring information to be displayed during WWW browsing, etc. This enables shorter waiting times and determines significant increases in customers' satisfaction with the service provided.

2.2.2.4 Other Interactive Services

Typical interactive services offer the possibility for the viewer to interact with the television set in multiple ways. VoD and AoD are interactive services, but the customers could also play games, place bets, vote or provide immediate feedback to a program, debate on certain subjects, do banking and shopping, etc.

Among the previously mentioned services the large majority are based on multimedia delivery to the customers and the most complex of them is VoD. Some of the capabilities VoD offers to its viewers are high quality and extended choice in terms of content, playing time and control (VCR capabilities) making it the ultimate experience of home accessible entertainment. Unfortunately these kinds of personalised multimedia services are important bandwidth consumers and in order for the offered services to be cost-effective, significant effort has to be made to

²¹ "A Guide to digital radio (Digital Audio Broadcasting)", http://www.radio-now.co.uk/faq2.htm

increase the number of VoD customers that could be served by a limited infrastructure. Naturally the perceived quality of the provided service should not be affected by the increased number of customers and by the rest of traffic carried by the same infrastructure. At the same time the service must not severely affect the background traffic.

2.2.3 Distribution Solutions

In this thesis the distribution solutions for delivering these services to the customers are assessed in terms of their quality and the utilisation of the existing infrastructure. The quality depends very much on the service provided, varies with the technical solution chosen for the delivery of the service and is subject to subjective considerations.

In this section first the term "quality of service" is defined and its meaning in relation with the quality of broadband services is explained. Then, different solutions for providing certain level of quality of service are presented and choices for gracefully reducing it if and when needed are analysed such as the end-user perceived quality is maximised while better utilising the existing infrastructure. At the end different options for assessing the quality of the provided service are described, mainly in relation to multimedia presentations.

2.2.3.1 Defining Quality of Service (QoS)

The Quality of Service (QoS) is defined by the ITU in ITU-T R. E.800 [15] as the "collective effect of service performances that determine the degree of satisfaction by a user of the service", by the ISO/IEC 10746-2 [16] as "a set of qualities related to the collective behavior of one or more objects" and by IETF in RFC 2386 [17] as "a set of service requirements to be met by the network while transporting a flow". The ITU-T definition closely relates QoS to the users' perception and expectations related to a certain service wheareas the ISO/IEC's is more general, but with direct applicability in networking. The IETF's definition involves more the idea of a "flow" than individual or group of packets suggesting QoS usage in connection with streams. At the same time the industry leaders define QoS closer to their object of activity. For example for Cisco QoS "refers to the capability of a network to provide better service to selected network traffic over various technologies" [18], while for Microsoft QoS "refers to the ability of the network to handle the traffic such that it meets the service needs of certain applications" [19].

Detailing the QoS definition, the ISO/IEC's view, which was also shared by ITU-T in R. X.902 [20], is that QoS concerns characteristics like the rate of information transfer, the latency, the

probability of a communication being disrupted, the probability of system failure, the probability of storage failure, etc. They also mention possible constraints that may affect the QoS, which include temporal ones (e.g. deadlines), volume constraints (e.g. throughput) and dependability, involving aspects of availability, reliability, maintainability, security and safety (e.g. mean time between failures). ITU-T R X.641 [21] and ISO/IEC 13236 [22] define the QoS Framework and associated concepts, which are described in order to highlight the QoS management. They present also the QoS characteristics and how QoS requirements drive the selection and use of QoS management functions and QoS mechanisms.

QoS is very complex and as there is not a widely accepted definition for QoS, there are not general solutions for assessing, providing and quantifying QoS. However different aspects of QoS are explored according to certain interests that have driven extensive research in a direction or another. Since the main interest of the research presented in this thesis is to provide high QoS levels in broadband IP-networks while having certain constraints, details are given about research directions that have proposed solutions in this context. Therefore different solutions for providing QoS are assessed next along with diverse proposed parameters that are associated with QoS.

2.2.3.2 Providing QoS

The IP-networks provide a single type of service often named "best effort", because "best effort" is undertaken to deliver packets as quickly as possible, treating all of them equally, in a perfect impartial and fair approach. This service is suitable for many applications and services such as WWW-based document retrieval and FTP-based data transfers. However many of the services meant to attract the customers like VoD are flow-based, have high resource requirements, and mainly are time-sensitive, generating a different type of traffic for which the "best effort" is not good enough to ensure certain QoS level at all times. At the same time these different types of services have to co-exist and be able to be served by the same network infrastructure. In consequence different methods for providing QoS are required in order to support both existing and emerging applications and services, which have different characteristics.

Extensive research was focused on providing QoS in different conditions, for various technologies and architectures and with different approaches. Among these the best known are bandwidth over-provisioning, traffic engineering, QoS architectures and application-level adaptive solutions.

2.2.3.2.1 Bandwidth Over-provisioning

One way to overcome the limitations of the "best effort" networks in providing QoS is by over-provisioning, that is by allocating more bandwidth than the expected network peak requirements [23]. Still, even though over-provisioning of network increases the probability of having enough resources available for real-time applications, it still does not guarantee the desired QoS at all times. The problem is that data, and especially multimedia data, are inherently bursty and regardless of capacity, congestion is very likely to occur for short periods of time. Another consideration is that the normal routing protocols do not know about load levels, so congestion will build up on some paths while others have bandwidth to spare. Also bandwidth alone does not ensure low and/or predictable delays, as even with huge bandwidth, there is still the possibility that large file transfers will interfere with real-time application traffic. On the other hand there will always be a waste of resources in an over-provisioned network, for instance during off-peak times, which is not economically justifiable. However, even a perfect solution based on over-provisioning is only temporary as a corollary of Moore's Law²² states that "as one increases the capacity of any system to accommodate user demand, user demand will increase to consume system capacity".

These considerations lead to the idea that other ways of providing QoS should be found.

2.2.3.2.2 Traffic Engineering

Traffic Engineering (TE) is concerned with "performance optimisation of operational networks" and "encompasses the application of technology and scientific principles to the measurement, modelling, characterisation, and control of network traffic" [24]. Its goal is to apply different proposed techniques in order to achieve certain performance objectives.

Different TE solutions for providing certain QoS were proposed and several bodies have shown interest towards their standardisation. Among the best known are IETF's working groups that have proposed Integrated Services (intserv)²³, Differentiated Services (diffserv)²⁴ and Multiprotocol Label Switching (mpls)²⁵ which will be discussed next. Also ISO and IEEE have standardised the IEEE 802.1p eight-level priority tag-based solution in ISO/IEC 15802-3 [25] and

²² Gordon Moore, "Moore's Law", Intel, http://www.intel.com/research/silicon/mooreslaw.htm

²³ IETF Integrated Services Working Group, http://www.ietf.org/html.charters/OLD/intserv-charter.html

²⁴ IETF Differentiated Services Working Group, http://www.ietf.org/html.charters/OLD/diffserv-charter.html

²⁵ IETF Multiprotocol Label Switching Working Group, http://www.ietf.org/html.charters/mpls-charter.html

IEEE P802.1D [26] respectively and ITU is working towards standardisation of its ITU-T R. Y.1541 [27].

Integrated Services (IntServ) - RFC 1633 [28] is a reservation-based mechanism that reserves resources explicitly for individual flows using a dynamic signalling protocol and employs admission control, packet classification, and scheduling to achieve the desired QoS.

There are two services defined in this model: i) *Guaranteed Service* - RFC 2212 [29] offers quantifiable firm delay limits to flows and ii) *Controlled Load Service* - RFC 2211 [30] offers delay and packet loss like in a light loaded "best-effort" network. IntServ requires state information to be saved in each router along the path in order to ensure QoS guarantees. Usually, but not compulsory, IntServ uses Resource ReSerVation Protocol (RSVP) - RFC 2205 [31] for signaling.

Signaling, processing power, the need for storing per flow information in each participating node and possibility of unauthorized reservations lead to complexity, scalability and security concerns of IntServ applicability - RFC 2208 [32].

Differentiated Services (DiffServ) - RFC 2475 [33] is a reservation-less based framework introduced to overcome some of the IntServ limitations while provides certain QoS. In order to solve the scalability problem, DiffServ does not differentiate the traffic per flow, but defines a small number of classes for which differentiated services are provided. It divides the network in DiffServ domains (DSD) that consist of nodes that support a common policy and requires state awareness only in edges of such domains. At the edge, packets are classified into flows and marked accordingly in order to ensure their differentiated treatment. Then the flows are aggregated and sent across the DSD cloud. DiffServ Codepoints (DSCP) identify classes and their per-hop behaviours (PHB) and they are set in packet headers (DS-field that consists of six bits of the former ToS byte of the IP header) - RFC 2474 [34]. The PHB determines the forwarding to be applied to the packet in each node of the DSD. The mapping between DSCPs and PHBs depends of the DSD and is not always 1:1.

Three important PHB are: i) Class Selector PHB - RFC 2474 [34] uses the IP precedence field to indicate relative forwarding priorities. ii) Expedited Forwarding (EF) PHB - RFC 2598 [35] guarantees that packets will have a well-defined minimum departure rate which, if not exceeded, make the associated queues empty. This intends to support services that offer tightly bounded loss, delay and delay variation. iii) Assured Forwarding (AF) PHB - RFC 2597 [36] offers different levels of forwarding assurances for packets belonging to an aggregated flow. Each AF

group is independently allocated forwarding resources and their corresponding packets are marked with one of three-drop precedence. Those with the highest drop precedence are dropped with lower probability than those marked with the lowest drop precedence.

Multiprotocol Label Switching (MPLS) - RFC 3031 [37] is a strategy for streamlining the backbone transport of IP packets across a network [38]. In MPLS routing, the assignment of particular packets to classes is done just once, as the packets enter the MPLS network. This is unlike in the conventional IP routing when each packet is sent by each router along the path to the next hop after two functions were performed. First all the packets were assigned to Forwarding Equivalence Classes (FEC), that define a certain forwarding manner (e.g. same path, same treatment, etc.) and then each FEC is mapped into a next hop.

At the edge of the MPLS network Label Switched Router-s (LSR) analyse the IP headers to determine the desired service levels and the addressing information. Then 32-bit (4-byte) labels are distributed by a dynamic Label Distribution Protocol (LDP) and added to the IP packets. These labels allow the LSRs to forward the packets along predetermined paths named Label Switched Path-s (LSP) according to, for example specified QoS levels. The forwarding is performed very efficiently along the LSPs by LSRs since the forwarding engines look only at the labels and not at the entire packet headers. The labels are removed when the packets are leaving the MPLS network. The LSPs can be set up in a variety of ways [39] for example the path could represent the normal destination-based routing path, a policy-based explicit route, or a reservation-based flow path. MPLS also permits explicit routing, where the hops a packet will take are specified in advance and the label is used to indicate this route. Explicit routing is a useful capability for allowing QOS and enabling network managers to set up defined paths through the MPLS network that apply to certain traffic streams. This is when DiffServ could be used in conjunction with MPLS to provide certain QoS. Therefore even if MPLS and DiffServ are perceived as rivals, they are in fact complementary to each other.

TE-based solutions help directly or indirectly for providing certain QoS, but they have also limitations. There are significant concerns regarding the complexity of the solutions, some security issues, size of the targeted networks, reaction in really congested conditions and deployment costs. These concerns have to be traded carefully against the advantages the solutions provide in order to really benefit from providing certain QoS.

2.2.3.2.3 QoS Architectures

QoS Architectures are integrated models that include both end-systems and networks and offer QoS support for a wide range of services, including multimedia applications [40, 41]. Different QoS architecture frameworks have been proposed with various goals but similar motivations. The ideas behind some of the most important ones are presented next for an overview on different alternatives for QoS architectural support.

The Lancaster's QoS-Architecture (QoS-A) [42, 43] is "a layered architecture of services and mechanisms for QoS management and control of continuous media flows in multiservice networks" [42] that coherently apply different QoS concepts across all architectural layers and integrate them into a complete framework.

Looking from a functional point of view the QoS-A is composed of a number of layers and planes. Distributed systems-related issues are addressed by the two highest layers: the distributedapplications platform that provides services for multimedia communications and QoS specification in an object-based environment and the orchestration layer which provides jitter correction and multimedia synchronization services across multiple related application flows [44]. End-to-end related problems are dealt with at the transport layer, which contains a range of QoS-configurable services and mechanisms. Lower layers-related issues are solved by the network, data link and physical layers that offer the basis for end-to-end QoS support. QoS management is realised in three vertical planes in the QoS-A. The protocol plane, which consists of distinct user and control subplanes, divides the protocol profiles for the control and media components of flows because of their different QoS requirements. The QoS maintenance plane contains a number of layer-specific Quality of Service Managers (QM). These are responsible for the fine-grained monitoring and maintenance of their associated protocol entities, at each layer. For example, at the orchestration layer [45], the QM is interested in the tightness of synchronization between multiple related flows, whereas the transport QM is concerned with intra-flow QoS such as bandwidth, loss, jitter and delay. The flow management plane is responsible for flow establishment (including end-to-end admission control, QoS-based routing and resource reservation), QoS mapping (which translates QoS representations between layers) and QoS scaling (which constitutes QoS filtering and QoS adaptation for coarse-grained QoS maintenance control).

The OSI QoS Framework Model is based on an early contribution [46] and a long-term effort and was standardised by both ITU in ITU-T R. X.641 [21] and ISO in ISO/IEC 13236 [22]. The model defines the architectural principles, the concepts and the structures that underlie the

provision of QoS, but does not specify any of the QoS parameters or QoS information that are exchanged during the functionality. It relies on the concepts of the OSI Basic Reference Model, and the OSI Management Framework and is built around QoS-related key concepts like: QoS requirements, QoS characteristics, QoS categories, and QoS management functions. The management of QoS is performed by entities and two classes of entities are defined: system QoS entities (entities which have a system-wide role) and layer QoS entities (entities associated with the operation of a particular subsystem). The system QoS entities coordinate the response to the requirements imposed on the system and interact with layer QoS entities to monitor and control the performance of the system. The layer QOS entities implement direct control of other entities (e.g. protocol entities, etc.) that are necessary for support the system's QoS-related activities. The collaboration of layer QoS entities will, in real open systems, typically be supported by stored information and processing functions that are not specific to individual OSI layers. These information and functions following are not modelled as entities in open systems but are left to be determined by implementation choice.

The Tenet Approach [47] is a real-time communication services model with emphasis on network support for continuous media applications proposed at University of California at Berkeley and the International Computer Science Institute - Berkeley. In this approach the main elements are performance guarantees (mathematically provable, but not necessarily deterministic), contractual relationships between client and server, parameterised user-network interfaces with multiple traffic and QoS bounds and large heterogeneous packet-switching networking environments. The key mechanisms Tenet relies on are: connection-oriented communication, per-connection admission control, channel rate control and priority scheduling.

Heidelberg HeiProjects [48] at IBM's European Networking Center in Heidelberg have provided a distributed multimedia platform that includes a comprehensive QoS model that offers guarantees in the end-system and network capabilities. The model includes *HeiTS* (the Heidelberg Transport System) for transporting multimedia streams across the network [49] and *HeiRAT* (the Heidelberg Resource Administration Technique) [50]. HeiTS provides the ability to exchange streams of continuous-media data with QoS guarantees. In order to do this HeiTS makes use of both protocols for transport, network, and data link layers and components for resource management, buffer management, and operating system abstraction. HeiRAT manages all the resources, on a path from source to destination(s), both in the local systems and the network, making use of admission control, resource reservation and scheduling mechanisms. It offers two types of QoS: *guaranteed* and *statistical*. For guaranteed QoS, the resources are reserved for the maximum demand, whereas

for the statistical QoS they are slightly overbooked. Applications are allowed to specify *QoS* requirements in terms of maximum end-to-end delay, minimum throughput needed, and reliability class (loss-related) values expressed from desired to the-worst-acceptable values and HeiRAT will answer with the best QoS it could guarantee.

Among other QoS architectures proposed are the Extended Integrated Reference Model (XRM) [51] at Columbia University, OMEGA [52] at University of Pennsylvania, TINA QoS Framework [53], NU-NET [54] and NetWorld [55] at University of Pittsburg, Server/Broker/Client System at Carleton University - Canada [56] and QoS Framework [57] at Swiss Federal Institute of Technology (EPFL). More detailed information about QoS architectures, including a comparison of the presented approaches is given in [40, 58]. However designing any complex QoS architecture involves sustained effort and implementing and deploying it require significant costs. Prior to using any QoS architecture it is very important to balance its benefits for delivering different services or a particular service and the associated effort and to compare the outcome with the one if other solutions (e.g. bandwidth over-provisioning, adaptive applications, etc.) are employed.

2.2.3.2.4 Application-Level Adaptive Solutions

Another approach that tries to provide certain QoS, although without any guarantees, is by using application-level adaptive solutions. Named also application-layer QoS control-based solutions, their goal is to avoid congestion and consequent packet loss and maximise QoS [59]. This is realised by adjusting the bandwidth used by the applications according to the existing network conditions without any QoS-related support from the networks. Extensive research, mainly interested in multimedia deliveries over "best-effort" networks, has tried to explore different directions in order to offer best algorithms and mechanisms that would allow to achieve high adaptiveness and responsiveness to network conditions and high quality for the provided services. The design alternatives explored differ on how some important issues are taken into consideration. Some of these issues are:

- Signalling or feedback mechanism used to inform the applications about the current network state
- Specific adaptive mechanisms used in response to this information
- Localisation of the adjustment mechanism
- Responsiveness of the congestion control scheme in detecting and reacting to network conditions

- Capability of the scheme to accommodate heterogeneous receivers that may differ in their connectivity to the network, the amount of traffic to their delivery paths, their need for quality
- Scalability of the control mechanism to a high number of receivers
- Sharing of bandwidth with competing traffic of different type (particularly with TCP)
- Perceived quality of adapted multimedia streams.

Each of these could constitute a good starting point for a categorisation of the existing approaches and especially those used for multimedia adaptation. In [60] the authors have chosen the localisation of the adjustment mechanism in response to information about the network delivery conditions and they distinguish sender-driven, transcoder-based and received-driven adaptations. In [59] the application-layer QoS control-based adaptive solutions are divided into congestion control-based solutions, which involve congestion control mechanisms that help reducing packet delays and loss rates and error control-related solutions, which help increasing the robustness to errors, recovering after errors or minimising the effect of errors, mainly caused by transmission.

Since these application-level adaptive solutions require no help from the networks' point of view, the efforts required by these solutions for deployment and exploitation are in general very low. It is also noteworthy that upgrades are much easier to be performed. Also a very low intrusiveness, due to the fact that the networks are use as they are, is very important, because the designed adaptive solutions can be deployed in networks that are owned by third parties, increasing the generality of the deployment. As it is the case for the other solutions for providing QoS in order for the application-level adaptive solutions to become effective, they have to be deployed in networks with potential for congestion, otherwise the effort is not paid off. Also similar to other QoS solutions if the networks become really congested and almost 100% of packets sent is lost, they become ineffective. However their tendency to "back-off" would increase the chances of a quick recovery for the network and such congested situations are generally avoided. The most important limitation of these adaptive solutions is that they cannot guarantee any QoS level, although they try to maximize it.

Taking into consideration these comments that highlight the advantages and disadvantages of these solutions and underscoring that their deployment and their operational costs are very low while still providing certain QoS, the research this thesis is focused on deals with the application-level adaptive solutions. More details about related research will be given in section 2.4.

2.2.3.3 Assessing QoS

Looking at the process of providing diverse services to the customers via a broadband IP multi-service network, in order to assess its overall QoS it is necessary to analyse the QoS from different points of view.

Looking at QoS from a traditional - network engineering - point of view the work of the IETF's IP Performance Metrics (IPPM) Working Group (WG)²⁶ [61] is significant. It has proposed a set of standard metrics that can be applied to the quality, performance and reliability of data delivery over networks. This set defined in the IPPM WG-proposed RFC 2330 [62] offer some solutions for unbiased quantitative measures of network performance. These metrics are connectivity, one-way delay, round-trip delay, delay variation, loss rate, loss pattern, packet reordering, bulk and link transfer capacity. As one could see they describe the network performance, but are not directly related to the quality of the service provided. Also they significantly depend on the type of this service. However they can be used in order to assess the network condition and suggest measures to be taken by an eventual QoS-aware mechanism for certain application domain. In the context of adaptive multimedia applications some related solutions are presented in chapter two. General recommendations and not hard limits are given in ITU-T R X.641 [21] and ISO/IEC 13236 [22] for values of these network-related parameters in relation with different types of traffic such as:

- bulk data high throughput, low error rate;
- interactive low delay, low error rate;
- isochronous high throughput, constant delay;
- time sensitive constant delay, fixed throughput.

In relation to QoS, the ITU-T R. E.800 [15] defines "network performance" as a set of parameters that are meaningful to the network provider, but are expressed in terms that can be easier related to users' QoS expectations. Among the defined terms are: service support performance, service accessibility performance, mean service access delay, service integrity performance, time between interruptions, interruption duration, reliability performance, maintainability, bit error ratio, transmission performance, primary failure, execution error. However, although they are close

²⁶ IETF IP Performance Metrics Working Group, http://www.ietf.org/html.charters/ippm-charter.html

related to the users' QoS, it is difficult to directly quantify the effect of all these parameters or of some of them on end-user perceived quality.

Looking at QoS from the end-user perspective, an ideal assessment of QoS would be a totally objective one that would use trustworthy, general accepted metrics. However, first a problem would be the fact the quality assessment greatly differs with the service provided and special metrics would have to be provided for each service. Another difficulty comes from the fact that the service users are very subjective by nature and the metrics have also to take this into account. A third source of problems is the type of the content the service is providing which varies significantly due to encoding scheme, nature of the content or other issues. This makes even more difficult the efforts that aim at the determination of an objective metric for assessment of end-user perceived QoS. Therefore research is still on-going in this domain and different approaches are proposed with various advantages and disadvantages. Current state-of-the-art in objective video quality assessment is presented in detail in section 2.5.

Since there is not a general accepted objective metric, in order to assess and to compare the QoS as provided by the existing systems and the newly proposed delivery schemes, efforts were made to define subjective quality metrics and methodologies for accurate measurements. For example ITU-T R. P.910 [63] presents recommendations about methods, systems, clip contents and environment conditions for testing and scales for assessing the end-user perceived quality while viewing multimedia clips. Similarly ITU-T R. P.800 [64] presents recommendations about conditions, systems, content of the sequences, noise levels, methods, assessment scales and data analysis related to subjective testing of audio content. The former standard is used for the subjective assessment of the quality of the VoD services provided, using the approach proposed and presented in this thesis.

2.3 Compression Techniques

One of the major problems associated with storing and transmitting of multimedia-related digital data is that the huge volume of uncompressed data may easily overwhelm the available communication channels and storage systems. For example, a digital video sequence that has a resolution comparable to the National Television System Committee (NTSC) analog video signal (720 x 486 pixels/frame, 30 frames/second and 16 bits/pixel), has an uncompressed data rate of 168 Mb/s that simply cannot be coped with for both transmission and storage (a typical two-hour movie would require approximately 150 GB of disk space). This did not take into account the extra

bandwidth and consequent storage space required by one or more audio components associated with this video sequence.

Therefore compression is necessary and different algorithms and techniques were developed, having different performances according to the requirements of the applications that use them. In [23] two types of applications are distinguished: *dialogue* (interactive) and *retrieval* (non-interactive) and therefore two sets of different requirements that differ mainly in timing-related and interactivity-related issues.

The compression solutions are also subject to certain constraints. Some of the most significant constraints are:

- The *quality* of the multimedia data reconstructed after decoding should be as good as possible in order to offer high quality of the services to the end-users.
- The *compression rate* should be as high as possible in order to minimise the storage space and/or the bandwidth for transmission
- The *complexity* of the technique should be minimal to allow for a cost-effective implementation
- The *delay* due to the coding and mainly decoding should be as short as possible not to interfere with time sensitive applications.

All modern methods used for multimedia data compression are compromises between the degrees in which they follow the requirements and respect the constraints.

In order to perform their tasks, the compression methods take into account some important facts related to sets of data and to multimedia data in special:

- some subsets of data are randomly repeated within a set of data
- some multimedia-related data is more significant than other from the human perception point of view
- there is very much redundancy in the set of multimedia data, spatial (between the pixels of the same frame), spectral (within the color components of the same frame) and temporal (between different frames).

The authors of [65] divide the basic compression techniques into *lossless* and *lossy* methods, whereas in [23] they are classified as *entropy-coding* and respectively *source-coding* based. The *hybrid-coding* techniques are more complex and make use of a combination of the above-mentioned elementary methods.

2.3.1 Entropy-Coding (Lossless) Techniques

The entropy-coding methods are lossless techniques since the data set obtained after decoding is identical with the one that has been used for encoding. These methods do not take into account any specific characteristics of the streams to be encoded, ignoring the semantics of the data. They consider the data streams to be compressed as simple sequences of bits and base their operation on the observation that many sets of data, and especially audio and video streams' data, often contain sequences of identical bytes (symbols). Apart from their usage in hybrid video and audio compression techniques, the entropy-coding methods are also used for compressing data in file systems and still images [23]. Some of the best-known and widely used entropy-coding techniques are: run-length coding, Huffman coding and arithmetic coding.

Run-Length Coding's main idea is to replace the repeating symbols with the pattern that is repeated and with the number of times this happens and to signalise this with a special flag that does not constitute a part of the stream. For Run-Length Coding to be really efficient, the data stream must contain long sequences of identical characters. However, hybrid coding could employ other techniques in an earlier phase that produce such long runs and then by using Run-Length Coding very significant compression will be achieved.

Huffman Coding is a variable-length encoding technique that makes use of the occurrence probability of repeating symbols in order to produce an optimal code by assigning the shortest bits to the most frequently occurring symbol. This code is built using a bottom-up approach in a tree-like structure whose leaves are the symbols to be encoded. Huffman Coding requires the same tree (table) to be available for both encoding and decoding in order to decode correctly the compressed set of data.

Arithmetic Coding, proposed by IBM researchers in [66], is another variable-length encoding method that encodes symbols using a non-integer numbers of bits per codeword. Unlike Huffman Coding, Arithmetic coding does not encode each symbol of a set of data separately, but computes a code representing the entire set of data, achieving better performance. A significant disadvantage is that an encoded data stream must always be read from the beginning, making the

random access difficult. However, Arithmetic Coding is patented and its use is not free whereas Huffman coding is.

2.3.2 Lossy Techniques

The lossy compression techniques introduce a one-way relation between the original set of data and the decoded data set, which is similar, but not identical. They take into account the semantics of the data and in consequence the degree of achievable compression depends on the data contents. Good techniques make extensive use of the characteristics of the streams (e.g. spatial and temporal redundancies in multimedia streams). Some of the best-known lossy techniques employed in multimedia compression are *transform-based* (e.g. Discrete Cosine Transform, Fourier Fast Transform Wavelets Transform), *prediction-based* (e.g. Differential Pulse Code Modulation, Delta Modulation), *layered coding-based* (e.g. sub-sampling, sub-band coding) and *vector quantization*.

Transform-based techniques are based on the observation that if the set of data is represented into another mathematical domain is more suitable for processing. However, a very important condition is that the inverse transformation must also exist. The most used transformations in multimedia-related compression are Discrete Cosine Transform (DCT) and Fast Fourier Transform (FFT), although lately Wavelets-based transform is also used for specific applications. For example, DCT is applied in image compression on a N x N image block transforming the data from spatial domain into Discrete Cosine (DC) domain, resulting N² DC coefficients. Regularly N = 8, which ensures low memory requirements, low computational complexity and high spatial correlation of the neighborhood related to the pixel in the center. It was proven that the higher the order of the DC coefficients, the more sensitive their influence on the human visual system is. Therefore in order to both reduce the data quantity and have as little as possible influence on the perceived quality, another step called quantisation may delete some of the low-order coefficients. The inverse DCT restores the data into the spatial domain. FFT transforms data into the frequency domain in which either the complexity of the computation is lower or information easier available. For example the audio-based compression uses 512 or 1024-point FFT for getting detailed information about the spectrum of the original signal. Based on the psychoacoustic model of the human ear it is decided which of these samples has a lower impact on the quality of the overall stream and by masking them, the quantity of data is reduced. Unlike DCT or FFT that are applied on homogeneous sets of data, wavelets are transforms characterised by strong locality and could be very successfully applied to determine local specifics of signals or

images. If applied on the whole image, other techniques such as quantisation and entropy encoding have to be used in conjunction in order to achieve compression. [65].

Prediction-based methods, also known as differential or relative encoding techniques, are based on the idea that if a set of values are clearly different from zero and these values do not differ much, encoding the differences from the previous values leads to compression. These methods are best suited for encoding audio data because audio signals change rather slowly. Differential Pulse Code Modulation (DPCM) and its variation Delta Modulation (DM) use this approach. DM uses only one bit to indicate whether the new value increases or decreases from the previous one, achieving high compression, but lower accuracy in case that high variations occur.

Layered-coding-based techniques consider that the data to be compressed can be divided in different layers that could be treated differently, according to their importance. For example in video compression Sub-sampling takes into account the fact that the human eye is more sensitive to differences in brightness than in color and therefore it divides the images in YUV components (i.e. luminance Y and two chrominance difference components U and V), instead of using the RGB components (i.e. Red, Green and Blue)²⁷. The real benefit is achieved by another step that compresses differently these components. Sub-band encoding, which is mainly used in audio compression, divides the frequency spectrum into pre-defined bands. Different quanitisation processes are then used, finer for more audible bands (e.g. between 100 Hz and 16000 Hz) and coarser for the rest.

Vector Quantisation is an asymmetric compression method having the decoding process much simpler than the encoding one. It achieves very good compression and could be performed quite fast, working directly in the spatial domain. Image compression is achieved by dividing the input image in non-overlapping $N \times N$ blocks, seen as N^2 -dimensional vectors, and matching each of them to a codeword from a codebook, so that the distortion between them is minimum.

2.3.3 Hybrid Techniques

Hybrid techniques make use of a number of lossless and lossy compression techniques in conjunction in order to achieve better data compression. Next some of the best known standardised encoding schemes and some proprietary solutions are presented. Since video accounts for the large

²⁷ Y=0.30*R+0.59*G+0.11*B;U=B-Y; V=R-Y

majority of data to be stored and transmitted, this section will give more details about hybrid video compression techniques, but it will also present briefly the ideas behind image and audio-related ones.

2.3.3.1 The JPEG Standards

The **JPEG**, standardised in ITU-T R. T.81 [67] and ISO/IEC 10918-1 [68], is an image compression hybrid technique that offers the flexibility to either select very high picture quality with low compression ratio or a very high compression ratio with low picture quality. The latter is caused mainly by "blockiness" artifacts. To achieve compression JPEG employs DCT followed by a quantization phase. The **motion-JPEG** (M-JPEG), an extension to JPEG standard for video, uses a series of still pictures and achieve low compression by not reducing any temporal redundancy.

The **JPEG2000**, standardised in ITU-T R. T.800 [69] and ISO/IEC 15444-1 [70], uses Wavelets-based transformation instead of JPEG's DCT, increases the compression ratio, but also the complexity and reduces the "blockiness" artifacts, but produces a slight "fuzzy" picture. The **motion-JPEG2000** (M-JPEG2000), although standardised in ISO/IEC 15444-3 [71], suffers from the same problems as M-JPEG.

2.3.3.2 The MPEG Standards

The goal of the ISO/IEC Moving Pictures Expert Group (MPEG)²⁸ [72] was to develop international standards for compression, decompression, processing and coded representation of moving pictures, audio and their combination.

MPEG-1 [73], the first standard generated by this group, defines the coding of the combined audio-visual signal at a bit-rate around 1.5 Mbps with VHS-quality (320 x 240 video resolution) and was initially developed to operate from storage media, but it can be used more widely than this. In different parts of the ISO/IEC 11172 document [73] the video, the audio and the system components of the standard are described.

MPEG-1 Video uses a number of lossy and lossless compression techniques in order to achieve high compression ratios while still providing good quality for the decoded video stream. First, an appropriate spatial resolution is selected and then the algorithm uses block-based motion

²⁸ ISO/IEC Moving Pictures Expert Group (MPEG), http://www.chiariglione.org/mpeg/index.htm

compensation to reduce the temporal redundancy. Motion compensation is used for causal prediction of the current picture from a previous picture, for non-causal prediction of the current picture from a future one, or for interpolation-based prediction from past and future pictures. The difference signal (prediction error) is further compressed using DCT to remove spatial correlation and is then quantised. Finally, the motion vectors are combined with the DCT information, and coded using variable length codes. *MPEG-1 Audio* uses filters and sub-sampling to map the input audio stream into a representation in the frequency domain and a psychoacoustic model that creates a set of data to control the quantisation and the coding processes. These processes create a set of coding symbols from the mapped input samples. The coded bitstream is then obtained from the output data and other information (e.g. error correction) if necessary. *MPEG-1 System* combines one or more MPEG video and audio streams, with timing information, to form a single stream well suited to digital storage and/or transmission.

MPEG-2 [74] was standardised by both ISO in ISO/IEC 13818 [74] and ITU in ITU-T H.262 [75] and defines the following components:

MPEG-2 Video, although similar to MPEG-1 Video, is targeting very high bit-rates of up to 20 Mb/s with full size pictures and very high quality. It also allows for higher flexibility in terms of applications, bitrates, resolutions and qualities with the introduction of "profiles" that define subsets of the MPEG-2 syntax and semantics and within each profile of "levels" that describe a set of constraints imposed on parameters in the stream. MPEG-2 Audio is very similar to MPEG-1 Audio having added multi-channel extensions. However, MPEG audio is backward and forwards compatible. MPEG-2 Program is similar to MPEG-1 System and aims at combining one or more elementary streams, which have a common time base, into a single stream. MPEG-2 Program is designed for use in relatively error-free environments and is suitable for applications which may involve software processing. Program stream packets may have variable and large size. MPEG-2 Transport combines one or more elementary with one or more independent time bases into a single stream. The Transport Stream is designed for use in environments where errors are likely, such as storage or transmission in lossy or noisy media. Transport stream packets are 188 bytes long.

MPEG-4, standardised as ISO/IEC 14496 [76], targets the extremes from the bitrate range point of view to the world of possible applications. It provides features like extended scalability, error resilience, interfaces to digital rights management systems and interactivity. It aims to achieve robustness in any kind of environment, compression efficiency coding, allow for transmission flexibility and provide support for objects with both natural and synthetic content.

It is significant to mention that a new scalable coding mechanism, different than classic quality (SNR), spatial and temporal scalability, called **fine granularity scalability (FGS)** was proposed for MPEG-4 and was described in [77]. The idea is that each stream should consist of a base layer ("must have") and an enhancement layer. Parts of the latter could be optionally transmitted to increase the overall quality if available bandwidth permits this. **Progressive fine granularity scalability (PFGS)** [78] extends the FGS idea, allowing for the existing of more than two layers.

2.3.3.3 The ITU-T Standards

The ITU-T has defined in the ITU-T R. H.320 [79] a standard for multimedia telecommunications that ensures compatibility among terminals produced by different vendors. It specifies certain standard protocols for video, audio, control, security etc. and provides mandatory requirements to make sure all H.320 compatible systems can communicate with one another. There are also optional requirements that can allow systems to provide additional functionality. However, this functionality is sacrificed for compatibility when communicating with systems that only meet the minimal requirements for H.320. The standards directly related to video and audio compression are presented next.

ITU-T R. H.261 [80] is a video-coding standard, designed originally to suit ISDN lines, that has output bit rates multiples of 64 Kb/s, between 40 Kb/s and 2 Mb/s. The encoding algorithm employed is a mixture of temporal and spatial coding to remove the redundancies in the video in a similar fashion MPEG does. However, H.261 offers lower compression ratio and provides lower flexibility in exchange for lower processing delays that may be required in videoconferencing. This is because H.261 was targeted at teleconferencing and videophone applications.

ITU-T R. H.262 [75] is common with MPEG-2 standard ISO/IEC 13818 [74].

ITU-T R. H.263 [81] is a video encoding standard that was originally designed for low bitrate communication, less than 64 Kb/s, a limitation that has now been removed. The coding algorithm of H.263 is similar to that used by H.261, but there are some changes that improve its performance and flexibility. Among other features H.263 supports five standard source formats instead of two and uses half pixel precision for motion compensation rather than full pixel. It also makes use of 3-D variable-length coding and median motion vector prediction. H.263 also offers a wide variety of optional modes that can be added to the baseline algorithm to improve the coding

performance or to broaden the application range. Further improvements have been proposed and have lead to H.323+ and H.323++.

Lately significant effort is directed towards the emerging standard "Advanced Video Coding (AVC)", widely known by its working title, H.26L or by the ITU-T document number H.264 [82].

ITU-T R. G.711 [83] provides telephone quality sound at rates between 48 and 64 Kb/s and is the only audio protocol required for a system to be H.320 compliant.

ITU-T R. G.722 [84] provides stereo quality sound with output between 48 and 64 Kb/s.

ITU-T R. G.728 [85] is meant to be used videoconferences at speeds lower than 256 Kb/s and requires only 16 Kb/s, allowing more bandwidth for video.

2.3.3.4 Proprietary Solutions

Different commercial companies have proposed hybrid compression solutions that make use of both lossy and lossless techniques in conjunction in order to maximise the benefit from their usage according to the companies' interests. Unfortunately majority of these solutions are proprietary and very little information is offered about them. This makes their usage outside the multimedia systems they were initially designed for extremely difficult. Among the best-known proprietary solutions are the Microsoft's Windows Media (WM)²⁹, Progressive Networks' Real Media (RM)³⁰ and Apple's QuickTime³¹.

2.3.4 Conclusion

Different basic and hybrid methods proposed for multimedia data compression were presented. The latter, using a combination of the former, achieve higher compression ratios while providing high quality for the reconstructed multimedia data and are best suited for using while delivering multimedia presentations to the residential homes. However, since the goal of this research is to deliver very high quality multimedia streams with very little effort, the chosen

²⁹ Microsoft, "Advanced Systems Format Specification", http://www.microsoft.com/windows/windowsmedia/format/asfspec.aspx

³⁰ Progressive Networks, http://www.realnetworks.com

³¹ Apple, QuickTime, http://www.apple.com/quicktime/

solution has to provide very good compression ratios, to be designed for very high quality streams and to be standardised and therefore offering good documentation and reducing the maintenance costs. MPEG-2 was selected because it was standardised by both ITU-T and ISO bodies, it is already popular through the DVDs that use it for storing very high quality multimedia, it is supported by a wide scale of software and hardware products and it is already used by current VoD and multimedia broadcast providers. The latter increases the chances for the proposed application-level adaptive solution to be accepted and used in existing systems, that requires only incremental changes.

2.4 Adaptive Solutions for Delivering Multimedia

Bursty loss and excessive and extremely variable delays have a devastating effect on multimedia deliveries, severely affecting the end-users' perceived quality. In consequence any effort that aims at reducing these delays and at lowering the loss rate helps to increase the quality of the remote multimedia presentations. This is also the main objective of the adaptive solutions (or adaptive control schemes) for multimedia deliveries.

Extensive research has focused on proposing different solutions for adaptive multimedia streaming and various directions have been taken. These directions differ due to a number of options taken when designing the adaptive solutions. Among the most significant are the manner the information about the delivery conditions is collected and used, who takes the adaptive decisions and what adjustments are performed, how fast and how appropriate are the adaptive measures taken and what is their effect on the end-user perceived quality. These are the characteristics that will be presented in relation to the proposed approaches.

Adaptive control schemes have been mainly classified in the literature [59, 60] according to the place the adaptive decision is taken and this thesis uses also this approach. The source-based adaptive control techniques require the sender to respond to variations in the delivery conditions or to changes in the quality of the reception. The receiver-based adaptive control schemes provide mechanisms that allow for the receivers to select the service quality and/or rate. The hybrid adaptive control mechanisms involve both the sender and the receiver in the adaptation. However the authors of [60] distinguish another category - the transcoder-based adaptive control solutions - that focus on matching the available bandwidth of heterogeneous receivers through transcoding or filtering. These schemes imply an active involvement in the adaptation process at the level of intermediary network nodes. Although adding intelligence to the network introduces supplementary

costs and has to have the acceptance of the network's owner or the administrator, which does not conform to the goal of this research, this approach is also presented along some other very interesting solutions.

There are also works like [86] that have considered other criteria to classify the adaptive schemes such as whether they are unicast or multicast, single-rate or multi-rate, end-to-end or router-supported, TCP-friendly or not. Although complex, unicast solutions that look for the adaptation of the delivery to one receiver at a time are far less difficult to design than multicast ones. The problem arises for multicast schemes from the fact that they have to scale to large number of receivers, often heterogeneous from both network and capabilities point of view. If there is a common transmission rate for all the receivers, there is very difficult to determine how to adjust it according to the available information about delivery, since for example these receivers may experience uncorrelated loss. Different problems related to multicast congestion control are presented in detail in [87, 88]. Single-rate solutions involve the data transmission at the same rate for all the receivers and this is the case for all unicast schemes and for some multicast solutions. Multi-rate schemes do not restrict the transmission rate to that of a bottleneck receiver, allowing for more flexibility. They require the existence of more than one multicast group and provide the choice for the receivers to join or leave the multicast groups according to their particular delivery conditions. However the latency of the process of leaving a multicast group is a reason of concern and it may take several seconds to complete. The end-to-end approach, chosen also in this thesis, is designed for best-effort IP networks and does not rely on any support from the networks that are taken as they are. Its biggest advantage is the low cost that makes it popular. The router-supported solutions rely on additional functionality from the networks and some proposals were presented in the first chapter. The schemes' positive results come with increases in the costs of their deployment. A more detailed discussion about the advantages and disadvantages of these solutions is published in [89]. Another issue is the solutions' degree of TCP friendliness, which measures the effect the adaptively controlled flow has on competing TCP flows. Although there is not a general accepted definition for TCP friendliness and even a general agreement for the necessity of strong TCP friendliness for time-sensitive flows, certain "social behaviour" from the solution is definitely required to allow for other traffic to have its share of the bandwidth, especially during increased traffic conditions. Definitions of TCP friendliness are given in [6, 86].

Apart from serving live content, which can be more flexible encoded at the required bitrate, in order to be able to provide adaptively on-demand multimedia services, including VoD, solutions for distributing pre-recorded multimedia streams are required. In the literature [90] there are

suggested several ways of providing quality adaptation for a pre-recorded stream, including adaptive encoding, switching among multiple pre-encoded versions and hierarchical encoding. Adaptive encoding involves re-encoding of the existing content on-the-fly, based on the available bandwidth for transmission. However, this is computationally complex and has high CPU and memory requirements and is very unlikely for the servers to be able to do this for a high number of clients. However the transcoder-based solutions use this approach, but do not involve the server. Switching among multiple quality pre-encoded versions with the same content require increased disk storage at the servers for these versions. However, lately the decrease in the price per gigabyte makes this solution more popular. Hierarchical/layered encoding requires the server to use layered encoding for the streams. As more bandwidth becomes available, extra layers can be delivered.

Next a review of work representative for the direction the research on adaptive control schemes have taken is presented, solutions classified according to the location where the adaptation decision is taken.

2.4.1 Source-based Adaptive Control Techniques

In the source-based adaptive techniques the sender is responsible for adapting the transmission rate to the delivery conditions. It is significant to mention that two main policies have been adopted for adjusting the rate: the additive increase and the multiplicative decrease (AIMD) that slowly increases the rate in good reported conditions and sharply decreases it otherwise and the multiplicative increase and the multiplicative decrease (MIMD) that uses roughly the same approach upwards and downwards. In relation to the adaptation approach, two explored directions have been distinguished in [59, 91]: a probe-based approach and a model-based approach. The probe-based solutions are based on probing experiments that try to detect the available bandwidth of the network and try to maintain the loss rate below a certain threshold. The model-based solutions follow a throughput model that determines the transmission rate in certain conditions. However a third direction, which could be named heuristic-based and relies on heuristic knowledge and experimental testing, encompasses many of the proposed schemes.

Kanakia, Mishra and Reibman present in [92] a heuristic-based unicast scheme that relies on periodically received feedback by the server about the bottleneck queue's buffer occupancy and service rate received by the connection. The latency of the feedback is taken into account while estimating the current buffer occupancy and the service rate. These estimates are used to calculate

the most appropriate transmission rate for each video frame, according to its type, before it is transmitted. The scheme uses MPEG encoding scheme and the adjustment in the quantity of data to be transmitted is performed by reducing the streams' quality by varying the encoder's quantisation factor (Q). The biggest problem related to this solution is that it is very difficult to expect to receive the required type of feedback from the bottleneck link router for many connections. Apart from scalability problems, not being end-to-end, the solution involves an increased deployment cost.

Jacobs and Eleftheriadis [93, 94] propose a protocol for transmitting multimedia that uses the TCP's congestion control window and hence TCP's acknowledgement messages to acquire information about the state of the network. The goal of their research is to find a TCP-friendly solution for multimedia streaming and therefore aim at adapting to network conditions in a similar manner with TCP. Before being sent to the receiver, the packets carrying multimedia data are placed into a local buffer at the sender. This buffer's occupancy is used by a dynamic rate shaping mechanism, which was described in [95] to control the encoder's output rate. The encoding rate is reduced when necessary by eliminating a set of DC coefficients using a Lagrangian optimisation.

Bolot, Turletti and Wakeman propose in [96, 97] a heuristic-based adaptation scheme that involves a server that is informed about the network conditions through feedback from the receivers. Once congestion is detected, the server varies the output rate of a H.261 encoder by adjusting the frame rate, the quantization factor and the movement detection threshold. This scheme was extended to multicast [98] and in order to reduce the load on the server, the receivers send feedback only if they were selected by a probabilistic polling and only if they experience congestion. The feedback is initiated by the server that sends probe messages with a random generated key of a length that decreases logarithmically in time. These keys have to be matched by the clients' own key in order for the latter to be allowed to answer. The decrease in the key length is performed in order to address more clients until the server receives an answer. The very good scalability of this multicast scheme has a drawback in the fact that the congestion is discovered with certain delay and until dealt with affects the quality of delivery.

Sisalem and Schulzrinne have designed the Loss-Delay based Algorithm (LDA) [99] that makes use of RTP [100] to deliver data and RTCP [100] to provide feedback. The scheme looks at the overall multimedia delivery taking into account all the LDA adaptive streaming processes. For each process, depending on whether was reported loss or not the transmission rate is either additively increased or multiplicatively decreased. The additive rate increase is performed with a parameter AIB_i whose value depends on its former values, on the value of the transmission rate and on the estimated bottleneck bandwidth. However the rate cannot exceed the rate suggested

for these conditions by the TCP model proposed in [101]. The decrease brings the transmission rate to a value equal to the current minimum rate experienced by the LDA receivers. The authors do not suggest how these variations in the rate could be performed from the multimedia encoding point of view and nor what would be their effect on the end-users' perceived performance. The scheme was tested for multicast deliveries and has shown certain TCP fairness, although no further tests have been done to prove this. The main drawback of the scheme is that it has several parameters that have to be set by the users.

Busse, Deffner and Schulzrinne have presented in [102] another adaptive scheme that bases its operation on how the sender perceives the receiver state according to feedback-received information. The receiver could be in "unloaded" state with no loss experienced and in consequence the server increases the transmission rate additively until it reaches the "loaded" state when the maximum rate is matched and the rate is not varied anymore or in "congested" state when loss is reported and the sender has to multiplicatively reduce the transmission rate until the loss rate decreases. A low-pass filter is used to smooth the reported packet loss rate and the resulting value helps the sender to decide the receiver's state. The multimedia data is transferred using RTP [100] and the feedback information using RTCP [100]. When applied to multicast, a significant problem is according to how many reports that inform about congestion the decision of setting the common rate for transmission has to be taken. A solution is to take into account the poorest receiver's report, another to consider a fraction of the total number of receivers' reports. The authors have examined both. The scheme suffers from the same problem of the dependency of loss rate only as previously mentioned.

Rejaie, Handley and Estrin have proposed in [103] the Rate Adaptation Protocol (RAP), an unicast adaptive solution that follows the TCP AIMD approach. In consequence each data packet require an acknowledgement from the receiver, according to which both the loss rate and the round-trip time (RTT) are estimated. In case that congestion is detected the transmission rate is halved, otherwise the rate is increased by one packet per RTT. RTT is also the interval between two possible decisions of rate adjustment. RAP additionally provides a fine-grained delay-based congestion avoidance mechanism based on short-term and long-term RTT averages that modify the interval between consecutively sent data packets. More information about RAP and its application in multimedia quality adaptation is given in [104].

Pahye, Kurose and Towsley have used the TCP model presented in [101] to propose the TCP-Friendly Rate Control Protocol (TFRCP) [6], a model-based solution that controls the sending rate in a similar manner to TCP. The sender computes the sending rate at the beginning of

every round of duration M units (recomputation interval), rate that is used to send data for the duration of the round. If no loss was recorded the rate is doubled, otherwise the rate is computed based on the TCP model. The data packets are acknowledged in a similar fashion with TCP, but each ACK packet gives supplementary acknowledgment information about the previous 8 data packets, protecting against ACK losses. In this manner the loss is detected and its rate computed. Also, the sender maintains estimates of the round-trip time and of the base timeout as TCP does, necessary for the rate computation. The protocol was tested for unicast transmissions and has shown high TCP-fairness. However the tests have not included any reference to eventual user perceived quality if used for delivering multimedia data and have not addressed the link utilisation.

Floyd, Handley, Padhye and Widmer have improved TFRCP [6] proposing a new TCP-Friendly Rate Control Protocol (TFRC) [105] that was designed for unicast communications, but could also be adapted for multicast. Like TFRCP, TFRC uses the same equation of the TCP model for determining the transmission rate, but uses more complex methods to determine the values of its parameters. The scheme regularly computes the loss intervals taking into account the number of packets between two consecutive losses. A weighted average is computed from a certain number of loss intervals allowing for newer loss intervals to contribute more to the result and increase its accuracy. The loss rate is measured then as the inverse of the weighted average loss interval, taking into account that it should not react strongly to single loss events and should adapt quickly to long periods of no loss. TFRC provides additional delay-based congestion avoidance by adjusting the time between two consecutive packets sent. As result of these improvements, the scheme's sending rate is more stable, while still provides high responsiveness to changing traffic conditions. Unfortunately the authors have not assessed the effect of using their proposed scheme on the endusers' perceived quality.

Sisalem and **Wolisz** have improved LDA [99] and have proposed the **Loss-Delay-based Adaptation Algorithm** (**LDA+**) in [7]. The adaptive scheme was designed for unicast transmissions and it bases its functionality on using RTP [100] for data delivery and RTCP [100] for feedback. LDA+ is an AIMD algorithm that changes its transmission rate with values dynamically computed based on the current network situation and the share of the bandwidth a flow is already utilising. In loss situations, the rate is decreased by the factor 1-lossrate ^{1/2}, but the final values should not be lower than the one the TCP model equation [101] would suggest for the transmission rate in these conditions. In no loss cases, the additive value is computed as the minimum between three values. One is computed in inverse relation with the share of the bandwidth the current flow utilises. A second value is meant to limit the increase to the bottleneck link

bandwidth as it converges to 0 when this happens. The third value is determined in such a manner that, at no time, the rate should increase faster than a TCP connection sharing the same link would increase its rate. Although performance tests were performed, they were mainly focused on tuning LDA+, determining fairness to other flows and comparing it to other schemes and the effect on the end-users was not taken into account.

Rejaie, Handley and Estrin have proposed in [90] the Layered Quality Adaptation (LQA) scheme for unicast transmissions that, based on a layered approach, provides the ability to control the level of smoothing. LQA consists of two mechanisms: a coarse-grain mechanism for adding and dropping layers and a fine-grain inter-layer bandwidth allocation mechanism. The sender performs the coarse-grain adaptation by changing the number of active layers and varying therefore both the quality of the delivered stream and the quantity of data buffered at receiver. The fine-grain adaptation is performed at the level of an active layer. For example if there is spare bandwidth, the sender could increase the rate the data is sent for an active layer, increasing the quantity of data buffered at the receiver for that layer. Later on, if there is receiver buffered data for a layer, the sender could reduce temporarily the allocated bandwidth for that layer under that layer's regular necessity by reducing its sender buffer's drainage rate. Additionally a smoothing mechanism was introduced that trades short-term improvements for long-term stability of quality.

Apart from the techniques presented, different other sender-based adaptive control schemes have been proposed, including a FGS-based solution [106], an adaptive TCP friendly scheme [107] and a mechanism based on priority drop [108], that try to find the best answer to streaming multimedia over best-effort IP networks. Unfortunately it was not reported any objective or subjective testing of the effect these solutions may have on the end-users' perceived quality and therefore it is difficult to be assessed from the existing publications. Other sender-based adaptive solutions are presented or discussed in [59, 109, 110].

2.4.2 Receiver-based Adaptive Control Schemes

The receiver-based adaptive control solutions involve the receivers as the main actors in the rate adaptation, while the senders either do not participate at this process or do not have a significant contribution. Currently they are built around the idea of layers and take advantage of it. Like the source-based rate control, the existing receiver-based adaptive control mechanisms were classified in the literature [59, 91] in two approaches: the probe-based approach and the model-based approach. The **probe-based approach** relies on adding and dropping layers. When no

congestion is detected, a receiver does the probing for the available bandwidth by joining a new layer, increasing its receiving rate. After the joining, if no loss occurs, the probing was successful, otherwise, the receiver drops the newly added layer. When congestion is detected, the receiver drops a layer, resulting in reduction of its receiving rate. The **model-based approach** attempts to explicitly estimate the available network bandwidth. Currently the only one model the solutions are based on is the throughput model of a TCP connection presented in [101]. Next the most significant receiver-based adaptive control solutions are presented.

McCanne, Jacobson and Vetterli have proposed the Receiver-driven Layered Multicast (RLM) in [111], a multicast adaptive solution especially designed to provide each receiver with the best possible video quality according to the available bandwidth between the sender and that receiver. In RLM, a probe-based approach, the sender splits the video into several layers and each is transmitted to a different multicast group. If a receiver joins the multicast group that transmits the first layer, it will receive the multimedia data associated with it. This process is named join experiment. If no packet loss will be experienced for a certain period of time, the receiver will subscribe to the next layer. When a receiver experiences packet loss, it unsubscribes from the highest layer it is currently receiving. The use of RLM to control congestion comes with many problems, some of which were reported in [112]. Among them is the coarse adaptation and the consequent variation of the end-user perceived quality when adding or dropping layers based only on the detection of packet loss. Also leaving a multicast group has certain inertia, taking time on the order of several seconds. Therefore a receiver, which has joined a higher layer immediately has to leave it, adds unjustifiable cost in terms of the additional bandwidth it may use. Furthermore, this increases with the number of receivers behind the same bottleneck link that take unsynchronised join and leave decisions.

Vicisano, Crowcroft and Rizzo address many of RLM's [111] problems when have proposed another probe-based solution, the Receiver-driven Layered Congestion Control (RLC) [113]. Their idea was such to dimension the layers that the bandwidth consumed by each new layer increases exponentially. The time that a receiver has to wait before being allowed to join a new layer also increases exponentially with each additional layer. However, a layer is dropped immediately when packet loss occurs, halving the overall receiving rate. This AIMD behavior is very similar to TCP's. Noticing that the receivers' synchronization is beneficial, synchronization points (SP) have been defined and the receivers may join a layer only at these SPs. Since the SPs are exponentially less frequent in higher layers than in lower layers, a low quality receiver is likely to catch up with receivers with a higher subscription level and, after some time, synchronization

will occur. In order to decrease the chance of a failed join experiment, the sender temporarily doubles the rate of each layer before every SP and only if a receiver does not report loss is allowed to join a higher layer. In spite of the efforts RLC has also problems [112] linked mainly by the coarse granularity of the rate adaptations, related to the fact that the transmitted data must support layering and in connection to the acceptability of the artificially introduced bursts. A general criticism that applies to all layered-based schemes is that they "abuse" of the network resources.

Byers et. al have solved some of the deficiencies of RLC and have proposed Fair Layered Increase/Decrease with Dynamic Layering (FLID-DL) [114], a model based solution.. The scheme makes use of a Digital Fountain [115] at the sender that encodes the original data and some redundancy information such that receivers can decode the data even if they receive only a certain number of distinct packets. In order to reduce the join and leave latencies associated with adding or dropping of layers, FLID-DL introduces the concept of Dynamic Layering (DL). DL involves constant decrease in time of the bandwidth associated with a layer. In consequence if a receiver wants to maintain the received quality, it has to periodically join new layers. The receive rate is reduced simply by not joining additional layers, whereas rate increase requires joining multiple layers. After a while every layer will carry no data and therefore layers are reused after a period of time of inactivity. DL is complemented by a Fair Layered Increase/Decrease (FLID) scheme that involves the receivers' subscription to additional layers only with a certain probability. These probabilities are chosen so as to achieve a rate compatible with TCP. FLID retains RLC's concepts of sender-initiated synchronization points to coordinate receivers but does not transmit packet bursts to probe for available bandwidth. FLID-DL is more flexible related to the data distribution between the layers, but involves major overhead for the underlying multicast routing protocol as more frequent join and leave decisions occur. It exhibits some rate oscillations caused by the use of TCP equation and has limited TCP-friendliness.

Other solutions have been proposed such as [116] that address some of the problems related to the other approaches, but currently FLID-DL is seen as the best existing receiver-based adaptive control scheme [59]. However no objective or subjective assessment was done in relation to the contribution towards the increase in the end-user perceived quality when streaming multimedia and therefore it is difficult to compare it to other approaches.

2.4.3 Hybrid Adaptive Control Mechanisms

The hybrid adaptive control mechanisms involve both the receivers and the senders in the adaptive control, the solutions being a combination of sender-based and receiver-based schemes. Next some of them are presented.

Sisalem and Wolisz have extended the functionality of the LDA+ [7] to multicast transmissions and have proposed Multicast LDA+ in [117]. MLDA is a typical example of hybrid congestion control mechanisms because it distributes its adaptive decisions between the sender and the receivers. MLDA uses RTCP for feedback as in LDA+, but also for signalling between the sender and the receivers. Unlike LDA+ MLDA bases its operation on layered multicast. It is very significant that although the AIMD principle for rate adaptation is maintained, the appropriate rate for each one of the receivers is computed by itself and feedback-ed to the sender. Exponentially distributed timers make sure that the sender is not inundated by feedback messages. The sender continuously adjusts the bandwidth distribution between the layers in order to support the receiver reported rates. Independently the receivers adjust their subscription level to the appropriate rate. The computation of RTT-s at the receivers was very difficult and a complex solution was found to estimate them accurately enough. The most significant benefit of this scheme is that by reducing the rate of a layer that causes congestion the adaptation is performed much faster than if the receivers are expected to leave this multicasting group. At the same time the scheme is very complex and requires further work related to the manner the data is distributed into the dynamic layers.

Rhee, Ozdemir and Yi have proposed in [118] the TCP Emulation At Receivers (TEAR) a hybrid adaptive rate control scheme that uses aspects of window-based congestion control and targets both unicast and multicast transmissions. The receivers maintain a congestion window whose size is modified in a similar manner with TCP. The main difference from TCP whose congestion window is located at the sender is that the TEAR receiver has to estimate the moments when TCP would increase or decrease this window's size. The receivers compute the transmission rate as roughly a congestion window worth of data per RTT. To avoid TCP's saw-tooth-like rate shape, TEAR averages this rate over an *epoch*, which is defined as the time between consecutive rate reduction events. To minimize the effect of noise in the loss patterns, the rate is then smoothed by weight averaging over a certain number of epochs. This value is then reported to the sender, which adjusts the sending rate. In the multicast case the TEAR sender sets the rate to the minimum of the reported rates. Although TEAR shows good TCP-friendly behavior while avoiding TCP's frequent rate changes, it is unclear what is the effect on the end-ser perceived quality if used for streaming high-quality multimedia.

Apart from the presented solutions, other mechanisms such as the destination set grouping [119] and another layered multicast scheme [120] have been proposed. However they are all very complex and mainly rely on the layered encoding which puts very much pressure on the applications that have to control the data division into these layers in order to be effective from the level of end-user quality point of view.

2.4.4 Transcoder-based Adaptive Control Solutions

The transcoder-based solutions provide a different approach than the end-to-end-based adaptation. They make use of one or more multimedia gateways placed at appropriate locations in the networks, which actively contribute to the adaptation process by modifying the bitrate to suit mainly heterogeneous receivers.

Yeadon, Garcia, Hutchison and Shepherd have proposed in [121] a multicast adaptive solution based on a QoS filtering model. A filter is a mechanism that operates within the network or at the network edge to control and/or modify some characteristics of transmitted media streams to support heterogeneous receivers in terms of their capabilities and associated QoS requirements. A significant part of their model is the filter propagation, which occurs when the levels of QoS requirements of all outputs of a node are lower than the QoS associated with the input stream. However, in order to avoid oscillations, the filters may only propagate when the difference of these levels exceeds a certain threshold. Different types of filters have been proposed such as codec, frame-dropping, color reduction, DCT-filters, mixing and splitting filters. More details about this solution can be found in [122, 123].

Amir, McCanne and Zhang propose in [124] a transcoder-based solution that, placed at the level of a multimedia gateway in the network, can convert the multimedia stream from the input format into an intermediary representation by a decoder. This intermediary representation is supposed to be encoded easily in a number of output formats supported by the transcoder. The transcoder is configured by an external control interface that can select parameters such as the input and output formats, streams' characteristics etc. However the effort involved by these encoding/decoding processes is significant and hardly can be accommodate for example at a router level where a very high number of time-sensitive operations have to be performed.

Other works have been proposed that complement or improve the already existing solutions. For example in [125] a control scheme meant to configure the transcoders from a multicast tree to support receivers with low QoS is presented, in [126, 127] faster transcoding-based

solutions are described and in [128, 129] rate control schemes for MPEG or H.263+ transcoding are proposed. However in [128] the authors have distinguished three major directions for very high quality transcoding: Cascaded Pixel Domain Transcoding (CPDT), DCT-Domain Transcoding (DDT) and Open-Loop Transcoding (OLT) [128]. CPDT involves first the total decoding of the input stream and later on its encoding at a different rate and has high computational and memory requirements that makes it inefficient for real-time streaming. Since many complex encoding schemes use DCT as a phase for achieving lossy compression, DDT and OLT decompress partly the input stream and use different level of DC coefficients requantisation.

2.4.5 Conclusions

This section has presented different solutions for adaptive control while delivering multimedia streams. Some of them were specially designed for multicast deliveries, others involve fixed or propagating filtering or transcoding deployed at network nodes' level, some use feedback from receivers to inform the sender about the network situation, others do not, some rely on layered coding, others are more general solutions.

Multicasting solutions distribute the same content or a limited set or versions to all the receivers, being very efficient. These solutions may be accepted today when non-interactive multimedia broadcasting still accounts for the large majority of multimedia based services. However, in the future the deployment of rich, high quality services require interaction with the customers, personalisation of services and extended VRC control, difficult to be provided through multicasting. Having access to the network nodes and adding intelligence to the network require the permission of the network operators and the service providers and add costs to the solution which may affect the generality of the services and final price the customers may have to pay, both influencing the chances for successful large scale deployment of these services to residential homes. A very significant criticism of all the existing solutions is that they do not directly relate their adaptive decisions to the end-user perceived quality, whose maximisation should be the goal of any adaptive control-based solutions employed in delivering multimedia-based services. Regardless of the adaptation mechanism, the results were mainly analysed in term of network-related metrics and only in very few cases have been assessed from the end-user perceived quality point of view in terms of an objective metric or after subjective testing.

All these observations explain our choice towards an unicast adaptive solution that offers one-to-one relationship with the receiver, offering high degree of flexibility, with no support from

the network. Relative to the location where the adaptive decision is taken, the server was chosen, although monitoring of the effect the delivery conditions have on the end-user quality and some related computation was distributed to the clients. This allows for relative independence from the encoding scheme (although MPEG-2 encoding is used), lower complexity of the application (that does not have to distribute multimedia data into layers for example) and reduces the server's load. On the other hand, the sender has to be informed about the quality of delivery and feedback is employed to carry the information from the clients. In relation to the solution for varying the quantity of data to be delivered according to the existing delivery conditions for the pre-recorded streams, switching among multiple quality pre-encoded versions was chosen due to the late high decrease in costs of the storage capacity that makes it more attractive. In both live and pre-recorded cases, the performed adaptive adjustments aim at modifying the transmission rate in an AIMD manner.

In consequence QOAS is a server-based adaptive control scheme for multimedia deliveries, that monitors the effect delivery conditions have on the end-user perceived quality at the clients and report it to the sender via feedback in order adaptive decisions to be taken. For building the confidence in its results, the scheme has to be assessed from end-user perceptual point of view both using an objective metric and using subjective testing.

2.5 User Perceived Quality (Research, Metrics, Testing)

2.5.1 Necessity of User Perceived Quality Assessment

The assessment of quality is a very important issue with the increasing use of digital multimedia as a significant part of the emerging rich services such as digital TV, video on demand, videoconference, etc. The competing research groups, in order to allow for achieving different constraints (e.g. delay, complexity, etc.), have proposed schemes for compression, processing and transmission of digital multimedia that very much differ in the manner their performances affect the quality of the outputted streams. In general they introduce some impairments, which for example for video are strongly dependent of the levels of details and motion in the scenes. Moreover, the human perception of these impairments also depends on the characteristics of the content making traditional evaluation methods inadequate for their quantification. In consequence there is a need for evaluation methods that quantify the quality of digital streams in order to assess the performances of both the proposed algorithms and of the systems that use them.

Since the proposed solution - QOAS focuses both on modifying of the video component of the multimedia stream and on taking into account of the end-user perceived quality as an active actor in the adaptive control scheme, next the efforts of user perceived quality assessment in relation to digital video are presented in detail.

There are two main directions the research that assesses the quality of digital video streams has taken: **objective methods** and **subjective testing**. Next both of them are presented, after the possible impairments of the remotely delivered video streams over IP networks are brief reminded.

2.5.2 Possible Impairments of Remotely Delivered Video Streams

Regardless of their categorisation the methods that assess the perceived quality of a remotely transmitted stream over IP-networks have to take into account different types of artifacts that may appear. The most likely source of such impairments are coding and transmission. The main types of impairments in a remotely delivered video sequence were presented and defined in [130, 131, 132] and are mentioned next.

Encoding artifacts are mainly caused by the lossy quantisation step applied in most of the existing encoding scheme which are based on Motion Compensation (MC) and block-based Discrete Cosine Transform (DCT). Blocking effect or tiling is defined in [131] as a distortion of the image characterised by the appearance of false blocks within a picture. Tiling is caused by the independent quantisation of blocks and is the most apparent visual impairment. Blurring is a global distortion over the entire image, characterised by reduced sharpness of edges and spatial detail [131]. It is the result of the suppression of higher-frequency coefficients by a coarser quantisation. The mosquito effect is defined as a form of edge busyness characterised by time-varying sharpness (shimmering) at the edges of objects. This temporal artifact is the result of different coding of the same area of the image in subsequent frames [132]. Jagged motion is the result of poor motion estimation, while jerky motion or jerkiness is defined in [131] as a continuous motion perceived as a series of discontinuous images. This is due to lost motion when video is transmitted at lower frame rates. Other possible artifacts are colour bleeding, random noise, and chrominance mismatch [132]. Some of these effects are unique to block-based coding, while others are prevalent in other compression algorithms. For example if using wavelets there are no block-related artifacts, but blurring may become more noticeable.

Transmission artifacts appear because the stream is fragmented into packets and sent over the network, subject to loss and variable delays that cause data unavailability at required time at the remote decoder and then player. The effect of data unavailability depends on the level of redundancy of the encoded stream (for example, intra-coded bitstreams are more resilient to loss). For MC/DCT codecs, like MPEG, interdependencies of syntax information can cause an undesired effect in which the loss of a macro-block may corrupt subsequent macro-blocks until the decoder can re-synchronise. These result in error blocks within the image (spatial propagation) and contrast greatly with adjacent blocks, having a major impact on perceived quality. Another problem arises when blocks in subsequent frames are predicted from a corrupted macro-block - they will be damaged as well and this will cause a temporal propagation of loss until the next intra-coded macro-block is available. More details about the propagation of errors in MPEG-2 streams are given in [133, 134].

2.5.3 Objective Assessment of User Perceived Quality

Objective methods aim at determining the quality of a video sequence in the absence of the human viewer. They are based on very different principles such as, for example: the comparison between the original and the distorted version of the same sequence, the statistical assessment of a large set of analysed cases, the analysis of the effect of possible interferences with the video streams and the relationships from the video contents and subjective testing results. The researchers divide the metrics associated to these objective methods into mathematical-based and model-based [135, 136]. The mathematical metrics rely on mathematical formulae or on functions based on intensive psycho-visual experiments. The model metrics are based on complex models of the human visual system. Apart from this categorisation, according to their possible usage, two approaches were distinguished in [137]: out-of service metrics and in-service metrics. The out-of service metrics base their operation on the fact that the full reference video is available and no time pressure is put to perform the computation. Although the majority of existing metrics belong to this category, their usage in real-time multimedia systems is limited. The in-service metrics are meant to operate while systems are in-service, allowing to perform measurements regularly and eventually to take actions if proved to be necessary. In general the original stream is not available and therefore the associated algorithms estimate the perceived quality based on a-priori knowledge about the encoding scheme, multimedia content, expected artifacts etc. In [132, 138] the authors list three approaches according to the requirement of the existence of the source video into: full reference methods (FR) (also called picture comparison-based), reduced reference solutions (RR) (also called feature extraction-based) and no reference methods (NR) (also called single-ended). Only the last category of methods is useful for in-service application.

Since 1997 the Video Quality Expert Group (VQEG)³² has studied extensively possible assessment of video quality. One of its goals was to propose a quality metric for ITU-T standardisation. It has tested proposals made by 10 different research groups and its aim was to check them in terms of: **prediction accuracy** (the ability to predict the subjective quality), **prediction monotonicity** (the degree the predictions agree with subjective quality ratings) and **prediction consistency** (if the prediction accuracy is maintained over the range of video test sequences, video systems, video impairments etc.). After over 26,000 subjective opinion scores were generated based on 20 different source sequences at bit rates between 768 kb/s and 50 Mb/s, processed by 16 different video systems and evaluated at 8 laboratories, they were compared to the tested objective metrics, among the conclusions drawn were the following:

- No perceptual model is able to fully replace subjective testing
- No perceptual model statistically outperforms the others in all conditions
- No method was recommended to the ITU for standardisation.

The proposals taken into account by VQEG are presented in its final report [139] and they will be briefly presented next along with some other models proposed by different research groups. Currently VQEG continues its work on FR quality assessment for television, aiming also RR and NR quality assessment for television and multimedia³².

2.5.3.1 Mathematical Metrics

The mathematical metrics rely on mathematical formulae or on functions based on intensive psycho-visual experiments. Among them the best known are PSNR and WSNR.

2.5.3.1.1 Peak-Noise-to-Signal-Ratio (PSNR)

A metric tested by VQEG was the peak-noise-to-signal-ratio (PNSR). PSNR is defined as in equation (2-1):

$$PSNR = 10 \log_{10} \left(\frac{255^2}{MSE} \right) \tag{2-1}$$

³² The Video Quality Experts Group (VQEG), VQEG Web Page, http://www.its.bldrdoc.gov/vqeg

where MSE represents the mean square error and is computed as in equation (2-2).

$$MSE = \frac{1}{(P_2 - P_1 + 1)(M_2 - M_1 + 1)(N_2 - N_1 + 1)} \sum_{p=P_1}^{P_2} \sum_{m=M_1}^{M_2} \sum_{n=N_1}^{N_2} (M(p, m, n) - R(p, m, n))^2$$
(2-2)

In equation (2-2), M(p,m,n) and R(p,m,n) are the values associated with the pixel located in frame p, row m and column n, of the modified and respectively the reference video stream.

2.5.3.1.2 Weighted Signal to Noise Ratio (WSNR)

The weighted signal to noise ratio metric (WSNR) takes into account some human visual system properties through weighting as for example the weighted noise power density as a function to the eye sensitivity.

Although they seem appropriate and are very simple, many studies [133, 135, 137] have shown that PSNR and WSNR are poorly correlated to human vision, not taking into account for example visual masking. For example this leads to similar decreases in scores regardless if the human subjects can or cannot perceive the difference from the original. Another problem is that these metrics are applied on frame-by-frame bases, not taking into account temporal correlation between frames. Also, being out-of-service metrics, the original set of frames has to be available.

2.5.3.1.3 Picture Appraisal Rating (PAR)

The Picture Appraisal Rating (PAR) [140] was proposed by Snell & Wilcox³³ as a noreference method of estimating the picture quality of an MPEG-2 video by measuring the distortion introduced by the MPEG encoding process. Based on PAR, Snell & Wilcox have launched Mosalina [141] an off-line, single-ended monitoring process that automatically detect possible picture quality problems in MPEG-2 streams and MVA200 [140], a real-time MPEG bit stream analyser. Although being a no-reference objective method PAR is best suited for in-service applicability, the algorithm was not built to directly detect artifacts that might be introduced by a decoder in response to problems in the transmitted stream for example³⁴. Another problem might be that PAR is based on PSNR that limits its correlation to the human visual system.

³³ Snell & Wilcox, Web Site, http://www.snellwilcox.com

³⁴ François Abbe, PAR Frequent Asked Questions, Snell & Wilcox, Sep. 2000, http://www.snellwilcox.com/reference/par_faq.html

2.5.3.2 Model-based Metrics

The model-based metrics are more complex and rely on human visual models in order to quantify the quality of a video sequence.

2.5.3.2.1 Image Evaluation based on Segmentation (IES)

The Image Evaluation based on Segmentation (IES) model was proposed by CPqD³⁵ to VQEG for assessment. It bases its operation on scenes' segmentation into plane, edge and texture regions, and on the assignment of a number of objective parameters to each of these components. A perceptual-based model that predicts subjective ratings based on the relationship between existing subjective test results and their objective assessment is used to obtain an estimated impairment level for each parameter. The final result is achieved through a combination of estimated impairment levels, based on their statistical reliabilities. An added scene classifier ensures scene independent evaluation. This model is very complex and its reliability limited which makes difficult its applicability, especially in real-time.

2.5.3.2.2 Picture Quality Rating (PQR)

The joint Tektronix³⁶/Sarnoff³⁷ VQEG submission, is the Picture Quality Rating (PQR) metric based on Sarnoff's Human Vision Model (HVM) that simulates the responses of human spatio-temporal visual system taking into account the perceptual magnitudes of differences between source and processed sequences. From these differences, an overall metric of the discriminability of the two sequences is calculated based on their proprietary JNDmetrix (Just Noticeable Difference)³⁸. The model was designed under the constraint of high-speed operation in standard image processing hardware and thus represents a relatively straightforward, easy-to-compute solution. Tektronix has already released two Picture Quality Analysis Systems PQA200 and PQA300 based on this. More details about PQR can be found in [142].

³⁵ CPqD, http://www.cpqdusa.com

³⁶ Tektronix, http://www.tek.com

³⁷ Sarnoff, http://www.sarnoff.com

³⁸ Just Noticeable Difference Metrics, (JNDmetrix), http://www.JNDmetrix.com

2.5.3.2.3 NHK/Mitsubishi Model

NHK Science and Technical Research Laboratories³⁹ and Mitsubishi Electric Corp.⁴⁰ have jointly proposed to VQEG a model that emulates human-visual characteristics using 3D (spatiotemporal) filters, which are applied to differences between source and processed signals. The filters characteristics are varied based on the luminance level. An output quality score is calculated as a sum of weighted measures from these filters. The hardware version is available and can measure picture quality in real-time. However being a FR method, it can be applied only when the presence of the original video source is available.

2.5.3.2.4 KDD Model

Kokusai Denshin Denwa (KDD) Research and Development Laboratories⁴¹, part of KDDI Corporation - Japan, has proposed for VQEG consideration a model based on mean square error (MSE) calculated by subtracting the test signal (Test) from the reference signal (Ref). MSE is then weighted by a set of sequential Human Visual Filters F1, F2, F3 and F4. F1 is a pixel-based spatial filter, F2 - a block-based filter, F3 - a frame-based filter and F4 - a sequence-based filter.

Lately KDD and Pixelmetrix⁴² have jointly launched VP2000 Series Picture Quality Analyzer, based on the KDD's model [143]. Unfortunately, very little information is provided about this full-reference proprietary model, which makes it difficult to be used.

2.5.3.2.5 Perceptual Distortion Metric (PDM)

The perceptual distortion metric (PDM) [134, 144] proposed by L'Ecole Polytechnique Federale de Lausanne (EPFL) - Switzerland, is based on a spatio-temporal model of the human visual system. It consists of four stages, through which both the reference and the processed sequences pass. The first converts the input to an opponent-color space. The second stage implements a spatio-temporal perceptual decomposition into separate visual channels of different temporal frequency, spatial frequency and orientation. The third stage models effects of pattern masking by simulating excitatory and inhibitory mechanisms according to a model of contrast gain control. The fourth and final stage of the metric serves as pooling and detection stage and computes

³⁹ NHK Science and Technical Research Laboratories, Japan, Web Site, http://www.nhk.or.jp/strl/aboutstrl/doc/introset01-e.html

⁴⁰ Mitsubishi Electric, http://www.mitsubishielectric.com

⁴¹ Kokusai Denshin Denwa (KDD) Research and Development Laboratories, KDDI Corporation, http://www.kddilabs.jp/english

⁴² Pixelmetrix, http://www.pixelmetrix.com

a distortion measure from the difference between the sensor outputs of the reference and the processed sequence. VQEG testing has also considered PDM. Although complex this metric has the advantage that many of its details are made public. However the fact that is a full-reference metric makes impossible its usage in real-time.

2.5.3.2.6 Digital Video Quality (DVQ)

The Digital Video Quality (DVQ) model and metric [145] was proposed by NASA - USA and is subject to U.S. patent no. 6,493,023 [146]. Being a full-reference model it requires as input a pair of color video sequences: the reference and the test. In the first step the sequences are sampled, cropped, and subject to color transformations that restrict processing to a region of interest and to represent the sequences in a perceptual color space. De-interlacing and de-gamma-correcting on the input video is also performed. The sequences are then subjected to blocking and DCT and the results are then transformed to local contrast. The next steps are a time filtering, a spatial filtering and a contrast masking operation. Finally the masked differences are used to compute a quality measure. Rhode and Schwarz⁴³ use this technique in the commercially available Digital Video Quality Analyzer DVQ. The greatest problem with this solution is the fact that is patented and using it requires a supplemental license cost.

2.5.3.2.7 Perceptual Video Quality Measure (PVQM)

The Perceptual Video Quality Measure (PVQM) [147] was proposed for VQEG assessment by KPN Netherlands and Swisscom CT Switzerland. It uses the same approach for measuring video quality as used in the Perceptual Speech Quality Measure (PSQM), standardised in the ITU-T R. P.861 [148], for measuring speech quality. The method was designed to cope with spatial, temporal distortions, and spatio-temporally-localised distortions like found in error conditions. It is a full-reference metric and therefore uses two input video sequences (reference and modified) and it bases its operation on the fact that the Human Visual System is much more sensitive to the sharpness of the luminance component than that of the chrominance components.

⁴³ Rohde and Schwarz, http://www.rohde-schwarz.com

2.5.3.2.8 Video Quality Metric (VQM)

The Video Quality Model (VQM) was proposed by the Institute for Telecommunication Sciences, the USA's National Telecommunications and Information Administration (NTIA)⁴⁴ and is subject to U.S. patent no. 6,496,221 [149]. Based on extensive research⁴⁵ and on an earlier model [150], VQM is a full-reference metric that uses reduced bandwidth features that are extracted from spatial-temporal regions of processed input and output video scenes. These features characterise spatial detail, motion, and color existent in the video sequences. Gain and loss parameters are computed by comparing two parallel streams of feature samples, one from the input and the other from the output. Gain and loss parameters are examined separately for each pair of feature streams since they measure fundamentally different aspects of quality perception. A linear combination of the results is used for the subjective quality rating computation. Although some publications have described the principle of VQM, detailed information about this patented metric was not revealed.

2.5.3.2.9 Full-reference Moving Pictures Quality Metric (MPQM)

Proposed by researchers from L'Ecole Polytechnique Federale de Lausanne (EPFL) - Switzerland and presented in [135, 151], the Moving Pictures Quality Metric (MPQM) is a full reference video quality metric based on a basic multi-channel human visual model that takes into consideration modelling of contrast sensitivity and intra-channel masking. It is based on the computation of a distortion metric E for each channel on 3D blocks that are defined by 2D blocks that cover two-degree visual angles and roughly 100 ms in time that accounts for the persistence of images on the retina. The formula is presented in equation (2-3).

$$E = \left(\frac{1}{N} \sum_{c=1}^{N} \left(\frac{1}{N_x N_y N_t} \sum_{t=1}^{N_t} \sum_{y=1}^{N_y} \sum_{x=1}^{N_x} |e[x, y, t, c]|\right)^{\beta}\right)^{\frac{1}{\beta}}$$
(2-3)

where e[x, y, t, c] is the masked error signal at position (x, y) at time t in channel c, N_x , N_y and N_t are the 3D blocks' dimensions and N is the number of channels. The exponent of the Minkowski summation β is 4.

⁴⁴ The Institute for Telecommunication Sciences, National Telecommunications and Information Administration (NTIA), USA, http://www.its.bldrdoc.gov

⁴⁵ Video Quality Research, The Institute for Telecommunication Sciences, http://www.its.bldrdoc.gov/n3/video/Default.htm

The quality rating Q is computed from E as in equation (2-4):

$$Q = \frac{5}{1 + N^* E} \tag{2-4}$$

where N=0.623 and was chosen on the basis of the human vision model.

The Color Moving Pictures Quality Metric [135, 152] applies the MPQM to the luminance and two chrominance components, after they were separated from the input sequences.

2.5.3.2.10 No-reference Moving Pictures Quality Metric (Q)

The biggest problem with the MPQM-based metrics is that they require the presence of the original video sequence. However, the researchers from L'Ecole Polytechnique Federale de Lausanne (EPFL) - Switzerland have also proposed a no-reference metric that estimates MPQM results based on a-priory knowledge about encoding scheme (MPEG) and the effect of the loss on the encoded stream [133]. Named Q the metric describes the joint impact of MPEG rate and data loss on video quality and has the formula presented in the equation (2-5).

$$Q = Q_0 + \chi_q * \left(\frac{\overline{R}}{\chi_r}\right)^{-\frac{1}{\xi_r}} + \chi_l * \overline{R} * PLR$$
 (2-5)

In equation (2-5) PLR is the packet loss ratio, \bar{R} is the stream's mean bitrate, the constant Q_0 has a value close to the maximum quality 5, χ_q , ξ_r and χ_r are related to the complexity of the sequence and χ_l depends also on the average bitrate of the stream. The fact that Q is a noreference metric, is not proprietary and has a simple formula makes it easy to be used for in-service monitoring of video quality.

2.5.4 Subjective Assessment of User Perceived Quality

Formal subjective tests as defined by ITU-R BT.500 [153] have been used for many years and lately the ITU-T R. P.910 [154] has specifically addressed subjective tests for multimedia applications. Among the advantages of subjective testing one could mention that the tests could be designed to accurately represent a specific application, direct users' opinions are gathered and valid results obtained regardless of the system used, the motion content of the sequences, the compression used, etc. Among the drawbacks of the subjective testing are a wide variety of possible methods and

test element parameters must be considered, complex setup and control are required, many observers must be selected and tested, and it very time consuming and costly. Therefore subjective tests are only applicable for development purposes and cannot be used for in-service quality monitoring. Details about testing conditions, testing methods, testing phases, rating scales, testing sequences, etc. are given in [153, 154].

2.5.5 Conclusions

Although many subjective tests and objective methods can quite accurately measure the impairments of digital video sequences, there are still many problems related to both subjective and objective testing and some of them were mentioned by the Internet2's QoS Working Group⁴⁶ in its QoS report [132]. Firstly subjective tests have been mostly defined for short video sequences (approximately 10s duration) which are not long enough to experience all the types of impairments that occur in a real video application and to allow the subject to assess them carefully. Secondly, the existing objective models cannot yet accurately grade all the impairments caused when digital video streams are transmitted over IP networks. On one hand, other mechanisms have to be used in conjunction to account for non-visual distortions (delay and delay variations-based), on the other hand, in-service metrics have to be proposed not to require the presence of a reference stream for quality level quantification. Thirdly, all the existing quality metrics are thought to grade only video sequences and does not account for audio components which are a significant part of the multimedia experience and may influence the assessment of the sequences' overall quality.

Therefore for best and confident results during the development phase, it is recommended to use both subjective and objective methods, in conjunction. For the in-service operation, no-reference methods are the only choice, although they are not fully matured yet.

2.6 Improving Performances of Multimedia Deliveries

Performance in terms of multimedia distribution can be looked at from different points of view, among which the most important are the customers' - on one hand and the service providers' and the network operators' - on the other. Taking into account all the interests involved, the performance of multimedia deliveries could be measured for example in terms of range of services offered, their corresponding end-user perceived quality and used bandwidth. Different solutions

⁴⁶ Internet2 QoS Working Group, http://gos.internet2.edu/wg/

were proposed in order to try to maximise one or some of these metrics, among which error control, protocols and delivery solutions are mentioned. Next these directions are explored, presenting different solutions, their advantages and disadvantages.

2.6.1 Error Control

Error controls aims at reducing the effect of loss on the quality of the remotely transmitted and played multimedia streams. Loss occurs due to the either network congestion that leads to the network routers' buffers overflowing and consequent dropping of the incoming packets or due to long or extremely variable delays that prevent the packets from arriving in time to be decoded and played at the receiver. Also some packets may arrive corrupted at the destination. Among the most significant for the effect of the correct data unavailability on the end-user perceived quality are the loss characteristics (pattern, duration, etc.), the compression algorithms used and the delivery solutions employed for data transmission. Among the worst affected by losses are the streams encoded using some compression schemes that achieve high compression efficiency like MPEG-1 [73], MPEG-2/H.262 [74, 75] and H.261 [80]. For these streams even small losses severely degrade their related video quality either due to the decoders' loose of synchronisation that make them to skip correctly received data until the next synchronisation point or because of error propagation that causes an error that affects a frame to affect also other frames that depend on the first frame's data in their decoding process.

Different error control mechanisms have partly addressed these problems and have proposed different solutions. In [60] the authors have distinguished four approaches: forward error control (FEC)-based mechanisms, retransmission, error-resilient encoding and error concealment. These directions are presented next in relation to video deliveries. A detailed survey on audio-related error control mechanisms was published in [155].

2.6.1.1 FEC-based Mechanisms

The principle behind forward error correction (FEC) is to add redundant information to the original data to be transmitted that would allow its reconstruction even in the presence of loss. In [91] FEC-based mechanisms are classified in three categories: channel coding, source coding and joint channel/source coding.

Channel coding involves the division of the continuous stream in segments and each segment is divided in k packets. For each segment a specific block code is applied on the K packets

resulting a group of N coded packets (N > K) that is send to the receiver. If any number of packets greater than K is received, the receiver can reconstruct completely the original transmitted data. The problem with channel coding solutions for error control is that the protection against errors they offer is provided at cost of increased used bandwidth for the whole duration of the transmission, regardless of the probability of the appearance of loss. Therefore these solutions are more likely to be used if the loss probability is high, severely affecting the video quality or constant and if bandwidth can be spared for transmitting the required extra information. Moreover, since channel coding-based schemes are applied in generally to a set of packets these solutions introduce delays and burstiness in traffic that may affect the performances of the multimedia delivery. Some channel coding-related error control schemes are involved in equal error protection (EEP), in which all bits are equally treated, others in unequal error protection (UEP) when extra protection is applied to more important data bits or in hierarchical FEC. Some mechanisms are presented in [156, 157].

Source coding is similar to channel coding in sense that it adds redundant information to ensure that the data can be recovered after loss. The difference constitutes the fact that for example the Nth packet contains the Nth group of blocks and a compressed version of the (N-1) th group of blocks that would allow the reconstruction of the Nth in case of loss, but at a lower quality. Source coding would suffer from the same problems as channel coding in terms of bandwidth requirement and invariability to loss variations, but would involve lower delays. An example is presented in [158].

Joint channel/source coding combines the channel coding and source coding approaches. More details are given in [91].

2.6.1.2 Retransmissions

Some researchers [159, 160] have proposed retransmissions of lost packets in order to provide error control. Unfortunately since the retransmitted packet arrives at least three times the one-way trip time after the transmission of the original packet, it is very likely for a retransmitted packet, part of a time-sensitive stream, to arrive at the destination after it is required and it will be discarded, making the effort futile. However there are solutions like delaying frame play-out times to allow for the arrival of retransmitted multimedia packets, selectively retransmit only those packets that, due to buffering, would have enough time to reach the destination before they are needed and selectively retransmit only important packets and only if it is estimated they will arrive in time. However, these solutions add significant latency and cannot be used for interactive or time-

sensitive media deliveries, unless the one-way delay is very short in comparison with the acceptable delays, which limits their applicability.

2.6.1.3 Error-resilient Encoding

The idea the error-resilient encoding is based on is to increase the robustness of the compressed stream to packet loss and is in general performed at the source of data. Classic error resilient encoding includes re-syncronisation marking, data partitioning and data recovery and they were standardised as part of MPEG-4 encoding scheme [75, 76]. Unfortunately these are targeting more error-prone environments like wireless channel and do not address low probability loss infrastructures such as wireline broadband IP-networks for example. Related to the latter, the authors of [91] distinguish two directions error resilient coding have taken: optimal mode selection and multiple description coding.

Optimal mode selection refers to the approach that tries to offer increased performance for video deliveries subject to loss based on characteristics of source, path and receiver behaviors [161]. Among them is for example the manner of choosing between intra-mode and inter-mode of coding blocks of video data in order to achieve both good compression and to limit error propagation [162]. A higher number of intra-mode-encoded blocks means higher robustness to loss, but lower compression ratio, whereas more inter-mode-encoded blocks mean higher compression ratio, but increased chance for error propagation.

Multiple description coding refers to the compression of a single video sequence into multiple streams (named also descriptions) in such a manner that each provides acceptable quality, but combined offer a better visual quality [163]. Among the advantages is the increased robustness to loss since the receivers do not have to get all the descriptions to view the content and additive quality as if the receivers get more descriptions their perceived quality becomes better. The disadvantage is that in order to allow for independent decompression, the descriptions carry redundant information relative to each other, decreasing the compression efficiency and therefore requiring higher overall bandwidth.

2.6.1.4 Error Concealment

Unlike error-resilient encoding that involves mainly the source of data and is performed prior to loss, the error concealment methods involve the receivers and are applied after the loss has occurred in order to minimise its effects on the quality of the displayed video stream. There are two

main approaches for error concealment: spatial interpolation and temporal interpolation [59]. Spatial interpolation refers mainly to the reconstruction of missing parts of frames from the neighboring blocks, whereas temporal interpolation involves the reconstruction of missing data with information from previous frames. Three simple error concealment methods were distinguished for block-based compressed streams subject to loss in [164]: EC-1 the current frame affected by packet loss is replaced by the previous frame, EC-2 the block corrupted by loss is replaced with the block from the same position from the previous frame and EC-3 the corrupted block is replaced by a block from a previous frame pointed by a motion vector. EC-1 and EC-2 have a lower complexity than EC-3, but the latter would achieve better quality. The same relationship is between EC-1 and EC-2 with the latter achieving better reconstructed image quality. Among the advantages of the error concealment methods are their reduced complexity and limited effort of their application in comparison to other error control methods taken into consideration. However detecting and repairing losses incur latency and in general the quality of the reconstructed image is not very high.

2.6.1.5 Comments

Error control provides means for either offering greater protection against errors or, if they have already occurred, for minimising their impact on the user-perceived performance. However error control comes with a cost in terms of necessary bandwidth, increased delays and computational requirements. These gains and costs have to be carefully balanced for each solution in order to determine the correct approach to be taken. For example in fully loaded broadband local IP-networks there is no spare bandwidth for using FEC-like methods or retransmissions. However, the latter are recommended anyway only in certain cases when delivering time-sensitive data such as multimedia. Different types of error-resilient encoding can be used as well as error concealment methods, which are the easiest to be deployed among these solutions.

2.6.2 Protocols

Few protocols have been proposed and standardised in order to support the deliveries of continuous multimedia streams. According to their functionality these protocols can be classified in three categories: network-layer, transport and session control [59]. Next each of these categories is presented and some standard protocols that belong to it are mentioned.

2.6.2.1 Network-level Protocols

Network-level protocols are supposed to provide basic network support such as network addressing. For video streaming over IP-networks, the Internet Protocol (IP) [14] provides these services.

2.6.2.2 Transport Protocols

Transport protocols were proposed in order to provide end-to-end network transport functions and among the best known are User Datagram Protocol (UDP) [165] and Transmission Control Protocol (TCP) [166] - lower-layer transport protocols and Realtime Transport Protocol (RTP) [100] and Real Time Control Protocol (RTCP) [100] - upper-layer transport protocols. The first ones support functions like multiplexing, error control, congestion control or flow control. The most significant for RTP is that it provides time-stamping, sequence numbering, payload type identification and source identification, whereas for RTCP is that it allows for providing QoS feedback to the participants of a RTP session, being a companion protocol to RTP.

2.6.2.3 Session Control Protocols

Session control protocols in relation to multimedia streams' deliveries aim at controlling multimedia sessions. Among the best known are the Real Time Streaming Protocol (RTSP) [167] and the Session Initiation Protocol (SIP) [168]. RTSP allows session establishment and control, as well as multimedia presentation. It offers VCR functions such as play, pause, rewind, stop, etc. SIP also initiates and controls multimedia sessions, providing also support for user mobility by proxying and redirection.

2.6.2.4 Comments

Since TCP uses retransmission to ensure reliable transport of data, it is not suitable for transmitting data with timing constraints like in the case of multimedia streaming. In these cases UDP is mainly used. Since UDP does not guarantee the arrival of packets at destinations, RTP is regularly employed to detect packet loss. In general RTCP is used in order to provide feedback about the QoS of the provided services. A session control protocol and more often RTSP, is used to initiate and control the multimedia session, including the data delivery. This is also the approach chosen in this thesis.

2.6.3 Solutions for Delivery Architectures

Among the best-known components or solutions for delivery architectures are proxy servers, caches, mirrors, content distribution networks and peer-to-peer solutions. Next they will be briefly presented.

2.6.3.1 Proxy Servers

As their name would suggest proxy servers are agents that intermediate between the real service provider – the server - and the receiver. In general they are used in more complex forms acting as gateways, firewalls, caches, etc., alone or part of a co-operating structure. Related to video deliveries they could help reducing network bandwidth requirement, delays and delay variations over WAN, decreasing the loads on the video servers, decreasing the start-up delays by storing the initial sequences of some video, smoothing the video streaming, improving VCR functionality, transcoding video to adapt to heterogeneous bandwidth or customers' requirements, etc.

2.6.3.2 Caching

Caching is based on the observation that some content is more popular than other is. If a copy of the already served data is placed closer to the customers, in the case that it will be requested significant benefits from the performance point of view will be achieved. The more the same cached content is requested, the more the benefit increases. This performance benefit is generally considered in terms of bandwidth, server load and service latency.

Although caching was well studied in relation to small, static, WWW-related objects [169], there are different issues related to caching that have to be taken into account in relation to storing and retrieving of multimedia content. The bandwidth requirements and the size associated with multimedia files put pressure on the caches that have to support high bandwidth, to have increased storage capacity and different policy since not all the content can be cached at a time. The popularity of continuous streams is not well defined and it is a generally agreed opinion about what parts of the video must be cached. Adaptive control schemes and different delivery techniques may have also a significant influence on the caches' design and functionality. Moreover the caches' location, their number and the relation between them greatly determine the performance of the cache-based solutions. Since there are many issues related to caching of multimedia sequences, next some of them are presented: segmentation of streaming objects, replacement policies, cache consistency, pre-fetching, caching architectures and co-operative caching.

Since the caches cannot store all the data associated to a multimedia file, they have to perform **segmentation of streaming objects** and to apply their policies at the level of a segment. This segment can be arbitrarily chosen from within the stream as in [170], can be the prefix of a popular stream as chosen in [171], can be the result of a content-aware segmentation process as in [172] or can belong to one of the layers of a layered encoded stream as in [173]. These segments can have a fixed size as in [174] or a variable size as in [172].

Replacement policy refers to the algorithm according to which a cache would choose an already stored content for deletion in order to create space for storing a more recent object. Traditional replacement policies are [175]: the least recently used (LRU) which evicts the least recently requested object, least frequently used (LFU) which evicts the least frequently used object and Pitkow/Recker which evicts objects in LRU order but between the objects accessed in the same day chooses the largest. However with video deliveries and their associated characteristics other policies have appeared like the eviction of segments in the descending order of their quality [173], according to the principle of temporal locality [172, 174] or according to the clients' bandwidth [176].

Cache consistency or coherency methods aim at making sure that the cached objects reflect existing objects from the server where they originated. Cache consistency was subject to extensive research in relation to regular Web objects and different techniques were proposed such as: client polling, invalidation callbacks, time-to-live and if-modified-since [177]. Since continuous multimedia streams are segmented prior to caching, the same techniques can be applied for ensuring multimedia streams' objects cache consistency [174].

Pre-fetching refers to retrieving data from original servers in anticipation of clients' requests [177]. Since bandwidth requirements and objects' sizes are larger, pre-fetching has to be performed cautiously [172], but is necessary because it reduces the user-perceived latencies. In [172] pre-fetching is performed based on a prediction algorithm that analyses the users' interaction with the video stream, the available bandwidth and storage space at the cache. In [173], apart from bandwidth and storage constraints, the pre-fetching algorithm takes into account the quality of the cached segments whose priority decreases with their quality level. However the gains of pre-fetching come with a cost in terms of increase in traffic and in its burstiness [176].

Caching architectures involve more caches that serve a higher community of users increasing therefore the probability for requested object to be found, increasing the performance of the caching process [169]. A hierarchical caching architecture was first proposed in the Harvest

project [178] and since then other works have used it [172]. A totally distributed caching architecture as in [179] has caches placed only at bottom level. Hybrid caching solutions allow for the existence of a complex architecture, with different levels. More information about caching architectures can be found in [175]. However there are several problems [180] related to placing the caches in the network, their co-operation, consistency, additional delays, bottlenecks, etc.

Co-operative caching refers to the collaboration between several caches in order to increase the performance of the system. The idea is that if the requested information is not found in a cache, it will be looked for in other caches. The Internet Cache Protocol (ICP) [181] was proposed to support this and involves fetching a document from a neighbouring cache with the lowest RTT. Different co-operation approaches were tried in [170, 172, 182] for caching video such as using a master cache in a hierarchical solution, distributed caches or hybrid architecture. However sometimes better performances are obtained if this co-operation is limited as proposed in [183] that suggests retrieving the content from the original server than using distant or slow caches.

2.6.3.3 Mirorring

Mirroring refers to placing copies of the original multimedia files on servers situated at different locations. In this way the clients can choose the location of the server they can retrieve the multimedia file from in order to have the best performances. This performance-related advantage comes with increase cost of the solution and complex administration of multiple copies of the same content.

2.6.3.4 Content Delivery Networks

Content delivery networks (CDN) are dedicated collections of servers located strategically across a wide area network (e.g. Internet) with the goal of offering services with very low latencies and high quality from closer locations to the users. In general the distribution of content is contracted by content providers and is performed by commercial distributors via their CDNs. Recent studies [184, 185, 186] have analysed the benefits and the costs of using CDNs. However, by bringing the data closer to the users significant advantage can be achieved, especially for distributing continuous, time sensitive content such as multimedia.

2.6.3.5 Peer-to-peer Systems

Peer-to-peer systems base their latest popularity on the fact that they allow exchanging information through a structure based on peers that act both as servers and as clients. They can be based on a centralised indexing architecture like Napster⁴⁷, a fully distributed solution such used by Gnutella⁴⁸ or on hybrid architecture like Kazaa⁴⁹. Former and current peer-to-peer systems use to non-interactively download of data, including multimedia, but very recent the interest has increased for using peer-to-peer solutions also for streaming multimedia like in an adaptive layered technique presented in [187].

2.6.3.6 Comments

Although not exhaustive, this presentation gave some idea about the delivery architectures that can be employed for distributing multimedia. Although none of them comes without some disadvantages, their advantages are significant in terms of bandwidth, delays and provided services. For delivering multimedia to home residences via broadband IP-networks proxy servers and caches could be used for large size networks as well as peer-to-peer systems.

2.6.4 Delivery Techniques

There are different delivery techniques that would allow multimedia delivery to home residences. Among them there are broadcasting, multicasting and unicast solutions that are presented next.

2.6.4.1 Broadcasting

Broadcasting refers to the delivery of a service to all the customers, regardless if they want it at the moment of delivery or not. Although broadly used today due to its reduce cost of implementation in relation to the number of customers served, the greatest disadvantages of broadcasting are that the resources, especially expensive bandwidth, are used anyway, it provides limited services and there is no interaction with the users. There are different proposals that increase

⁴⁷ Napster, http://www.napster.com

⁴⁸ Gnutella Protocol Specification v 0.4, March 2001, Clip2, http://www9.limewire.com/developer/gnutella_protocol_0.4.pdf

⁴⁹ Kazaa, http://www.kazaa.com

its friendliness to the users such as periodic broadcast proposed in [166], but its significant limitations have fuelled searches for other solutions.

2.6.4.2 Multicasting

Multicasting in IP-networks allows for the delivery of a service, including multimedia streaming, only to a group of customers that have chosen to join this group. In this manner multicasting solves some of the problems raised by broadcasting in terms of more efficient use of bandwidth, choice offered to the customers, more flexibility etc. Different algorithms have been proposed in order to achieve certain performances while delivering services in a tree like manner and addressing some limitations and introducing others. Among general limitations it is worth mention that some routers do not support multicasting and changing them involves costs, there is a significant overhead multicasting introduces through the group setup and maintenance, there is a certain latency of leaving a group affecting the use of bandwidth and restricts the possibility of choice for the users. Significant research has addressed multicast delivery of multimedia services to customers, including some adaptive solutions as presented earlier in this chapter (section 2.3).

2.6.4.3 Unicast

Unicast solutions involve a relationship one-to-one between the senders and the receivers of services. This allows for an increased flexibility for the personalisation of the services provided to the customers' and for the efficient bandwidth usage. However it also has drawbacks such as poor scaling to a high number of customers and the necessity to know the other end prior to the connection. Different proposals were made in order to address these limitations and a direction related to multimedia deliveries is through adaptive applications and associated schemes. Some examples of such adaptive schemes are presented in section 2.3 of this chapter.

2.6.4.4 Comments

Different techniques can be used to deliver information to the customers and in particular multimedia related data. Three different approaches were presented that have benefits and disadvantages that have to be taken into account when selecting one or another for applicability. For the success of delivering rich content, high quality multimedia to the residential customers, very significant is what services are offered, their facilities and their quality. Therefore a high degree of

flexibility is important for both the users and the providers, only unicast provides it and this is the reason for choosing an unicast technique for the delivery of multimedia to residential users.

2.6.5 Conclusions

In this section error control solutions, protocols that support multimedia deliveries, solutions for delivery architectures and techniques for deliveries were briefly presented as existing solutions for increasing the performances of multimedia deliveries. Also the advantages and the disadvantages associated with these solutions were highlighted and comments in relation to their applicability in multimedia deliveries over broadband IP-networks were presented. Moreover the solutions selected for usage as part of the designed QOAS-based multimedia streaming system were also mentioned.

2.7 Summary

This chapter presents significant works related to the proposed QOAS for adaptive multimedia streaming in local broadband IP-networks. The chapter starts with an overview of these directions among which very important for the design of QOAS are compression techniques, adaptive solutions for delivering multimedia, solutions for assessing end-user perceived quality and methods for improving performances of multimedia deliveries. Each of these directions is then detailed in a separate section that presents solutions, significant research, metrics, protocols, standards etc. related to the chosen subject. These sections also include conclusions and comments related to the usage of some of the presented solutions by the designed QOAS.

Chapter III Quality Oriented Adaptation Scheme in Local Broadband Multi Service IP Networks

Abstract

The third chapter specifies the context of the QOAS, its applicability and its localisation relative to the delivery of multimedia-based services to the residential customers via broadband multi-service IP-networks. It also aims at finding the best approach for designing QOAS, presents its problems and reduces them to sub-problems, simpler to be solved in a top-down manner.

3.1 Overview

The goal of this research, as was already mentioned in the first chapter, is to propose the Quality Oriented Adaptation Scheme (QOAS), an inexpensive application-level end-to-end adaptive mechanism for streaming multimedia, that helps offering high quality multimedia-based services to home residences via local broadband IP-networks.

Before designing the QOAS and other mechanisms it may rely on, significant issues related to broadband IP networks are considered. Among these are possible broadband IP-networks architectures and the manner they may affect QOAS's design. Next different proposed architectures for delivering services to home residences and business premises are presented and their associated advantages and disadvantages are mentioned and commented in relation to their influence on the QOAS's design.

3.2 Broadband IP-Network Architectures to Home Residences and Business Premises

Significant effort was involved in proposing different architectures for delivering information to home residences via cable TV or telephone infrastructure [188, 189, 190]. This includes multimedia data and on-demand multimedia-based services. Lately these solutions were reviewed, addressing broadband connectivity and targeting especially broadband IP-networks [2, 11, 191, 192, 193, 194]. However the principles behind these architectures are similar and very few details differ. They can also be applied successfully for distributing services to business premises.

In general, three main approaches for architectures aimed at distributing on-demand services, to home residences were taken into account in [185]: a distributed architecture, a centralised solution and a hybrid one. Next they are discussed in relation to the delivery of multimedia-based services.

3.2.1 Centralised Architecture

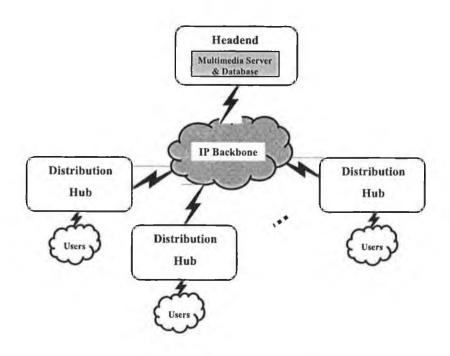


Figure 3-1 Centralised architecture distributing multimedia to residential users

Figure 3-1 presents a typical example of a centralised architecture. This architecture includes a centralised **headend** and a number of **distribution hubs** through which the headend is

connected to diverse groups of users. The headend, based on a multimedia server (or a pool of servers) with access to a multimedia database, provides multimedia-related services to the residential customers via these distributions hubs and also other offered services, such as Internet connectivity. The distribution hubs have minimal responsibilities, which mainly concern data forwarding in both directions: from the headend towards the users and from the users to the headend.

The main advantage of this approach is that it requires only one multimedia server (or server farm) and only one multimedia database, with apparently low hardware costs (although they require high complexity) and, mainly, reduced location and maintenance costs. Also the security is easier to be provided for this approach since a single location has to be protected. The most important disadvantage of this solution is that very much pressure is placed on the IP backbone between the headend and the distribution hubs, pressure that increases significantly with the number of customers served.

3.2.2 Distributed Architecture

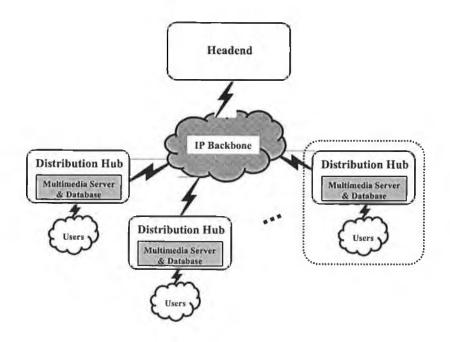


Figure 3-2 Distributed architecture for delivering multimedia to home residences

A typical case of a distributed architecture is presented in Figure 3-2. Similar to the centralised approach, the distributed architecture includes a centralised headend and a number of

distribution hubs through which the services are provided to diverse groups of users. However the headend is in charge only with offering other services such as Internet connectivity and is not involved in providing multimedia-based services to the customers. In this case the distribution hubs have a more important role since at their level there is a multimedia server (or a server farm, but less likely) and a multimedia database that contains multimedia content to be provided to the residential users.

The greatest advantage of this distributed solution is that it releases the pressure placed on the IP backbone in the centralised approach. In this case multimedia data, which have timing constraints and significant sizes and is expected to account for an important part of the total traffic, is served locally. The fact that multimedia-based services are being offered to a smaller group of customers helps at reducing the complexity of the multimedia server system, which could consist of a single server. However, since these simpler multimedia servers and their associated multimedia databases are placed at every distribution hubs, other issues appear which are not favourable to this solution. The disadvantages are mainly in relation to the costs involved in the maintenance of these distributed hubs (i.e. location, power, security) and to the multimedia databases' updates.

3.2.3 Hybrid Architecture

An example of a hybrid architecture for distributing multimedia-based services to residential homes is presented in Figure 3-3. It combines some of the issues provided in the centralised solution with ones that are associated with the distributed approach. The hybrid architecture includes a **headend** similar with the one that exists in a centralised solution and a number of **distribution hubs** with structure and functionality similar to those from the distributed case. The idea this solution relies on is that the multimedia server systems and their multimedia databases from the level of the distribution hubs serve the associated group of users for the majority of their requests. If for some reason a request cannot be fulfilled, the request is re-directed towards the headend whose server system and its database will answer to it.

Other versions of this hybrid approach involve caches located at the level of distribution hubs instead of local multimedia servers and they may be very useful. However although caches are very well studied and recommended for being used with Web content, for continuous media with different characteristics and different interaction with the users for example, the advantages are not yet fully balanced against disadvantages [192].

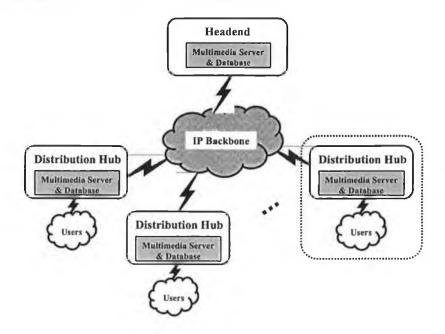


Figure 3-3 Hybrid approach for distributing multimedia-based services to residential users

By using the hybrid architecture, which is a combination of a centralised approach for lowdemand content and a distributed solution for high-demand content, a compromise is also made in terms of advantages and disadvantages.

3.2.4 Comments

A significant decision for the network operators and the service providers is whether to use a centralised, a distributed or a hybrid solution. A centralised architecture avoids the costs of installing and maintaining multimedia servers in remote distribution hubs, but offers limited scalability because of the additional load that is placed on the IP backbone network. If a centralised solution may be acceptable for voice and data-based services, for video-based services at least 10 times more bandwidth over the IP backbone is necessary for each subscriber, making congestion more likely to occur. The consequent delay, delay jitters and packet loss may severely affect the quality of the remotely delivered video services.

The requirements of bandwidth at the IP backbone level are minimised by locating the video servers as near as possible to the subscribers, like in the distributed approach. The claims that a distributed solution may cost more than a centralised one are contradicted by a comparison performed taking into account the distributed connectivity costs, core network bandwidth, headend and distributed server costs, storage and set top box costs. The results reported in [191] found out

that a distributed approach may cost less than a centralised solution, while it could have supplementary benefits in terms of performance for the end-users such as reducing delays and delay variations in video deliveries, for example.

Although the hybrid solution seems to make a compromise between the advantages and disadvantages of the previous two approaches, it also relies very much on the multimedia servers from the distribution hubs to serve a large majority of the requests in order to avoid the congestion of the IP backbone, which will affect both multimedia and other provided services.

As a direct consequence of these findings the decision was to focus the research on the local delivery of multimedia-based services in broadband IP-networks. This was since they carry the very large majority of the overall multimedia traffic in parallel with other types of traffic generated by the other provided services.

3.3 QOAS in Local Broadband Multi-service IP-Network

There are different ways in terms of architectural design for delivering the provided services from distribution hubs to residential homes [189, 191, 193]. Among them, Figure 3-4 shows as an example a pure horizontal distribution structure whereas Figure 3-5 presents a tree-like structure.

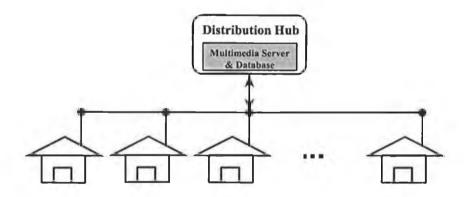


Figure 3-4 Horizontal solution for local distribution of services to home residences

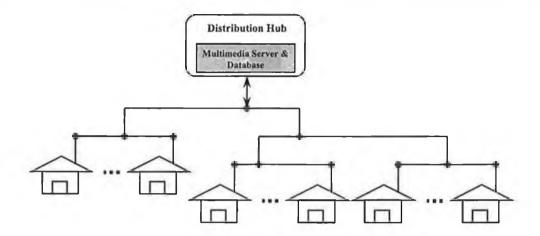


Figure 3-5 Local service distribution to home residences in a tree-like manner

Regardless of the chosen solution for the distribution of services to residential users, the infrastructure that connects the distribution hub with the users has to support all the traffic exchanged by them. A schematic representation of the architecture of the problem is presented in Figure 3-6.

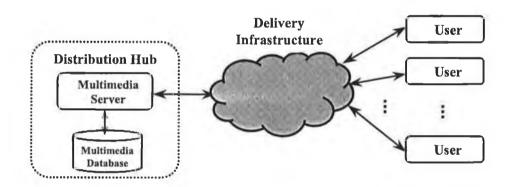


Figure 3-6 Architecture for local multimedia delivery to residential customers

The multimedia server located at the level of the distribution hub was named in [192] local video server, name that will be also used in this thesis.

It is in the interest of network operators, service providers and the customers to reduce the effort for providing these services. A significant source for decreasing it is to raise the number of customers served from a fixed infrastructure, while maintaining a good perceived quality. Therefore it is assumed that the design of the delivery system is such that - at least for periods of times - the infrastructure is overloaded, having potential for congestion.

In this context the problem of delivering high quality multimedia with little effort to home residences can be reduced at two simpler problems. Firstly, the solution has to allow for deliveries of multimedia streams at good quality in increased traffic conditions. The traffic can be of different types and could have various variation patterns since it originates from different services not only multimedia-based ones. Therefore the designed adaptive solution - QOAS, has to be tested taking into account the one-to-one approach and the other traffic as background traffic over the same delivery infrastructure. Secondly, multiple QOAS solution instances have to be deployed for delivering multimedia-based services to a significant number of users and its benefits have to be globally assessed, looking at the overall delivery process.

In consequence QOAS, a unicast adaptive technique for delivering high quality multimedia to the users, has to be designed starting from a one-to-one approach. Its deployment involves placing components at both the level of an adaptive server application (Ad. Srv. App.) and an adaptive client application (Ad. Cli. App.), as shown in Figure 3-7. The infrastructure that links the server-side QOAS component and the client-side one is a local IP network with short propagation delays and high potential for congestion.

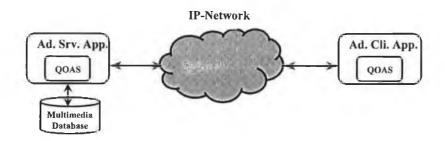


Figure 3-7 QOAS deployment at the level of an adaptive client-server system

3.4 Designing QOAS

QOAS's goal is to maximise the quality, as perceived at the receiver, of a multimedia stream transmitted over the local IP network. This quality is determined by the transmitted quality of the stream and is directly affected by eventual problems that may occur during multimedia transmission over loaded IP networks. Among the causes of these streaming-related problems are network transmission parameters such as packet loss, increased delays, extremely variable delays, etc. as suggested in [133, 195, 196]. Moreover, the perceived quality at the receiver (Q_R) , the server-transmitted quality of the multimedia stream (Q_T) and the network transmission-related parameters, as measured at the receiver, for example in number of N, $(P_i, 1 \le i \le N)$ have values

that vary in time. On top of this, their corresponding values are not recorded at the same time and transmission latency at that moment (D') has to be taken into account in order to achieve accuracy.

Equation (3-1) tries to formalise the relationship between these parameters, where f is the function that has to be maximised by QOAS.

$$Q_R(t) = f(Q_T(t - D^1), P_1(t), P_2(t), ..., P_N(t))$$
(3-1)

This function shows that in order to expect certain quality for the received multimedia stream, modification of the transmitted quality have to be performed in advance with a period of time equal with the transmission latency. Therefore equation (3-2) presents the same function, but written in a form that would allow adjustments to be made in time for their effect on the end-user perceived quality to be effective.

$$Q_R(t+D^*) = f(Q_T(t), P_1^*(t+D^*), P_2^*(t+D^*), \dots, P_N^*(t+D^*))$$
(3-2)

In equation (3-2) D^* is the estimation of the average transmission delay and P_i^* is estimation of the value of the network-related parameter i for the moment t+D* made at moment t. Therefore in order to have a closer representation of this dependency to the one from reality, good estimates of future values of these network-related parameters are necessary to be made.

Unfortunately the situation is very complex, since studies confirmed by prestigious bodies such as IETF through its IP Performance Metrics Working Group [61] have shown that not only immediately previous values or variations of network-related parameters determine their effect on the user perceived quality, but also their variation patterns. This introduces another dimension of complexity in the equation (3-2). Therefore, for each parameter P_i, its values during a certain period of time, prior to the current moment of the multimedia transmission, are taken into account in order to estimate its future value. The duration of this period is subject to debate, but existing research such as [90, 103, 118] suggests that the closer the period to the moment of estimation, the more accurate the estimation is. However, the variation pattern can be considered only during long-term monitoring of parameters and some of these works also suggest taking this into account, as well as short-term values and variations.

Since during streaming values and variations of all these network-related parameters can be measured only with a certain sampling rate and similarly for the transmitted quality of the

multimedia stream, the continuous function from equation (3-2) is replaced by a discrete one: f^d in equation (3-3).

$$Q^{d}_{R}(t+D^{*}) = f^{d}(Q_{T}^{d}(t), P_{1}^{d*}(t+D^{*}), P_{2}^{d*}(t+D^{*}), ..., P_{N}^{d*}(t+D^{*}))$$
 (3-3)

where for each of the N network-related parameters P_i , its value P_i^{d*} at moment t+D* is using an estimator Estim_i^d , according to the generic formula presented in equation (3-4). This estimator bases its computation on M_i previously recorded values for the parameter P_i recorded at moments t_i prior to the moment t, with $1 <= j <= M_i$.

$$P_i^{d^*}(t) = Estim_i^d(P_i^d(t_1), P_i^d(t_2)..., P_i^d(t_{M_i}))$$
(3-4)

Even if the number of very significant network-related parameters in relation to their effect on the multimedia end-user perceived quality is reduced to a minimum, proposing Estimi^d estimators for each of them is very complex. This is because each has different characteristics that have to be taken into account. Once these estimator functions have been defined, they are then used as parameters in the function f^d presented in equation (3-3). Since QOAS goal is to find the necessary measures for maximising the quality at the receiver based on existing information, the problem is reduced to trying to maximise function f^d. However, supposing that even a small number of estimator functions have been added to f^d, the function has a significant number of parameters and its maximisation has a very difficult analytical solution.

In consequence I have decided to follow the path taken by all the existing research in this domain. They have combined heuristic and experimental methods to design the proposed adaptive schemes and the following chapters give more details about the proposed solution.

3.5 Conclusion

After exploring the broadband IP-networks' architectures for distributing multimedia-based services to home residence, the decision was taken to deploy QOAS at the level of local video servers since these servers serve the very large majority of multimedia requests. The traffic generated by other provided services is considered background traffic in relation to QOAS. Similarly, multimedia traffic served by the centralised video server located at the level of the headend in hybrid architectures is considered background traffic. Due to the intended reduced cost

of the solution, no cache or other solutions for improving the quality of the delivery as presented in the previous chapter, is considered to be deployed. However, after the deployment of QOAS or in conjunction with it other mechanisms can be used (e.g. error concealment), but they are not addressed by this work.

In this context the QOAS's problem of delivering high quality multimedia-based services to home residences with little effort via the local broadband multi-service IP-network was formalised and an analytical solution was sought. Since the solution is very complex, the decision taken was to follow a hybrid heuristic-experimental approach. In order to allow for finding an easier solution to the QOAS's problem, two main sub-problems were distinguished. The first aims at finding a one-to-one adaptive solution for delivering multimedia in highly loaded networks, subject to other traffic of different type, size and variation pattern while maintaining a good perceived quality level. The second aims at using the solution already found for simultaneously delivering the highest possible number of streams, serving an increased number of customers from the same infrastructure, significantly reducing the associated costs in comparison to other solutions.

3.6 Summary

This chapter presents the context the QOAS is designed for and tries to determine the best manner for solving the problems it rises. The chapter starts with a presentation of possible architectures for broadband IP-networks and mentions their advantages and disadvantages. Next a possibility for improving the quality of the delivery of multimedia-based services by using QOAS is assessed in order to influence the most of the traffic involved and to bring the best possible benefits. Consequently, QOAS is then localised, having components placed at both the local video server and customers' equipment. The chapter presents at the end how the decision to find a solution for QOAS design using a hybrid heuristic-experimental approach was taken.

Chapter IV Quality Oriented Adaptation Scheme for Multimedia Streaming

Abstract

Network operators and service providers aim for high infrastructure utilisation and a large number of customers to increase their revenues. At the same time, the customers are interested in receiving high quality streamed multimedia, having access to diverse services, and paying a low price. This chapter presents the Quality Oriented Adaptation Scheme (QOAS) - an adaptive solution for high bitrate multimedia streaming that balances these opposing goals. First the QOAS's principle is described and the QOAS's main components are presented. The client-located Quality of Delivery Grading Scheme (QoDGS) is in charge with monitoring, grading and reporting of delivery quality; the server-situated Arbiter (Arb.) that implements the Server Arbitration Scheme (SAS) is responsible with the analysis of end-users' reported quality and with taking adjustment decisions; the Data Transmission and Feedback Mechanisms ensure the delivery of both multimedia data and control information - including feedback. This chapter also presents quality assessment criteria for the QOAS and QOAS applicability considerations.

4.1 QOAS Overview

During high-quality multimedia stream delivery, end-user perceived quality could be negatively affected by *server-related* problems (e.g. server load, software, etc.), by *network-related* problems (e.g. congestion, extremely variable traffic, equipment failures, last-mile low bandwidth connection, etc.), and/or by *client-related* problems (e.g. slow or incompatible software, old hardware, etc.). These problems directly or indirectly cause some *periods of unpredictable delay and/or loss* (PUDL) that affect the overall quality of delivery.

Receiver buffering may be a good solution for many cases of PUDLs, but, used alone, it does not always solve streaming-related problems that occur in difficult delivery conditions. Simply producing multimedia streams at a single lowest or respectively highest quality (and hence lowest or respectively highest bitrate) could also be a solution in case of homogeneous customers, but may leave heterogeneous clients permanently unhappy. If the multimedia stream for remote playback is stored on the server at a common lowest quality, high-bandwidth clients will receive poor quality despite the availability of a large amount of bandwidth. However, if the multimedia stream is stored at a single high quality encoding on the server, for many low-bandwidth clients the high loss rate will make the remotely played stream quality not acceptable. Another solution may be for the service providers to allow for their clients to choose between different already encoded streams at different bitrates at the beginning of streaming and to maintain the bitrate constant for the whole duration of the streaming process. Unfortunately the bottleneck that causes problems may be in the backbone for example at the provider's links to the multimedia server and therefore the user can not know the congestion level and its variation with the time. In consequence the customer cannot make a favorable choice between the offered different quality streams that would guarantee high quality reception for the whole duration of the streaming process.

Since static solutions seem to be unsuitable for a delivery environment in which the available bandwidth may change in time due to the presence of other traffic of various types (some of which has well-defined congestion control policies), it is necessary to propose a dynamic solution. The idea was to allow the server to adjust dynamically the quality of the multimedia streams it remotely plays so that the end-customers' perceived quality is as high as the available network bandwidth permits. This process was named in [197] quality adaptation.

The proposed Quality-Oriented Adaptation Scheme (QOAS) is an end-to-end quality adaptive solution for high quality - high bitrate multimedia streaming. It balances the network operators' and service providers' desire to increase their revenues by serving a high number of customers from limited network resources with the customers' interest in receiving high quality multimedia and paying a low price. Designing any quality adaptation scheme is a complex task. Normally, by increasing the number of remote simultaneous viewers of different content multimedia streams served by the same infrastructure, significant degradations in the end-users' perceived quality are expected mainly due to PUDLs. This is the case if the mechanisms employed to prevent or to minimise their effects are not highly effective (e.g. quality adaptation, post-processing techniques, etc.). QOAS was designed to try to prevent the PUDLs or to react to them. It

aims at maximising the end-user perceived quality and the links utilisation in the existing delivery conditions. It both varies the transmission rate and adjusts the streamed content if necessary.

In order to achieve high performance in variable delivery conditions by maintaining both good end-user quality and high links' utilisation, the QOAS is using the architecture presented in detail in section 4.2. The QOAS-based quality adaptation is performed by two mechanisms in conjunction: IntrA-stream QOAS and respectively InteR-stream QOAS. IntrA-stream QOAS is the main quality adaptation mechanism, involves streaming of the current multimedia clip only and creates between the QOAS server and the QOAS client a one-to-one relation. Its idea, similar to the classic quality adaptation, is to adjust dynamically the quality of the streamed multimedia, which in turn increases or decreases the quantity of multimedia data to be transmitted according to the feedback-reported state of delivery conditions. The adaptation is performed while maintaining continuity of the streaming process and the quality is varied in a well-controlled manner. In consequence the end-users benefit in their perceived quality compared to the alternative random losses that severely affect the quality of the streaming process [59]. InteR-stream QOAS, an extension of the IntrA-stream QOAS, aims at further improving both the end-user quality and the links utilisation. It involves a server-located controller module which is in permanent contact with all the individual IntrA-stream QOAS-s. The controller looks at the delivery process globally and makes fine adjustments to the IntrA-stream QOAS-s both in their initial stages and during their adaptive streaming. More details about the IntrA-stream QOAS and InteR-stream QOAS are given in sections 4.3 and 4.8.

4.2 QOAS-based System Architecture

This section presents the architecture of a multimedia streaming system that embeds the Quality-Oriented Adaptation Scheme (QOAS). First the overall architecture is presented at high level and then in more detail at block-level. Both the server's and the client's components are described separately later on.

4.2.1 High-Level Architecture

The high-level architecture of the adaptive multimedia system is presented in Figure 4-1. It includes multiple instances of QOAS-based adaptive client and server applications that communicate bi-directionally through the delivery network. They exchange multimedia data and control packets (including feedback messages). Since the system was aimed for streaming high

quality - high bitrate multimedia, the delivery network could be any broadband IP-based network such as broadband local area networks or all-IP multi-service delivery networks [1, 2, 11].

The QOAS client and server application instances implement the proposed adaptive multimedia streaming scheme, allowing for an adaptation process that involves the delivered stream only. The QOAS Client Application monitors some transmission-related parameters and the estimated end-user quality, allowing for its *Quality of Delivery Grading Scheme* (QoDGS) to compute accordingly scores that reflect the overall quality of the streaming process. These computed grades are then sent as feedback to the corresponding QOAS Server Application instance, whose *Arbiter* analyses them and proposes adjustment decisions in order to try to maximize end-user perceived quality in the given client-reported conditions. The QOAS Server Controller Application is in permanent contact with all the QOAS Server Application instances in order to allow for making fine adjustments to the adaptation processes by looking at the delivery process as a whole in this local delivery network. Its aim is both to improve the link utilisation and to minimise the transitory periods that may cause fluctuations in the end-user perceived quality. The Multimedia Database stores the multimedia streams in the pre-recorded streaming case, and some indexing information necessary to achieve high performance in the adaptation process.

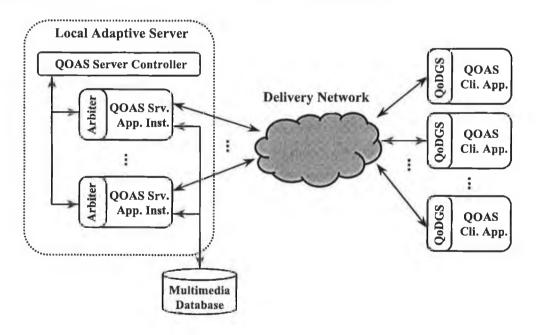


Figure 4-1 The architecture of the Quality Oriented Adaptation Scheme - based multimedia streaming system

4.2.2 Block-Level Architecture

A more detailed representation of the architecture of the QOAS multimedia streaming system at block level is shown in Figure 4-2. This architecture follows in general classic multimedia streaming system architecture designs that were presented or proposed in [197, 198, 199]. Its main components are the client-server communication modules, feedback-related units, multimedia (or audio/video) blocks, multimedia database system and quality-related components. As seen from above, the chosen architecture encompasses the major elements of distributed applications for multimedia streaming.

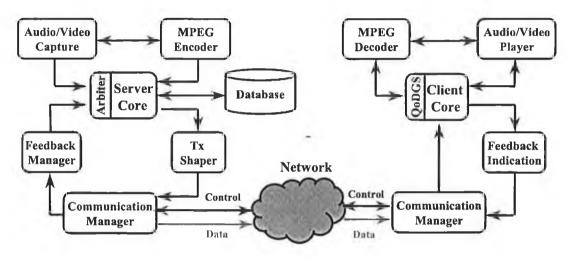


Figure 4-2 The block structure of the QOAS-based multimedia streaming system

The communication between the QOAS server application and the QOAS client application makes use of a double-channel link. A bi-directional connection is created between the client and the server when the first one sends a request to the server and the latter accepts it. This connection is used for the transmission of control messages, including the feedback ones. Next an unidirectional data link is established between the server and the client applications allowing for multimedia data transfer. This link is necessary to transmit time sensitive multimedia data faster, even though nonnecessary in a 100% reliable manner. The **Communication Managers** situated at both sides of the double communication link are in charge with requesting, respective accepting the requests and with establishing the double channel links. They control the transmission and play an important role in disconnecting the communicating partners upon request or when one of them is not reachable any more. They also perform buffering at both sender and receiver. To allow for the adjustments of transmissions if necessary, the server's Communication Manager co-operates with the Transmission Shaper whose functionality will be described later on. The client's Communication Manager

forwards both control information and data to the Client Core for processing and receives feedback data from the Feedback Indication Unit to be sent to the server.

For acquiring multimedia data, a special Audio/Video Capture Unit was provided as part of the server application. It works in conjunction with the Encoder Unit whose task is to control the multimedia data compression process using the chosen MPEG encoding scheme. These units are either actively involved in the live multimedia streaming case or work off-line prior to the actual streaming of pre-recorded clips. As part of the client application, the Decoder Unit is in charge with the decoding of the remotely streamed multimedia data, which is then played by the Audio/Video Player. This unit separation is only conceptual, allowing for a software and/or a hardware solution for decoding and playing.

The main goal of the **Feedback Indication Unit**, situated at the client, is to collect the grading scores computed by *Quality of Delivery Grading Scheme* (QoDGS) related to the quality of streaming. Details about QoDGS will be given in the next chapter. The **Feedback Indication Unit** assembles these scores in the feedback control messages and transmits them to the server, informing it about the quality of the reception. The **Feedback Manager**, located at the server, receives the feedback messages, extracts the scores and sends them to the *Arbiter (Arb.)* for the analysis. The Arbiter takes decisions concerning the transmission according to their values. The Arbiter's detailed functionality will also be presented in the next chapter.

One of the **Transmission Shaper**'s major goals is to reduce the burstiness of the transmissions (as the chosen encoding scheme MPEG is well known for its bursty streams [200, 201]), in environments that require a flat-like data flow (e.g Internet). However this can be switched off in situations when statistical multiplexing effect of data originating from multiple bursty sources does not have a disturbing effect on the end-quality of the multimedia streamed clips (e.g. local IP networks where multimedia accounts for the majority of traffic). The other main goal of the Transmission Shaper is to allow for the adaptive modification of the transmission rate after the analysis of the feedback received from the QOAS client. A modified token bucket mechanism [202] to which a variable token generation procedure has been added is used to allow for an adjustable transmission process. Both the token creation and the server side buffering (done prior to actual transmission of data) are feedback-controlled.

The **Server Core** is responsible not only for inter-connecting the other components, but also for applying the QOAS policy through its *Arbiter*. The Arbiter's role is to assess the quality of

streaming as reported by the QOAS client through its feedback messages, to take the adjustments decisions, if and when necessary and to apply them with the help of the Transmission Shaper.

The Client Core inter-connects the other client-located blocks in a similar manner with its counter-part, the Server Core. However the most important role in the QOAS-based architecture is played by the *Quality of Delivery Grading Scheme* (QoDGS). QoDGS monitors permanently both transmission related parameters and end-user quality and grades regularly the overall quality of delivery in terms of Quality of Delivery (QoD) scores. These scores are sent as feedback to the server and are used for adaptation.

More details about both the Server's Arbiter and the Client's QoDGS are presented later in this chapter.

The **Database Unit**, as part of the server application, is mainly used to store the multimedia streams for pre-recorded stream transmission, after the multimedia data acquired by the Audio/Video Capture unit is compressed by the Encoder. The Database is also used for saving some indexing information related to each stream, allowing for achieving high performance while applying the QOAS-based server adjustments. More details are presented in this chapter when describing QOAS and in the fifth chapter when presenting implementation issues.

4.3 IntrA-Stream QOAS

IntrA-stream QOAS has three main components: the client-located Quality of Delivery Grading Scheme (QoDGS), the Server Arbitration Scheme (SAS) and the Data Transmission and Feedback Mechanisms. Detailed information about these components and their functionality are given in separate sections of this chapter. Since the IntrA-stream QOAS is the main mechanism of QOAS, in this section the same abreviation QOAS is used to name it, unless explicitly stated otherwise. The following terms are also used: "server" - naming an instance of the QOAS Server Application and "client" - refering to the QOAS Client Application instance.

The IntrA-stream QOAS or QOAS is a feedback-controlled end-to-end adaptive scheme that relies on both long term and short term monitoring and assessment of some transmission parameters and of the end-user quality. This is performed by the QoDGS, which also regularly grades the quality of the ongoing streaming process in terms of Quality of Delivery (QoD) scores. These scores are sent using the QOAS's Feedback Mechanism to the server whose Arbiter processes them. The Arbiter takes into consideration the values of a number of recent feedback

reports, analyses them and suggests adjustment decisions to be taken by the server if necessary. These decisions affect in a controlled manner the quantity of streamed multimedia data and - in consequence - its quality. Figure 4-3 describes graphically the functionality of the QOAS, presenting also an example of a possible quality adaptation scenario during streaming of a multimedia clip.

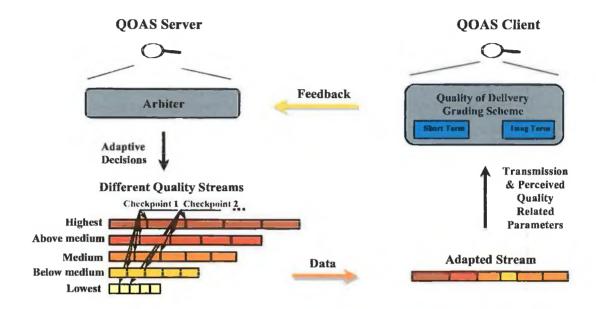


Figure 4-3 Schematic description of QOAS's adaptation principle

QOAS requires the definition of a number of different server states for each multimedia streaming process. For example the five-state model presented in Figure 4-4 was used for the experimental tests. Each server state is then assigned to a different possible stream quality version. The stream quality versions differ in terms of compression-related parameters (e.g. resolution, frame rate, colour depth) and therefore have different bandwidth requirements. They also differ in the consequent end-user perceived quality if presented as they are. For QOAS, the more server states are defined and therefore the greater the number of different stream quality versions associated with them, the better the adaptation process becomes. In the pre-recorded streaming case this is done at the expense of increased storage space in the server's Database. For live streaming, the granularity of the adaptation can be much higher and therefore there could be a high number of server states. The only limitation is introduced by the equipment or software that performs the real-time encoding.

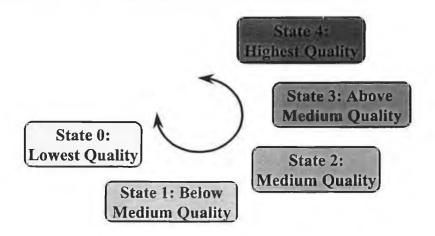


Figure 4-4 A five-state model that could be used by the QOAS's server

The server always has a current state which determines the quality of the multimedia clip to be streamed. During transmission the server dynamically varies its state according to the reported end-user quality of the streamed video. As previously mentioned the client-located QoDGS monitors and analyses the effect of delivery conditions on end-user perceived quality and quantifies it in terms of QoD scores. The detailed functionality of the QoDGS is presented in section 4.5. These computed QoD scores are sent regularly via a Feedback Mechanism described in section 4.7 to the server that takes the necessary adjustment decisions as presented in section 4.6. For example, when increased traffic in the network affects end-user quality, the server switches to a lower quality state which therefore also reduces the quantity of data sent, helping eliminate the congestion. If the client reports improved viewing conditions, the server gradually increases the quality of the delivered stream. The quantification of end-user quality is done using a metric that is described in section 4.4.

The QOAS also includes a mechanism to adjust the transmitted quantity of data. For a smooth play-out at the client, not only the starvation of the remote player (which forces it to stop play-out and start buffering) has to be avoided, but also jumps from one scene to another during adaptive measures. Therefore switching the quality of the transmitted source from the current one to a new one is done at well-determined checkpoints, as shown in Figure 4-5. This aims at keeping the skew between the two sequences at the client side as low as possible. At the same time, the mechanism is aware of the particularities of the encoding scheme. Since for testing the MPEG-2 encoding scheme is used, the implementation takes into account the MPEG I-P-B frame-based stream structure.

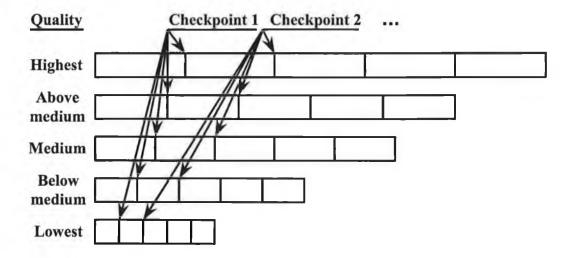


Figure 4-5 Switching between different quality streams with the same multimedia content is performed at certain checkpoints

The checkpoints are defined at the beginning of each Group of Picture (GOP), ensuring a high end-user perceived quality. In this manner the next to be transmitted is the I frame that can be decoded independently from other frames, and can constitute a good reference to the next temporal encoded frames. If the switch had been done in a position where a P or B frame is next to be transmitted, the remote decoder would have had problems re-creating the correct version of the actual frame referencing data which belongs to another quality stream. In the absence of the correct frames the temporal encoded data refers to (previous and/or next I or P frames), the decoder will use the existing frames, producing a low quality result.

Switching between different quality streams at the beginning of GOPs is simple for live transmissions and can be easily done during streaming. The determination and usage of checkpoints are more complex for the pre-recorded case and are performed in two phases. During a pre-processing phase performed only once for each stored multimedia clip, the server quality states are defined and the different quality streams are associated with them. Next the checkpoints' positions are determined and stored into the server Database. This process was named *registration*. The second phase is performed at every transmission and consists of fast retrieval of the checkpoints' positions from the Database if and when the server Arbiter suggests the adaptation to be performed. This is the *look-up* phase.

4.4 Q - End-User Quality Assessment

Different factors may affect the end-user perceived quality of the remotely streamed multimedia clips. By using the QOAS, the streams may suffer also bitrate variations that further affect the end-user's perceived quality. Therefore there is a need to quantify the perceived quality of the streams, affected both by bitrate variations and losses during transmissions, in order to determine the right balance between the server adaptations and end-user quality. Also it is significant to be able to assess the results of the adaptive streaming in terms of a well-known subjective scale, easy to relate to.

Rating	Impairment	Quality
5	Imperceptible	Excellent
4	Perceptible, not annoying	Good
3	Slightly annoying	Fair
2	Annoying	Poor
1	Very annoying	Bad

Table 4-1 Quality scale for subjective testing

In the second chapter many proposed quality metrics and some existing scales for assessing the multimedia streams' quality were presented. It was also commented on their relative advantages and disadvantages if used both during the QOAS adaptation process and for the final assessment of the scheme's quality-related performance. Due to the good balance between simplicity and information content, the five-point (1-5) subjective testing scale defined by the ITU-T-R P.910 [63], which is presented in Table 4-1 was chosen. Also, in order to map the end-user quality during the adaptive streaming on the selected 1-5 subjective scale and since the MPEG encoding scheme was used for testing, the multimedia quality metric Q proposed in [133] was used. Q describes the joint impact of MPEG rate and data loss on video quality. Its formula is presented in equation (4-1).

$$Q = Q_0 + \chi_q * \left(\frac{\overline{R}}{\chi_r}\right)^{-\frac{1}{\xi_r}} + \chi_l * \overline{R} * PLR$$
(4-1)

In equation (4-1) PLR is the packet loss ratio, \overline{R} is the stream's mean bitrate, the constant Q_0 has a value close to the maximum quality 5, χ_q , ξ_r and χ_r are related to the complexity of the sequence and χ_l depends also on the average bitrate.

Some of the most important advantages of using Q are:

- i) the possibility for its in-service usage based on fact that is a no-reference metric
- ii) its direct output on the ITU-T 1-5 scale without another mapping stage that may reduce the measurement accuracy,
- iii)its relatively simple formula requires few computations that can be performed very fast and without loading excessively the client machine during the grading process,
- iv) it uses parameters that are easy to monitor and
- v) it provides a good representation of the expected evolution of the perceived quality with the variation of loss rate and respectively stream bitrate (see Figure 4-6).

At the same time, some of the main disadvantages of using O are:

- i) using many constants that have values related to the streams' complexity, the same formula for Q may not describe best the quality of different multimedia clip types (e.g. high motion content, cartoons, etc.),
- ii) being simple Q may not fully describe the relation between (mainly) transmission related errors and the end-user perceived quality and
- iii) being proposed for MPEG-encoded streams, Q is not independent from the encoding scheme, requiring the MPEG compression for obtaining significant results.

Since the advantages overcome the disadvantages, it was decided to use Q for assessing the end-user perceived quality during both QOAS adaptation and adaptive streaming results analysis.

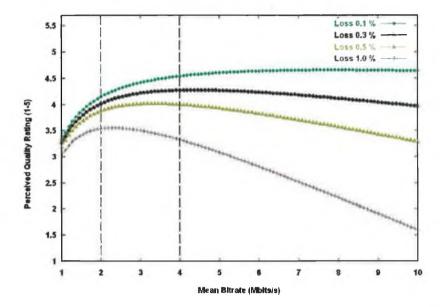


Figure 4-6 The end-user quality (Q) variation with the mean bitrate for a multimedia stream with average motion content, plotted for different packet loss ratios in the interval [0.001, 0.01]

The curves in Figure 4-6 show the evolution of the end-user perceived quality as measured by Q with the multimedia stream's mean bitrate for packet loss ratios between 0.1% and 1.0% when using the average values for parameters related to the stream's complexity suggested in [133]. As expected, the higher the loss ratio is, the lower is the end-user perceived quality. This fact supports the proposed adaptation policy of reducing the transmitted stream quality during congested periods. The resulting reduction in loss will yield improved end-user perceived quality. In normal traffic conditions, characterised by low loss ratios, the transmitted stream quality can be upgraded and an increase in the perceived quality is again obtained.

Since for very low loss rates (less than 0.1%), the benefit in the perceived quality with the increase in the stream bitrate above 4 Mb/s (and consequent bandwidth consumption) is not significant, the higher limit of interest for the MPEG encoding rate was chosen 4 Mb/s. Encoding multimedia below 2 Mb/s makes the perceived quality to drop below the "good" level even in very good delivery conditions and therefore 2Mb/s was selected as the lowest limit of interest for the MPEG encoding rate. Since the experimental testing was performed with MPEG-2 encoded streams with bitrates between 2 Mb/s and 4 Mb/s, the corresponding region of interest is delimited in Figure 4-6 by dashed lines.

4.5 Client-Located QoD Grading Scheme (QoDGS)

4.5.1 QoDGS Overview

One of the most important components of the QOAS is the client-located Quality of Delivery Grading Scheme (QoDGS), since on its functionality relies the performance of the whole adaptive scheme. Its goal is:

- to monitor continuously the streaming process,
- to collect both transmission performance related data and information related to the end-user perceived quality,
- to analyse the data gathered over a recent period of time,
- to grade the streaming process on a pre-defined scale computing QoD scores.

The resulted QoD scores are regularly sent by the feedback mechanism to the server whose Arbiter analyses them and takes adjustment decisions in order to improve the quality of delivery.

4.5.2 QoDGS Principles

In order to have a higher degree of confidence that the proposed QoDGS reflects the quality of delivery as accurate as possible, some design principles were formulated.

- 1. The QoDGS allows for both a long term and a short term monitoring of the monitored parameters related to the quality of the streaming process. Short-term variations are important for learning quickly about transient effects, such as sudden traffic changes or operating system/software problems, and for reacting as fast as possible to the resulting effects. Long-term variations are monitored in order to track slow changes in the delivery environment (e.g. new users in the system). The only difference between the two sets of collected data is the duration of the collection period. A suggested length is an order and respectively two orders of magnitude greater than the feedback reporting time. The analysis of the collected data and the corresponding partial grading are performed separately. As a result two different QoD grades are computed: one for long term: QoD_{LT}, the other for short term: QoD_{ST}.
- 2. The QoDGS takes into consideration all the parameters chosen to be monitored such as: the one-way delays, delay jitters, the loss rate and Q, but it is also very important to

allow for the possibility to consider also other parameters related to the quality of streaming. This is significant both for testing purposes and for eventual extensibility and improvement of the QoDGS.

- 3. The QoDGS allows for considering different types of parameters and for taking into account their characteristics. For some parameters it is very important to have a value as low as possible, for others a steady value is better. For some parameters there are no low and high limits for their values, which depend on the network topology, network state, streaming session, etc., for others, where a percentage describes their variation, such limits naturally exists. It is important for the QoDGS to provide a mechanism to accommodate all these different particularities of the monitored quality of delivery related parameters.
- 4. The QoDGS takes into consideration the relative importance of each of the monitored parameters in comparison to the others. The proposed solution uses a weighting mechanism. The best values for the weights associated with each monitored parameter have to be determined during a detailed tuning phase, prior to the deployment of the QOAS system that implements the QoDGS. The tuning aims at achieving high QOAS performance in terms of adaptiveness and stability.
- 5. The QoDGS allows for the consideration of different importance of the short term and long term grading processes. The solution proposed is to associate different weights to each of them according to their relative importance. As a result, the computation of the overall QoD score takes into account these associated weights. For good tuning of the QoDGS, a detailed testing phase is necessary to be performed.
- 6. The QoDGS grading process is to be performed very fast in order not to influence in a negative way the performance of the overall multimedia client application. Also it has to allow to be computed at any time, independent from the packet receiving process, the multimedia data decoding procedure or the stream play-out process. In this way the frequency of the feedback control packets that carry the computed QoD scores can be easily modified. A high value for the feedback frequency will overload both the client and the server, which have to be able to compute the QoD scores and to analyse them and take decisions, to send and to receive the control packets and to send or receive data packets. Apart from this, the network itself will be overloaded, making the stream's data transmission more difficult. A low value for the feedback transmission frequency

will not permit the system to react fast enough to the changes in delivery conditions. In consequence a compromise value must be found during the testing phase.

7. The QoD scores computed by the QoDGS should not be dependent on any other parameters apart from the ones related to the streaming process (e.g. the machine, the processor's load, the type of connections, etc.). This is to have a deterministic behaviour of the QoDGS and therefore of the QOAS-based multimedia streaming system that uses it.

4.5.3 Monitored Parameters

In order to build the QoDGS it was important to determine which parameters related to the performance of streaming process are correlated to the end-user perceived quality, in which way and how strong is this link. It was equally important to determine if some metrics exist that could be used to measure these parameters, how much effort their computation takes and whether by using them the QoDGS is closer to achieve its goals and to follow the stated principles.

Analysing the International Telecommunication Union (ITU)⁵⁰ and the Internet Engineering Task Force (IETF)⁵¹ proposals, one could notice that, although they have common goals, they tend to have different paths to achieve them. Since the ITU tends to evaluate services in general and their quality in particular, its metrics can be used for assessing the stream quality. The IETF is more network-oriented and since our interest is closer to the IETF's, we have tried to determine a working group within IETF that focuses on studying performance parameters that, if monitored by QoDGS, could give significant information about the state of the network and its potential effect on the end-user perceived quality.

IETF IP Performance Metrics (IPPM) Working Group²⁶ has proposed a set of standard metrics that can be applied to the quality, performance and reliability of data delivery over networks. The set of metrics defined in their framework that offer some solutions for unbiased quantitative measures of performance are: connectivity, one-way delay and loss, round-trip delay and loss, delay variation, loss patterns, packet reordering, bulk transport capacity and link

⁵⁰ The International Telecommunication Union (ITU), http://www.itu.int

⁵¹ The Internet Engineering Task Force (IETF), http://www.ietf.org

bandwidth capacity [62]. Their potential usage by the QoDGS was assessed and the conclusions are presented briefly next.

Connectivity refers to the fact that a host is reached or not by a data packet sent to it [203]. Although obviously very important for the multimedia streaming, the connectivity-related problems are taken care of by the Communication Managers, part of the QOAS Architecture in charge with establishing and controlling the communication and by the server-located Arbiter, part of the QOAS. This is described in details in sections 4.6 and 4.7.

One-way delay between two hosts is defined in [204] as the time between the moment when the first bit of a packet was sent from the first host to the second one and the moment when it reaches the second host. It is expected that one-way delay and especially delay variation to be correlated with packet loss, which in turn has a strong influence on the end-user perceived quality for the case of multimedia streaming as reported in [196, 205] and mentioned in the second chapter. This is because whenever packets are delayed in the network, they are regularly stored in either router queues or in buffers that have a finite capacity. In consequence if one-way delay increases high enough, loss occurs as those queues or buffers become full. However, if there is enough storage capacity to absorb considerable delay variations, this correlation weakens. The relationship between one-way delay and loss is further weakened by the fact that delay and delay variation are the result of a repeated concatenation of variations at each hop, while loss is caused by one or few overloaded elements along this path. Hence, many elements will contribute to delay and delay variation, but not also to loss. Many researchers have studied the one-way delays [154, 195, 205, 206, 207, 208, 209] and their conclusion is that although the linkage between delay on one side and loss and consequently quality of service is not very strong, it cannot be neglected. Especially when some of them [154] that have studied also the degree to which the delay reflects the available bandwidth found out that the ratio between the delay a packet incurred due to its connection's own loading of the network path, versus the total delay it incurred correlates very well with the overall throughput achieved by the connection. Others [209] have even found a direct connection between the increase of the one-way delay and the available bandwidth. These findings and the results of other works that have used previously one-way delay in adaptive streaming that were presented in the second chapter have suggested taking the one-way delay into account when choosing the parameters monitored by the QoDGS.

One-way loss related to a packet transmitted between two hosts is defined in [210] as 0 if the packet transmitted has reached its destination and 1 otherwise. In practice the one-way loss is measured over a period of time and is expressed as a percentage of the total number of packets sent.

The loss appears due to the fact that when the packets transmitted over the network are delayed, they are regularly stored in either router queues or in buffers with a capacity that should normally accommodate them. In highly increased traffic conditions, the storage space in these intermediate network elements is exceeded and the data that cannot be saved is lost. There is an important connection between data loss and different application performances that are significantly degraded with the increase in loss rate [154, 205, 207, 210]. This is especially valid for time-sensitive applications (including multimedia streaming ones) as reported in [133, 196]. The end-user quality is the most affected by a phenomenon called *error propagation* particularly lately when compression based on reducing both spatial and temporal redundancies [201] (e.g. MPEG-1 [73], MPEG-2 [74], etc.) is used to diminish the quantity of data to be transmitted. Based on these considerations and on results of other research that have successfully used one-way loss rate in quality adaptation schemes [211, 6, 7], loss was considered a significant input parameter for the QoDGS.

Round-trip delay between two hosts is defined in [212] as the time between the moment when the first bit of a packet was sent from the first host to the second one and the moment when it reaches again the first host, after the packet was received by the second host and immediately was sent back to the first host. This metric was introduced to complement the one-way delay since for some applications this is the quantity of interest, it is simpler to compute and it can be determined more accurately. Unfortunately in general the path from a source host to a destination may differ from the path from the destination back to the source ("asymmetric paths"), such that different sequences of routers are used for the forward and reverse paths. Therefore round-trip measurements actually measure the performance of two distinct paths together. Also, even when the two paths are symmetric, they may have completely different performance characteristics due to asymmetric queuing. On top of this, the performance of some applications, especially multimedia streaming ones, depend mostly on the performance in one direction and therefore the measurements of the round trip delay may not describe accurately enough the existing network situation the traffic of interest may have to face. In consequence the decision was taken not to monitor round-trip delay in the QoDGS.

Round-trip loss was defined⁵² as the percentage of the packets sent by a host to another host (meant to answer by sending a packet back) that were not followed by a corresponding received packets from the total number of packets sent. Round-trip loss, although was listed by the

⁵² BT Ignite, Web site, http://ippm.ignite.net/more info.html

IPPM Working Group as a metric of interest, has not made the working group members to propose a RFC yet, due to its lower usage interest in comparison with the one-way loss metric. Therefore, taking also into account the probability of dealing with asymmetric paths for which the round-trip loss will not give significant information related to the path the multimedia data takes to the destination, the round-trip loss was not considered important to be monitored by the QoDGS.

One-way delay variation for a pair of packets in a stream of packets was defined in [213] as the difference between the one-way-delays computed for the selected packets. As previously mentioned, in networks the one-way delays vary much due to the routers' queuing and/or the usage different paths by the packets to reach the destination. Since the network congestion is a phenomenon that builds up by increasing the number of packets trafficked trough the network forcing the routers to queue them and in consequence to introduce increasing delays, delay variation (or delay jitter [214]) is an important metric that signalises such a situation [215]. The effect of a highly variable delay jitter was also studied, especially in relation to time-sensitive applications, including those that stream multimedia. The conclusion was that although these applications would best perform if the delay was constant for all the packets, a certain variation can be coped with using receiver buffering. Unfortunately when the delay jitter exceeds a certain threshold, the received buffering is not enough and the performances of the applications are severely affected. For the case of multimedia streaming applications, a highly variable delay causes loss of data by either buffer over-run or under-run, significantly reducing the end-user perceived quality. This was confirmed by researchers that have studied the delay variation [154, 195, 205, 206, 207, 208]. Their conclusion that there is a certain correlation between the delay variation and loss rate (and consequent quality of service) that neither can be fully described, nor can be neglected, makes us to suggest to use delay jitter as one of the monitored parameters by the QoDGS.

Loss pattern (or loss distribution) is a key parameter for certain real-time applications (e.g. multimedia-based ones) that determines the performance observed by the users. For the same loss rate, two different loss distributions could potentially produce different perceptions of performance [216]. The impact of loss pattern is also extremely important for non-real-time applications that use an adaptive protocol such as TCP. Research results that demonstrate the importance and existence of loss burstiness and its effect on packet voice and video applications are published in [217, 218, 219, 220]. In [216] two metrics, named "loss distance" and "loss period", were defined to describe the loss pattern. The "loss period" metric captures the frequency and length (burstiness) of loss once it starts, and the "loss distance" metric captures the spacing between the loss periods. The QoDGS

takes into account the loss pattern through its short-term and long-term grading mechanisms. These mechanisms try to consider also the patterns of other parameters' variations.

Packet reordering was considered a performance issue of certain importance since has determined the IETF IPPM Working Group [61] to propose a metric subject of an Internet Draft [220] that seems likely to become a RFC soon. A reordering metric is relevant for many applications, but significant only for time-sensitive ones and only when the extent of reordering affects the applications' performance. In general packet order is not expected to change during transmission from a host to another one, but there are cases when it does change. For example when a single packet stream is sent from a host to another one between which there are two paths, one with slightly longer transfer time, the packets traversing the longer path may arrive out-of-order. The ability to restore order at the destination will likely have finite limits and mainly due to the receiver buffers' finite size in terms of packets, bytes, or time. Also it is important to quantify the extent of reordering, or lateness, in all meaningful dimensions. Since the percentage of out-of-order packets from the total number of packets sent in the measurements carried out and reported in [207, 221] was very low, our decision was for the QoDGS not take this parameter into account. This is also supported by the fact that the target network the multimedia streaming application the OoDGS is aimed for has little or no parallel paths that could constitute a cause for the out-of-order arrival of packets. However, the extensibility of the QoDGS's design should allow for adding the percentage of out-of-order packets as monitored parameter if QOAS-based multimedia system is meant to be deployed in the Internet where lately, due to the increase in the number of parallel paths, the packet reordering is more common than thought [222].

Bulk transport capacity (BTC) metric, as defined by IPPM WG in [223] was meant to measure a network's ability to transfer significant quantities of data with a single congestion-aware transport connection (e.g., TCP). The intuitive definition of BTC is the expected long-term average data rate (bits per second) of a single ideal TCP implementation over the path in question. Although there was some interest for BTC [224] and it may be useful for some applications, for multimedia streaming applications the transport capacity using a reliable transport protocol is of little interest, more significant being the timely arrival of data which affects more the quality of service. In consequence the BTC was not taken into account as a parameter to be monitored by the QoDGS.

Link bandwidth capacity is an important metric for many applications, but of more interest is the *available bandwidth* related to a link. There is significant research in this direction [209, 225, 226] that has to face problems of scalability, intrusiveness, accuracy and high computation related to the determination of the available bandwidth at any moment. Apart from

this, IETF IPPM WG's [61] Internet Draft referred in [227] suggests a method to measure the available bandwidth of a path using an active approach that probes the path using TCP New Reno. Since there is no general accepted metric or mechanism to determine the available bandwidth at any moment with significant accuracy, it was not taken into account directly as a parameter for the QoDGS. However, the InteR-stream QOAS uses an estimation of the link bandwidth capacity that is described in detail in section 4.8.

In conclusion, after assessing these performance metrics proposed by IETF IPPM Working Group, the decision was taken to monitor and to grade the one-way delay, the delay variation (jitter) and the one-way loss rate by the QoDGS. The loss pattern and the other parameters' variation patterns are taken into consideration in the grading scheme while the percentage of out-of-order packets is allowed to be taken into account in future, if the target network for QOAS-based systems is different. The QoDGS makes also use of the Q metric, which was described in detail in section 4.4 in order to assess the end-user quality during the streaming process.

4.5.4 Measurements Accuracy

The one-way delay, the delay jitter and the one-way loss rate were the parameters taken into account for monitoring by the QoDGS. As mentioned in [204, 210, 213] there is a significant problem when measuring these metrics due to their sensitivity to clock-related errors and uncertainties. There are two types of errors and uncertainties: i) due to the difference between the wire-time and the hosts' clocks times or between the real time (UTC) and hosts' clocks times and ii) due to uncertainties in the clocks of the source and the destination hosts' clocks. These problems are summarised next, according to their source: clock wire-time, clock offset, clock synchronisation, clock accuracy, clock resolution and clock skew.

The wire-time was defined as the time at which a packet appeared on a link, without exactly specifying whether this refers to the first bit, the last bit, etc. Unfortunately there are metrics defined using wire-time, which has to be related to the host's clock time, process that may introduce errors. QoDGS uses only IETF IPPM WG - defined metrics that do not introduce this kind of errors.

If there is a significant interest in the high accuracy of the results related to the real time (universal time clock - UTC), another source of error may be caused by the **clock offset** which represents the difference between the time reported by the clock and the "true" time as defined by the UTC at a particular moment. Since the QOAS does not relate its results to the UTC, there will not be such errors.

If the source host's clock and the destination host's **clock** are not **synchronised**, this will cause an error in the delay measurement. The source clock and the destination clock have a synchronisation error of Tsynch if the source clock is Tsynch ahead of the destination clock. Thus, if the value of Tsynch is known exactly, the clock synchronisation error could be corrected by adding Tsynch to the uncorrected value of Tdest-Tsource. In practice the synchronisation error is not known precisely (and varies with the time) and therefore the synchronisation of the two hosts' clocks is recommended prior to the measurements.

The **clock accuracy** is important only in identifying the exact time at which a given delay or loss was measured. Clock accuracy as is has no importance to the accuracy of the measurement of delay or loss. When computing delays, including in the QoDGS case, only the differences between clock values are interesting and not also the values themselves.

The **clock resolution** adds a certain uncertainty about the time measured with it. For example if the source clock has a resolution of 10 msec, an uncertainty of 10 msec is added to any time measured with it, including the ones that are used to compute the one-way delays.

The **skew of a clock** is not so much an additional issue as it is a realisation of the fact that Tsynch is itself a function of time. Thus, if Tsynch is to be measured or bound, this needs to be done periodically.

Since both the hardware and the software computer clocks of both the source and the destination hosts are poor timekeepers [228], a good practical solution that both keeps the clocks' skews to minimum and maintains them synchronised is to periodically synchronise the clocks with a third, more reliable clock. For a very precise synchronization, special arrangements that include GPS, a local atomic clock or an ISDN synchronous clock board are needed [229]. However multimedia streaming deals with millisecond order delays, so NTP protocol [230] can be used for synchronizing both clocks separately with a third external clock. This is the approach QOAS takes for maintaining a good level of accuracy in the measurements related to the one-way delay, the one-way loss and the delay jitter and it uses the U.S. atomic clock located in Boulder – Colorado, USA^{53,54,55} or any other public NTP time server⁵⁶.

⁵⁴ National Institute of Standards and Technology, USA, Atomic Clock, http://www.boulder.nist.gov/doc-tour/atomic_clock.html

⁵³ The Official U.S. Time, http://www.time.gov

4.5.5 QoDGS Design

The client-located QoDGS was designed according to the principles previously mentioned in section 4.5.2. Figure 4-7 presents the QoDGS block structure. This figure also shows the parameters taken into consideration by the QoDGS for monitoring, as presented in section 4.5.3, in order to assess the quality of the streaming process and grade it in terms of QoD scores. It is assumed that the server and the client clocks are synchronised all the time (see section 4.5.4), assuring therefore the measurements accuracy.

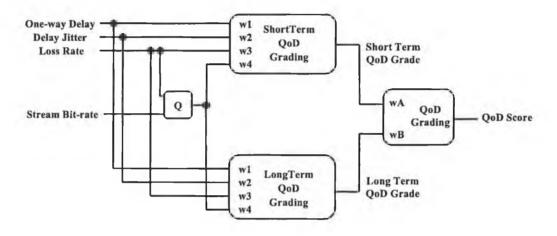


Figure 4-7 QoDGS takes into consideration both traffic-related parameters and end-user quality

The QoDGS consists of three stages. In the first stage QoDGS both grades the instantaneous values of the monitored parameters (one-way delay, delay jitter and loss rate) and saves session-specific information related to each parameter. This allows for the corresponding partial scores to be more precisely computed next time during the grading process. The first stage also involves the computation of the multimedia quality metric Q whose formula was presented in equation (4-1). The computation of the Q metric makes use of the bit-rate of the streamed multimedia clip and the loss rate. The partial scores computed during this first stage are saved in different length sliding windows. Based on them, the parameters' short-term and long-term variations are assessed in the second stage of the QoDGS. This second stage takes into account the relative differences in the importance of the monitored parameters in relation to the characteristics of the delivery architecture by weighting their contributions. Finally short-term (QoD_{ST}) and long-

⁵⁵ Atomic Clock Time Server, time-a.timefreq.bldrdoc.gov (132.163.135.130, 132.163.4.101), NIST Boulder Laboratories, Boulder, Colorado, USA

⁵⁶ Public NTP Time Servers, http://www.eecis.udel.edu/~mills/ntp/servers.html

term (QoD_{LT}) grades are computed. In the third stage, QoD_{ST} and QoD_{LT} scores are used to compute the overall score (QoD_{Score}) in a weighted process that accounts for their relative importance.

Next sections present more information about the QoDGS, describing in detail each of the three QoDGS grading stages.

QoDGS - First Grading Stage

The one-way delays, the delay jitter, the loss rate and the estimated end-user quality (Q) are the parameters under permanent monitoring by the QoDGS. Therefore the QoDGS watches out for all the events that influence their values such as arrivals of data packets and modifications in the streamed clips encoding bitrates. These events trigger computations of one-way delays by taking into consideration the packets' timestamps as suggested in [204], of delay jitters based on the computed one-way delays as presented in [213] and of loss rate by looking at the packets' sequence numbers as in [210]. Changes in the streams' encoding rates and the computed loss rates are used to compute the Q metric as presented in equation (4-1). These measured values are used in this stage both to grade the monitored parameters' variation and to update statistical information related to this variation. Details about these grading processes are presented next for each monitored parameter that were individually taken into account as stated in the third QoDGS principle (section 4.5.2).

One-way delay

One-way delay is computed for each packet carrying multimedia data that has arrived at client as in equation (4-2). The resulted value is used as input for two similar Delay grading schemes whose block structure is schematically presented in Figure 4-8. These grading schemes consist of the *Delay Grading* unit that grades the one-way delay based on historic information related to its variation and a *Delay Statistics* unit. The historic statistics stored in the *Delay Statistics* unit are updated each time when a new packet arrives and new one-way delays values are computed.

$$Delay = TimeStamp_{Dest} - TimeStamp_{Source}$$
 (4-2)

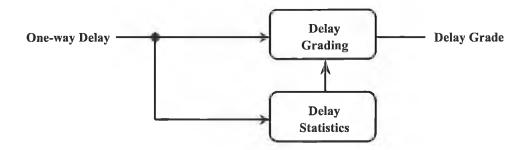


Figure 4-8 DelayGrade computation in the QoDGS first grading stage based on historic statistics about oneway delays

The fifth principle presented in section 4.5.2 states that QoDGS should take into account both short-term and long-term monitoring of parameters for better determination of their variation pattern. Since the statistical information necessary has to be collected during different periods of time for the two types of monitoring, it was necessary to have two Delay grading schemes.

Since there are no standards for relating one-way delay values to end-user perceived quality and there are not even general accepted recommendations for high and low limits for the one-way delay, building a Delay grading scheme is difficult. To make the situation worse, even research that have studied the one-way delay and have reported some acceptable values for it [208, 231, 232] could not give valid suggestions for any type of application or for any target network. In consequence for the Delay grading scheme a variable grading interval was used that spans between minimum and maximum delay values recorded by the *Delay Statistics* unit. Since the one-way delay values are subject to noise, the decision was to take into account the delay average computed for the duration of the monitoring (i.e. short-term and respectively long-term).

Equation (4-3) presents the formula used by the one-way *Delay Grading* unit for computing the *DelayGrade*. It considers the average of the one-way delay values (*AvgDelay*) in relation to minimum (*MinDelay*) and maximum (*MaxDelay*) delay, as recorded by the *Delay Statistics* unit. It also takes into consideration the minimum (*MinG*) and the maximum (*MaxG*) grades on the chosen scale (e.g. for ITU-T R P.910 five-point scale [63] they are 1 and respectively 5).

$$DelayGrade = MinG + (MaxG - MinG)*(1 - AvgVar)$$

$$AvgVar = \frac{AvgDelay - MinDelay}{MaxDelay - MinDelay}$$
(4-3)

Figure 4-9 shows the linear variation of the *DelayGrade* when *AvgDelay* varies from *MinDelay* to *MaxDelay* and consequently *AvgVar* varies from 0 to 1. The fact that the *DelayGrade*

decreases towards minimum *MinG* when the *AvgDelay* tends to the maximum recorded value is meant to punish increases in delay that indicate a possible build up of a network congestion.

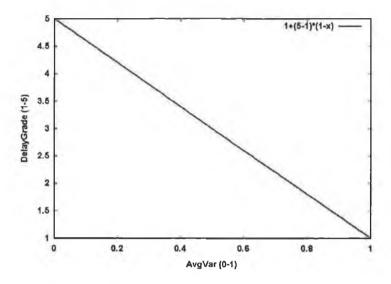


Figure 4-9 DelayGrade linear variation when AvgDelay varies between MinDelay (AvgVar=0) and MaxDelay (AvgVar=1)

Equation (4-4) lists the statistical information updates performed by the *Delay Statistics* unit. α is the update factor suggested in [233] for best performance in average delay estimation (the implementation uses a value of 0.9), *Delay* is the instantaneous value for the one-way delay as measured in equation (4-2) and *AvgDelay'*, *MinDelay'* and *MaxDelay'* are updated values for the indicated statistics.

AvgDelay '= AvgDelay *
$$\alpha$$
 + Delay * $(1 - \alpha)$
MinDelay '= min(MinDelay , Delay)

MaxDelay '= max(MaxDelay , Delay)

(4-4)

The AvgDelay is initialised every time when the monitoring interval (short-term or long term, respectively) elapses. The MinDelay and MaxDelay maintain their values during the whole streaming session in order to learn from historic behaviour and achieve good adaptive performance. MinDelay and MaxDelay are initialised for the first time with corresponding min-max values if there is enough information about the target network or with the first Delay value computed otherwise. AvgDelay is initialised for the first time with the first Delay value and then with the latest AvgDelay value recorded in the previous period.

Concluding the first grading stage of the QoDGS in relation to the one-way delay, one could say that its aim is to compute the *DelayGrade*, which is then used in the second grading stage.

Delay variation (jitter)

Delay variation (jitter) is estimated as in equation (4-5) after the arrival at client of each packet that carries multimedia data. This is performed with the *Delay Statistics* unit's help (part of the Delay grading scheme) that provides the average value for the one-way delay - *AvgDelay*. The resulted value for jitter is used as input for two Jitter grading schemes that allow for both short-term and long-term monitoring and grading, similar to those presented for the one-way delay grading. Their block structure is presented in Figure 4-10. The *Jitter Grading* unit grades the delay jitter based on delay jitter historical statistics as recorded by the *Jitter Statistics* unit. The information stored by the latter is updated every time when new delay jitter values are computed.

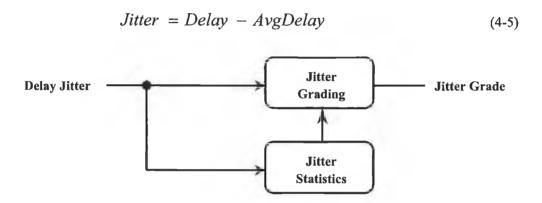


Figure 4-10 Delay jitter grading scheme that computes Jitter Grades in the first stage of QoDGS

Although there are works that take delay jitter into consideration [195, 205, 207, 231, 232] in relation to the end-user perceived quality, there is neither a widely accepted standard for the levels of jitter valid for any type of application or for any target network and nor graphs that would describe how the perceived quality decreases with the increase of jitter. However, there are works such as [208] that have suggested that there is a certain value for jitter after which the performance of the application (including multimedia) decreases sharply. This value depends on both the network and the application. In building the Delay jitter grading scheme this suggestion was taken into account and the squared Butterworth formula [234] shown in equation (4-6) was used, with median value *JThresh* - a threshold value for jitter - and n=3 in the *Jitter Grading* unit in order to compute the *JitterGrade*. Since the instantaneous values for delay jitter are subject to noise, the

average AvgJitter as computed by the Jitter Statistics unit for the duration of the monitoring (i.e. short-term and respectively long-term) is taken into account.

$$JitterGrad e = MinG + \frac{MaxG - MinG}{1 + \left[\frac{AvgJitter}{JThresh}\right]^{2 * n}}$$
(4-6)

In equation (4-6) *JitterGrade* is expressed on the ITU-T R P.910 five-point 1-5 scale [63], where MinG = 1 and MaxG = 5 are minimum and respectively maximum possible grades.

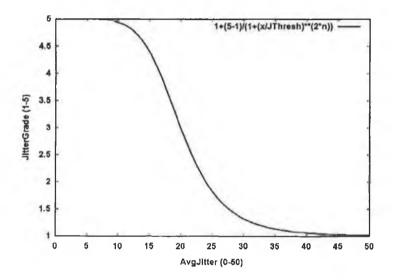


Figure 4-11 JitterGrade with AvgJitter variation between 0-50 ms (JThresh =20 ms, n =3)

In order to exemplify the effect of the squared Butterworth formula on the values of JitterGrade, Figure 4-11 shows the variation of the JitterGrade when AvgJitter varies from 0 msec to 50 msec and JThresh is 20 msec. It is significant to mention that one could divide the plot into three regions. For AvgJitter values smaller than JThresh the JitterGrade remains close to the maximum grade MaxG = 5. For values greater than JThresh the grade is close to the minimum value MinG = 1. In the threshold neighbourhood, the JitterGrade decreases sharply with the increase of AvgJitter.

$$AvgJitter' = AvgJitter * \alpha + Jitter * (1 - \alpha)$$
 (4-7)

Equation (4-7) presents the statistical information updates performed by the *Jitter Statistics* unit. α is the update factor suggested in [233] for best performance (the implementations use a

value of 0.9), *Jitter* is the instantaneous value for the one-way delay variation (jitter) as shown in equation (4-5) and *AvgJitter*' is the updated value for the jitter-related statistics.

In conclusion, the first grading stage related to the delay jitter aims at computing the *JitterGrade* which is then used in the second grading stage of the QoDGS.

Loss rate

Loss rate is computed with the simple formula presented in equation (4-8), each time when a multimedia data packet arrives at the client. This is performed simultaneous by two Loss Rate grading schemes that consider short-term and long-term evolution of the loss rate respectively. Their common block-structure is presented in Figure 4-12.

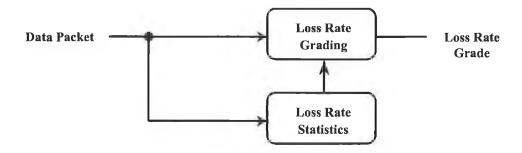


Figure 4-12 Loss Rate grading scheme computes Loss Rate Grades in the first stage of QoDGS

$$LossRate = \frac{TotalTxB - TotalRxB}{TotalTxB}$$
(4-8)

The LossRate is computed and stored for the duration of the short-term and of the long-term monitoring period respectively by Loss Rate Statistics units. They make use of the total number of bytes received by the client - TotalRxB and the total number of bytes sent by the server - TotalRxB in these periods. The units update these values every time when a new data packet arrives at the client using its sequence number and size fields. The Loss Grading units make use of the LossRate-s as computed in the equation (4-8) in order to perform the computation of the LossGrade-s as shown in equation (4-9).

$$LossGrade = MinG + \frac{MaxG - MinG}{\frac{4*LossRate}{3*LTarget}}$$
(4-9)

In this equation LossGrade is expressed on the ITU-T R P.910 five-point 1-5 scale [63], where MinG and MaxG are minimum grade 1 and respectively maximum score 5. The grading formula for LossGrade was chosen in a manner that allows for flexibility in deploying the QoDGS (by choosing the target loss rate LTarget), while maintaining the same policy of severely punishing loss rates that tend to get closer to the LTarget rate, regardless of the value of the target loss rate. The aim was also for the grades to tend to the MinG value, while being very close to it, once the loss rates have exceeded the LTarget value. For the target network QOAS is going to be deployed on and in the absence of any post-processing techniques that may accommodate greater loss rates, the LTarget was set to 1%.

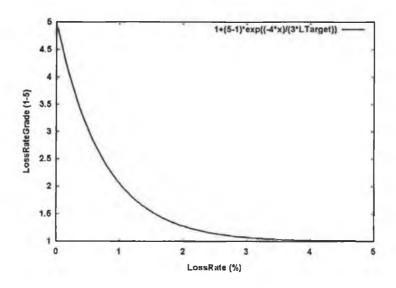


Figure 4-13 LossGrade variation when LossRate varies between 0 and 5 % for LTarget 1%

Figure 4-13 shows the variation of the *LossGrade* when *LossRate* varies between 0 and 5%, for this value of *LTarget*. One could notice that the *LossGrade* has a value of 2 on the 1-5 scale when the *LossRate* reaches *LTarget*, sharply dropping from maximum toward *LTarget* and then tending to minimum 1. This grading mechanism was designed for loss rate since existing research such as [235, 236], although extensive, does not agree on certain limits that would be applicable to all applications and any network.

In brief, the first grading stage that focuses on the loss rate computes the *LossGrade* which is then used in the second grading stage of the QoDGS.

End-user quality

End-user quality is measured by the metric Q [133], which describes the joint impact of MPEG rate and data loss on video quality whose formula is presented in equation (4-1). Figure 4-6 plots the variation of the Q value with the variation of video bitrate for certain loss rates. Section 4.4 also highlights the reasons Q was chosen for measuring the end-user quality.

For taking into consideration both short-term and long-term variations of end-user quality as estimated by Q, the computation of Q requires values of the loss rates for the corresponding monitoring periods. For retrieving these values, the Q grading scheme co-operates with the short-term and long-term Loss Rate grading schemes' *Statistics* units. This information and instantaneous streamed multimedia bitrates are used to compute the Q metric values that are used in the next grading stage of the QoDGS.

QoDGS - Second Grading Stage

The second stage in the QoD grading process is focused on taking into account the relative difference between the importance of the monitored parameters and on computing both the short-term score $QoD_{ST\ Score}$ and the long-term grade $QoD_{LT\ Score}$. Short-term variations are important for learning quickly about transient effects, such as sudden traffic changes or operating system/software problems, and for reacting as fast as possible to the resulting effects (e.g. high loss, excessive delays). Long-term variations are monitored in order to track slow changes in the delivery environment (e.g. new users in the system). Their effects are not evident on short-term and therefore longer monitoring periods are necessary.

Figure 4-14 presents graphically the short-term second grading stage of the QoDGS. The only difference between the short-term and the long-term grading procedures is the duration of the period the statistics related to the monitored parameters was collected for. The short term grading scheme focuses on the changes that occur on short term, regardless to what happens on a greater scale, whereas the long-term grading scheme, presented in Figure 4-15, grades variations that happens on longer time scale. These short-term and long-term periods are considered, respectively, an order and two orders of magnitude greater than the time between consecutive feedback reports.

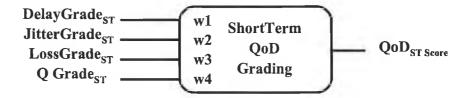


Figure 4-14 Short-term QoDGS second grading stage

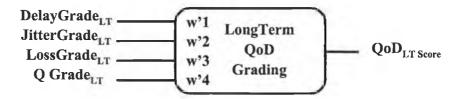


Figure 4-15 Long-term QoDGS second grading stage

Figure 4-14 and Figure 4-15 also show how each grade computed in the QoDGS's first stage in relation to a certain monitored parameter has associated a different weight w_i and respectively w_i^* . The more important the parameter i is, the higher the value of the corresponding weight w_i or w_i^* is, and therefore the higher the contribution of its grade in the overall computed scores: QoD_{ST Score} or QoD_{LT Score}. Equations (4-10) and (4-11) present the formulas according to which the QoD short-term and long-term scores are computed. The formulas are similar, but they use different values for both the weights and the grades that were computed for two different sets of statistically collected data.

$$QoD_{ST\ Score} = w_1 * DelayGrade + w_2 * JitterGrade$$

+ $w_3 * LossGrade + w_4 * QGrade$ (4-10)

$$QoD_{LT\ Score} = w'_1*DelayGrade + w'_2*JitterGrade + w'_3*LossGrade + w'_4*QGrade$$
 (4-11)

For accurate results, it is necessary to respect the conditions from equations (4-12) and (4-13).

$$\sum_{i=1}^{4} w_i = 1 \tag{4-12}$$

$$\sum_{i=1}^{4} w'_{i} = 1 \tag{4-13}$$

This QoDGS design provides a high degree of flexibility by defining two different sets of monitored parameters-associated weights for short-term and respectively long-term monitoring. In practice tuning such a system with so many variables is very difficult and we have suggested the use of identical sets of weights $(w_i = w'_i)$ in this second stage of the grading process.

QoDGS - Third Grading Stage

The third grading stage in this three-stage grading process combines the short-term and the long-term QoD grades computed in the previous stage, taking into account their relative importance. In order to allow for fast computation as stated in the sixth principle (section 4.5.2), two different weights w_A and w_B were associated with these scores. The final QoD_{Score} is calculated according to the formula presented in the equation (4-14), with values for w_A and w_B that respect the condition presented in equation (4-15).

$$QoD_{Score} = w_A * QoD_{ST Score} + w_B * QoD_{LT Score}$$
 (4-14)

$$w_A + w_B = 1 \tag{4-15}$$

The computed QoD_{Scores} are sent to the server via the Feedback Mechanism that is described in section 4.7 of this chapter and used by the server Arbitration scheme as described in section 4.6 to assess the quality of delivery and take adaptive decisions when necessary.

Extensive testing was performed in order to tune the QoDGS and to determine values for w_A and w_B , and for w_I , w_2 , w_3 and w_4 that best achieve the QOAS's goals in local broadband IP-networks. Good adaptiveness and responsiveness to network traffic variations, significant quality stability, high link utilisation and good end-user quality were obtained for the following set of weights: $w_A = 0.75$, $w_B = 0.25$, $w_I = 0.4$, $w_2 = 0.3$, $w_3 = 0.2$, and $w_4 = 0.1$. Tests that involve a simulation model of a QOAS-based multimedia streaming system and their results are presented in the sixth chapter that focuses on experimental testing.

4.6 Server Arbitration Scheme (SAS)

4.6.1 SAS Overview

Apart from QoDGS, another important component of the QOAS is the Server Arbitration Scheme (SAS) that has to analyse the feedback-reported information and to take adaptive decisions if and when necessary. Its goal is:

- to collect the feedback transmitted QoD scores computed by the QoDGS,
- to analyse the QoD scores received during a recent period of time,
- to take decisions in relation to the reported quality of delivery and to trigger quality adaptations.

By determining quality adaptations based on feedback-received QoD scores, the SAS aims at improving the quality of delivery in the existing streaming conditions.

4.6.2 SAS Principles

In order to achieve the SAS's goals, the following principles related to the SAS design were formulated.

- 1. The SAS takes into account the QoD scores as received via feedback from the QoDGS located at the client. SAS should differently consider the positive feedback reports and negative ones in relation to the current quality of the streamed multimedia clip. An asymmetric behaviour should ensure a fast reaction during difficult delivery conditions that affect the end-user quality, reducing their length and a slow reaction to feedback that indicates improved streaming. In this manner the SAS helps in the elimination of the cause of the increased traffic condition, fast reducing its contribution to the overall transferred data. By cautiously reacting to positive reports, SAS intends to allow for the network to recover before upgrading the quality of the streamed multimedia and therefore to increase its contribution to the overall traffic.
- The SAS takes into account more than a single feedback report in order to reduce the influence of eventual noise in the received QoD scores that may cause temporal instability in selecting a quality for the streamed multimedia.

- 3. The SAS is allowed to suggest only quality changes adjacent to the current quality of the streamed multimedia clip. This is in order to reduce the eventual negative influence in the end-user perceived quality.
- 4. The SAS has to be able to help in the quality adaptation process even if the feedback reports do not arrive at the server. This is considered as an indication of network congestions and determines SAS to suggest quality degradations, trying to help in solving this problem. If the situation continues, the stream is transmitted at the lowest possible quality.
- 5. The SAS analysis and decision taking has to be performed **very fast** in order not to influence negatively the performance of the multimedia server application. Also it has to be able to suggest quality variations **at any time**, independent from multimedia streaming and from the process that effectively performs the quality variation at media level. In this way the latter can be performed in such a manner that is the least disturbing for the remote viewer.
- 6. The SAS decisions have to be dependent only on QoD scores received from the QoDGS and the arrival or not of the feedback messages. They must not be dependent on other parameters in order to have a deterministic behaviour of the QOAS and QOAS-based multimedia streaming system that uses it.

4.6.3 SAS Design

The server-situated SAS whose block structure is presented in Figure 4-16 was designed according to the principles previously mentioned in section 4.6.2. The figure also shows the input parameter taken into consideration by the SAS - the QoD scores computed by the QoDGS and sent to the server via feedback. Due to the fact that SAS asymmetrically assesses the feedback, SAS consists of two similar modules: the *Downgrade Module* – in charge with the analysis of feedback and assessing the opportunity of a downgrade in the stream quality and the *Upgrade Module* – which analyses if an upgrade in stream quality is beneficial. These modules suggest changes in quality to the *SAS Decisions Module* that takes the decisions. The *Timer*'s role is to driven degradation decisions if the feedback does not arrive at the server, suggesting that there is a delivery problem.

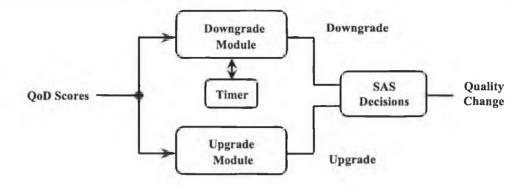


Figure 4-16 SAS block-level structure

The *Downgrade* and the *Upgrade* modules are similar, the difference being the time scale on which the assessment of the necessity to suggest quality adjustments is performed. They consist of a circular buffer in which the QoD scores are stored. A sliding window that encompasses the most recently received scores provides the data source for the analysis. The average value of these received QoD scores is compared with the current server quality state that determines the quality of the streamed multimedia clip (see section 4.3). This comparison allows the *Decisions Module* to take or not to take the upgrade and respectively the downgrade decision, as suggested by the other modules. This affects the quality of the streamed multimedia clip, increasing or decreasing it.

The *Timer* module allows a period equal to four times the round trip time measured during the session for the arrival of the expected feedback. If this does not happen, automatically it suggests degradations decisions to be taken.

4.7 Data Transmission and Feedback Mechanism

The QOAS architecture presented in section 4.2 includes Communication Managers in charge with client-server communication establishing, controlling and ending. The QOAS makes use of a double-channel for exchanging both multimedia data and control information. A bidirectional link is meant for transmission of control messages in charge with session control. This link is also used for sending feedback messages that carry the client-computed QoD scores to the server. An unidirectional link from the server to the client transports multimedia data to the latter where is decoded and played. Figure 4-17 shows schematically this double-link communication channel between the QOAS server and the QOAS client.



Figure 4-17 Multimedia data transmission and control data exchange between QOAS server and client applications

For selecting protocols for multimedia data transport and session control, the IETF Multiparty MUltimedia SessIon Control (MMUSIC) Working Group's⁵⁷ documents were consulted, the IETF Audio/Video Transport (avt) Working Group's⁵⁸ works and the ITU-T publications⁵⁹. The Real-time Transport Protocol (RTP) [100] and its companion Realtime Transport Control Protocol (RTCP) [100] were selected for transporting data and the Real Time Streaming Protocol (RTSP) [167] for controlling data delivery session, respectively. Next these protocols' characteristics are indicated and the choice for them justified.

The Real-time Transport Protocol (RTP) is both an IETF Proposed Standard - RFC 1889 [100] and an International Telecommunications Union (ITU) Standard - H.225.0 [237] and currently seems to be "the standard" for transporting time-sensitive data. It is a higher-level transport protocol, which provides transport functions for applications that involve transmissions of data with real-time or interactive characteristics (e.g. audio, video, simulation data, etc.), over unicast or multicast networks. RTP provides support for payload type identification, sequence numbering, time stamping and delivery monitoring. The data transport protocol is complemented by a control protocol, the Realtime Transport Control Protocol (RTCP), which allows for monitoring of the data delivery and provides support for minimal control and identification functionality [238]. RTP and RTCP are designed to be independent from the underlying transport and network layers, although UDP/IP [14, 165] is preferred. RTP does not provide resource reservation and does not guarantee any quality of service for real-time services. It is the most used protocol for time sensitive

⁵⁷ IETF Multiparty MUltimedia SessIon Control Working Group, http://www.ietf.org/html.charters/mmusic-charter.html

⁵⁸ IETF Audio/Video Transport Working Group, http://www.ietf.org/html.charters/avt-charter.html

⁵⁹ International Telecommunication Union - Telecommunication Standardisation Sector (ITU-T), http://www.itu.int/publications/main_publ/itut.html

data, including multimedia. Many commercial companies make also use of it for delivering data with their products (e.g Microsoft's NetMeeting⁶⁰, Apple's QuickTime⁶¹).

The Real Time Streaming Protocol (RTSP), a IETF proposed standard - RFC 2326 [167], is a control protocol for initiating and directing the delivery of streamed multimedia, acting like a "network remote control" protocol. It does not typically deliver the continuous streams itself, although interleaving of the continuous media stream with the control stream is possible. RTSP provides the following specific benefits to its users: enables full bidirectional stream control, offers high reliability over current infrastructure, ensures low overhead data delivery, fully exploits emerging technologies and protocols (e.g. IP Multicast, RTP, etc.), offers support for security and intellectual property rights protection. Very important is also that it is scalable, working well both for large audiences as well as single-viewer media-on-demand. Although other proposed standards like SIP [168] or H.323 [239] could also be used, their high complexity in comparison to RTSP and mainly the general tendency of their applicability in audio/voice transmissions [60], made us to preferred the latter. This is especially since important commercial companies such as Progressive Networks³⁰ use RTSP for controlling streaming sessions.

QOAS uses RTP for transporting multimedia data and therefore it uses RTP packet format. The most important fields for the OQAS operation are: the "Sequence number" that allows the receiver to detect eventual packet loss and to restore packet sequence in case of out-of-order arrival of packets and the "Timestamp" which permits the computation of one-way delays and jitter delays during streaming.

RTCP is used to both transmit feedback from the clients to the server and to transmit adaptation-related information from the server to the client. Since the standard allows for the definition of new packet types (with the reservation of the definition of the associated packet types with the Internet Assigned Numbers Authority (IANA)⁶²), a new RTCP packet type was defined. This packet respects the RTCP packet structure, but it is shorter due to the size of the information it carries.

⁶⁰ Microsoft's NetMeeting, http://www.microsoft.com/windows/netmeeting

⁶¹ Apple's QuickTime, http://www.apple.com/quicktime/tools_tips/tutorials/rtp.html

⁶² Internet Assigned Numbers Authority (IANA), http://www.iana.org

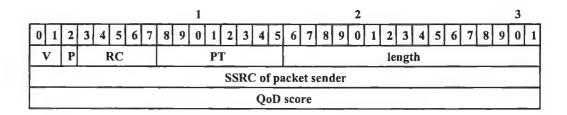


Figure 4-18 RTCP addition - QOAS receiver report packet type

Figure 4-18 presents the structure of the proposed QOAS Receiver Report packet type (QOAS-RR) that has all the fields common for all RTCP packets:

- version (V): 2 bits which identifies the version of RTCP (same as in RTP data packets),
- padding (P): 1 bit that indicates, if the padding bit is set, that the RTCP packet contains some additional padding octets at the end which are not part of the control information. The last octet of the padding is a count of how many padding octets should be ignored. Some encryption algorithms with fixed block sizes may need padding. In a compound RTCP packet, padding should only be required on the last individual packet because the packet is encrypted as a whole.
- reception report count (RC): 5 bits which contains the number of reception report blocks contained in this packet. Since QOAS has none, a valid value of zero is set.
- packet type (PT): 8 bits which contains a constant that identifies the packet type. Since 200-204 are used by RTCP, higher values can be used by QOAS.
- *length*: 16 bits which shows the length of this RTCP packet in 32-bit words minus one, including the header and any padding. (The offset of one makes zero a valid length and avoids a possible infinite loop in scanning a compound RTCP packet, while counting 32-bit words avoids a validity check for a multiple of 4.)
- SSRC: 32 bits The synchronisation source identifier for the originator of this SR packet.

Apart from these, a 32 bit field *QoD score* that stores the QOAS quality of delivery grading score as computed by the client-located QoDGS is part of the QOAS-RR packet structure.

RTSP is used for session establishment, control and disconnect. The most important RTSP methods, used also by the OOAS are: SETUP, PLAY, PAUSE and TEARDOWN.

- SETUP: Causes the server to allocate resources for a RTSP streaming session and starts it.
- PLAY: Requests streaming of a stream and starts data transmission (performed regularly using RTP/RTCP).
- PAUSE: Temporarily halts the streaming process without freeing the allocated resources.
- TEARDOWN: Frees resources associated with this RTSP streaming session.

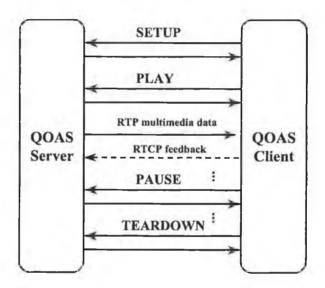


Figure 4-19 Example of a RTSP session

Figure 4-19 presents a possible RTSP session using QOAS approach that consists of a SETUP and a TEARDOWN method invocation, at least one PLAY method call that starts the streaming process and an indefinite number of pair calls for the PLAY and PAUSE methods. The data transmission is performed using RTP as a transport protocol and feedback is sent via RTCP.

4.8 InteR-stream QOAS

The InteR-stream QOAS (IR-QOAS) is an extension of the IntrA-Stream QOAS (IA-QOAS)-based adaptation and aims for a finer adjustment in the overall adaptation process to yield

better utilisation of network resources. The IR-QOAS is also responsible for preventing the IA-QOAS-driven adaptations of the multimedia stream transmissions from reacting simultaneously to variations in the delivery network. Such an eventual synchronisation may trigger the IA-QOAS's over-reaction and determine both under-usage of the available bandwidth and reduced perceived quality for the remote viewers. The IR-QOAS is meant to work in conjunction with the end-to-end IA-QOAS aiming at achieving both high end-user perceived quality and high utilisation of the shared network resources.

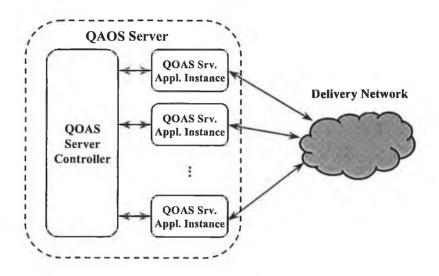


Figure 4-20 QOAS Server Controller in permanent contact with QOAS server application instances in charge with the deployment of the inter-stream QOAS

Figure 4-20 presents the localisation of the IR-QOAS, which is deployed at the Server Controller Application level. The QOAS Server Controller Application (SCA) is in permanent contact with all QOAS Server Application instances part of the same QOAS Server, communicating with them. They exchange information in order to allow for the SCA to have an overall view of the multimedia delivery process and to actively contribute in the adaptive process as a whole. At the QOAS Server Application level, the IR-QOAS-driven adaptation involves the Server Arbiter whose IA-QOAS-based adaptive decisions (presented in section 4.6) are being filtered before they are taken. The filtering is performed at the SAS Decisions block level that was presented in Figure 4-16.

IR-QOAS requires the definition of a control state, named SCA state, on whose value depends the overall adaptive decisions suggested to be taken by the IA-QOAS-s. Based on the history of all the IA-QOAS multimedia streaming processes in progress, IR-QOAS estimates the network transmission conditions and sets its SCA state. If all the IA-QOAS-based remote

multimedia deliveries are performed at maximum quality, the SCA state is set to "normal" suggesting that the network traffic conditions are good for streaming. If some IA-QOAS-based streaming processes have adjusted downwards their transmission quality (and consequently have decreased their IA-QOAS server state), this suggests that there are some delivery problems and the SCA state is set to "difficult". During the overall streaming of multimedia streams SCA state bounces between "normal" and "difficult" affecting the IR-QOAS adaptive decision.

IR-QOAS interferes with the IA-QOAS individual adaptive streaming processes in three occasions: during the initialisation stage and after ending the streaming, as well as when any IA-QOAS-suggested adaptive measures are taken.

The initialisation stage for IA-QOAS-based streaming is very important since the initial transmission quality (and in consequence the corresponding transmission rate) of a multimedia stream requested to be remotely delivered should not be different than the possibility of the network to deliver. If it is higher, during the transitory period in which an IA-QOAS-driven quality decrease is performed the end-user quality may be severely affected by loss for example, not only for the current streaming process by also for others. It also may trigger adaptive over-reaction from the other concurrent streams, reducing also the network utilisation. If the initial quality of the streamed multimedia takes a lower share than the network potential available bandwidth, the current remote viewer perceives a lower quality than he/she could see and the network utilisation is not as high as it could be.

During the initialisation stage, based on its SCA state, the IR-QOAS suggests to the individual IA-QOAS streaming processes a starting quality for their multimedia clips. In the "normal" SCA state any newly requested stream is going to be delivered at the highest quality. In the "difficult" state, however, the new stream's initial quality and in consequence its corresponding rate is computed averaging the rates the other streams are currently being delivered with, as in equation (4-16).

$$InitRate = \frac{\sum_{i=1}^{N} CrtRate_{i}}{N}$$
 (4-16)

In equation (4-16) N is the total number of already existing concurrent streams and $CrtRate_i$ is the rate the stream i is being streamed with.

In order to prevent any loss from occurring, the IR-QOAS forces a quality reduction on some of the IA-QOAS streaming processes that are performed at a higher quality than the average. A similar "imposed" adaptation process, but positive for the end-viewers, is performed when any streaming process has ended. In this case some of the IA-QOAS streamed multimedia clips that are delivered at a lower quality will benefit from a quality increase. The number of streams affected by this IR-QOAS forced adaptation is determined from the estimation of the available bandwidth, the number of the on-going multimedia deliveries and their quality.

$$NoStreams = \left(\frac{CrtBwd + InitRate - EstimBwd}{AvgRateDif}\right)$$
(4-17)

$$CrtBwd = \sum_{i=1}^{N} CrtRate_{i}$$
 (4-18)

Equation (4-17) presents the formula used for the determination of the number of streams that are affected by this "forced" adaptation. In this formula the current bandwidth at the moment the "forced" adaptation is required - CrtBwd is computed as in formula (4-18) from the current rates of the total number of the existing parallel streams - N. InitRate is the suggested initial rate by the IR-QOAS as computed in equation (4-16), the EstimBwd is the estimated bandwidth of the connection and the AvgRateDif is the average rate difference between the different quality levels defined for all the multimedia streams. The result is more accurate if this difference is the same for all streams taken into account. The total bandwidth estimation EstimBwd is obtained by averaging an under-estimation UEstimBwd which saved the current bandwidth CrtBwd, computed as in equation (5-19) last time SCA state was "normal" and a supra-estimation SEstimBwd computed similarly when the SCA state was set to "difficult".

$$EstimBwd = \left(\frac{UEstimBwd + SEstimBwd}{2}\right)$$
 (4-19)

During the adaptive streaming involving IA-QOAS, IR-QOAS selectively permits some of the multimedia streaming processes to react to the received feedback, in a step-by-step process, aiming for achieving near maximum link utilisation and long-term fairness between the clients.

In order to reduce the eventual synchronisation between the IA-QOAS-based streaming processes, the IR-QOAS has introduced a mechanism that spreads their reaction over a period of

time, introducing random delays in their adjustment decision taking process. This is performed with the hope that if some of the IA-QOAS-based adaptive processes that have noticed problems in the delivery network decrease their contribution to the traffic by adjusting the quality of their streamed multimedia, the traffic problem will be solved and the other will not have to adjust anymore. This behaviour also depends on the IR-QOAS associated SCA state.

If the SCA state is "normal" any downward adaptation suggested by individual IA-QOAS is performed without interference from the IR-QOAS. Once the adaptation is performed this affects the SCA state that changes to "difficult" and the IR-QOAS reacts differently to quality adjustments. A separate timer with a random timeout period is associated to each request for quality degradation issued by individual IA-QOAS schemes, involving their SAS modules. When any of these timeout periods expires, the IR-QOAS allows the SAS's *Decisions* module presented in Figure 4-16, to perform the suggested degradation. If the QoD scores received via feedback by the IA-QOAS do not indicate an improvement in the quality of delivery when the first adaptive measures were taken, the IA-QOAS downgrading in the quality of the streamed multimedia will continue when next timeout periods expires. However if the delivery situation improves and the IA-QOAS-s that have requested downwards adjustments in their streamed quality stop their requests, the timers are reset and the decrease in quality will not take place. A similar process occurs when the IA-QOAS-based streaming processes request increases in their streamed multimedia quality.

This IR-QOAS driven adaptation ensures not only increased quality of the end-user quality mainly during the IA-QOAS initial stage, but also higher available network resource utilisation and, very important, higher stability of the QOAS adaptive process in terms of quality variation that may affect the end-user perceived quality.

4.9 Applicability Considerations

The QOAS relies on feedback in order to learn about the quality of the streaming process and to take the necessary adjustment decisions. The existing research like [197, 6, 7] that took into consideration feedback for performing adaptations shows that the faster the feedback messages arrive at the server, the better the results of the adaptation process are. If the feedback takes too long a time to arrive, the information the server has about the system does not reflect the current reality anymore and the scheme may react too late to make the difference or out of synch. For example the feedback-controlled scheme may not react in time to prevent losses from occurring once increased delays have been reported or if the loss is already affecting the streamed multimedia. Also there is

the possibility that the adjustment measures to be taken when the cause of PUDL may have already passed, decreasing the quality of the streamed multimedia when it is not required anymore. Therefore the QOAS is best recommended to be applied in local or metropolitan area networks, local cable delivery networks, or local all-IP broadband networks where **fast feedback** is feasible. Experimental test results presented in the sixth chapter show how the performances of a QOAS-based multimedia streaming system are influenced by network latencies.

The applicability of any adaptive scheme, including QOAS is most recommended in **networks with a potential for congestion**. This is because this scheme offers significant benefits in comparison to a non-adaptive approach only if shared bandwidth is limited. The benefits are also significant in networks with highly increased traffic conditions, even if compared with other feedback-based adaptive schemes like for example TFRCP [6] and LDA+ [7]. The results of tests that have studied this comparison are also presented in the sixth chapter.

Very important is that the multimedia streams' viewers targeted by the QOAS can tolerate a certain degree of quality variation. In consequence QOAS does not target multimedia systems whose viewing quality has life-threatening or research-quality consequences as for example some areas of Medicine (e.g. Surgery), Physics (e.g. atomic phenomena) or Transport (e.g. Radar systems). QOAS can be successfully applied in the entertainment industry, business for video-on-demand applications, commercial presentations, video-conferencing in which a slight decrease in the quality is not disturbing and is even preferred to interruptions in the play-out for buffering performed by many existing solutions, as reported in [240].

The QOAS usage was considered in the absence of any error-concealment techniques' [241] deployment that could improve the end-user perceived quality of a streamed multimedia clip affected by loss during transmission. In principle any error concealment technique could be taken into account in conjunction with QOAS to further improve the end-user perceived quality in the tested conditions. However, further tests have to be performed to see the benefit of using QOAS in conjunction with such error-control mechanisms if the multimedia transmissions are subjected to higher loss rates.

4.10 Summary

The fourth chapter focuses on the detailed presentation of the Quality-Oriented Adaptation Scheme (QOAS) for multimedia streaming. It starts with a general description of the scheme and of the architecture of the QOAS-based multimedia streaming system that implements it. The chapter then continues with the presentation of the QOAS's main mechanism: the IntrA-stream QOAS (IA-QOAS). Detailed information about the IA-QOAS's main components and their functionality are given in separate sections of this chapter: the client-located Quality of Delivery Grading Scheme (QoDGS), the Server Arbitration Scheme (SAS) and the Data Transmission and Feedback mechanisms. While multimedia data is being streamed via the Data Transmission mechanism, QoDGS monitors and assesses both long term and short-term variation of some transmission parameters and of the end-user quality. The QoDGS also regularly grades the quality of the ongoing streaming process in terms of QoD scores in a three-stage process presented in detail. These scores are sent using the Feedback mechanism to the server whose SAS processes them. The SAS takes into consideration the values of a number of recent feedback reports, analyses them and suggests adjustment decisions to be taken by the server. Detailed information is also offered about the parameters taken into account by the QoDGS in its quality assessment process and about the metric Q used to estimate the end-user perceived quality during multimedia streaming. The InteR-stream QOAS (IR-QOAS) mechanism used to complement IA-QOAS in order to achieve better end-user perceived quality and higher network utilisation when streaming multimedia was also described in this chapter. At the end, QOAS applicability considerations are presented, indicating both the recommendations and the limitations for the QOAS potential deployment.

Chapter V Implementation Details

Abstract

Since the proposed Quality Oriented Adaptation Scheme (QOAS) needs extensive testing, both simulations and emulations are employed in order to produce real-life like network delivery conditions. In these conditions, both a simulation model system and a real prototype system that instantiate QOAS have been implemented and tested. This chapter presents details about both implementations of the proposed QOAS, the simulation model and the real prototype system.

5.1 Implementation of the Simulation Model System

5.1.1 Network Simulator version 2

The simulations are performed using the Network Simulator version 2 (NS-2) [246], which is an object-oriented discrete event simulator, written in C++, with an OTcl [247] interpreter as front-end. In NS-2, the simulations are performed according to simulation scenarios that consist of several components [242]. The most important are: *a network topology*, that specifies the physical inter-connections between nodes and the characteristics of links and nodes, *traffic models* which define the senders and the protocol(s) of the packet transmission and *test scenarios* which generate traffic causing network dynamics designed to test a certain implementation.

The NS-2 simulator supports two class hierarchies: the compiled hierarchy, consisting of C++ classes and the interpreted hierarchy written in OTcl. Extensions to the first hierarchy are done through C++ classes if changes in the manner the exchanged packets are processed are required and the behaviours provided by the existing C++ classes are not enough to solve these problems. The second set of classes is appended with scripts written for configuration, setup and single-use modifications of the overall NS-2-provided behaviour. In general the latter manipulate existing or newly built C++ objects.

5.1.2 Simulation Model's Implementation Overview

For fully testing QOAS NS-2 provides both network topology and components and different traffic models for building various test scenarios. In consequence the implementation involves only building the QOAS client-server model system that follows the architecture presented in the fourth chapter.

Apart from QOAS client and server applications, some other mechanisms had to be implemented in order to allow for extensive QOAS testing. Among these mechanisms are the RTP transport of multimedia data packets model, the enhancement of the drop-tail router queue model and the QOAS server controller application model. The implementation of these models is presented next.

5.1.2.1 RTP-based Transport of Multimedia Data Packets

In order to allow for the RTP-like handling of the multimedia data packets, the UDP-related classes are extended or used, both in C++ and in OTcl, as recommended in [249]. In this context a MultimediaHeaderClass was defined in C++ and was associated with the OTcl hierarchy name "PacketHeader/Multimedia". Also a RTP agent class named UdpMmAgentClass that inherits the TclClass base class was associated with the OTcl hierarchy name "Agent/UDP/UDPmm" and with the C++ class UdpMmAgent that was implemented as an extension of the UdpAgent class. The most important methods provided are the sendmsg() that sends a number of bytes received from the application level to the UDP level, after attaching the RTP header and the recv() which is automatically called by the underlying UDP agent when a packet is received in order for the RTP header to be removed and the data to be sent to the application level.

5.1.2.2 Drop-Tail Router Queue

Although a drop-tail queue is defined by the NS-2, since during simulations extensive statistical information is needed for fully assessing the QOAS's performances, an enhanced drop-tail queue was implemented. It performs statistical-multiplexing of incoming data, drops packets if they exceed the storing capacity and records statistics.

The StMuxSingleQueueClass was defined in C++ inheriting the TclClass and was associated with both the OTcl hierarchy name "Queue/StMuxSingleQueue" and with the C++ class StMuxSingleQueue that was implemented as an extension of the Queue class. The most important

methods are *enque()* which enqueues an incoming packet if there is storage space left or drops it otherwise and *deque()* which retrieves the packets in a FIFO manner. Both methods update also the statistics information related to each data flow by a call to *update_statistics()*. This makes use of a specially designed complex list of statistic-related structures that was implemented by the *PacketQueueList* class that extends the C++ *TclObject* class. Regularly the statistics are written to a log file by calling the *write_statistics()* function. This is performed by a specially defined timeout timer: *StMuxStatisticsTimer*.

5.1.2.3 QOAS Server Controller Application

The QOAS server controller application implements the Inter-stream QOAS as it was described in section 4.8 of this thesis.

The defined AdSrvCtrlClass class inherits the TclClass and is associated to the OTcl hierarchy name "Application/AdSrvCtrl" and to the C++ class AdSrvCtrl that was implemented which inherits the Application class. The most important methods of the AdSrvCtrl class and their roles are presented next.

The function *attachApp()* adds the indicated QOAS server application to a specially built list of applications registered with the server controller application. Only these applications will be affected and affect the functionality of this controller and only when they are active. The activation of a registered application is performed when a new streaming process starts by a call to the *activateApp()* method and ends when the streaming process has ended by a call to the *deactivateApp()* function.

Function computeStartRate() computes the start rate for a new QOAS multimedia stream in the presence of similar other streams in existing delivery conditions as their average streaming rates, determined using the computeAverageRate() method. The computeStartRate() function also applies the "imposed" rate adaptation for the streams with the highest delivery rate in order to accommodate for the new stream by calls to the decreaseStreamStates() function.

The function *computeEndState()* triggers "imposed" rate adaptations to the existing QOAS streams after a multimedia stream generated by a registered application has ended. This involves an implicit increase in the transmission rate for the streaming processes with the lowest rate performed by the function *increaseStreamStates()*.

The functionality of both *computeStartRate()* and *computeEndState()* methods relies on the bottleneck link bandwidth estimation performed using the function *computeTotalBandwidth()*.

5.1.3 Implementation of the QOAS Server Application Model

The QOAS server application model relies on the definition of the *AdSrvAppClass*, an extension of the TclClass base class. This class makes the association between the OTcl hierarchy name of "*Application/AdSrvApp*" and the C++ *AdSrvApp* implemented class that inherits the *Application* class provided by the NS-2. The implementation of the latter follows the server application architecture presented in the fourth chapter and is described next in terms of the most important *AdSrvApp* class's member functions and their roles. This description is structured based on the server application's architectural blocks.

5.1.3.1 Multimedia Acquirer, MPEG Encoder and Multimedia Database

Both multimedia capturing and MPEG encoding are performed offline using a Canopus Amber MPEG hardware encoder card. For each multimedia content, different quality stream versions were encoded from various clips at five bit-rates equally distributed between 2 and 4 Mb/s. The resulting MPEG files are then parsed using a specially built application (named Read_IPB_Frames) that saves trace files in the following format: frame number, frame type (I, P or B), display time (ms) and frame size (bytes). These traces are used as input by the QOAS server when adaptively streaming multimedia data, acting like a multimedia database. The class that accesses this database and adaptively reads the frame-related information was named AdaptiveTraceFile and inherits the NS-2 NsObject class. Among its most important methods is setup() that associates indexes to each of the different quality version files according to their corresponding QOAS server quality states allowing for parsing of the correct quality file during adaptive streaming. Similarly important is the get_next() method which retrieves the information related to the next frame to be streamed given the existing QOAS server quality state.

5.1.3.2 Server Communication Manager and Transmission Shaper

AdSrvApp class through its command() method defines control functions that allow for the parameterised setup of the QOAS server application model via a OTcl script. Among these functions are "attach-agent" and "attach-agent2" that associate transport layer agents, which are in charge with the data transmission. In this implementation an RTP agent was already defined an is

used for this purpose, but since the implementation is flexible, it allows for other agents to be also used if desired. Another important OTcl method is "attach-tracefile" that allows for associating a trace file to a certain movie name and server quality state.

Other important methods of the AdSrvApp class are attachCtrl() that links the QOAS server application to the QOAS server controller application, initialise() that initialises the server application related structures and decide_start_rate() that determines the adaptive streaming starting rate based on the controller's suggestion.

The AdSrvApp class also implements the RTSP-based session control mechanism and processes the feedback messages using a set of methods presented next. recv_msg() is the function called automatically when the underlying transport level receives any packet for this application, and according to its type, a different method is called. The RTSP server-side SETUP is performed by the process_connect() method, PLAY – by the process_request() member function and SHUTDOWN – by the process_shutdown(). These are followed by server application answers by calls to the send_ackconnect_pkt(), send_ackrequest_pkt() and send_ackshutdown_pkt(). The feedback messages received by the recv_msg() are processed by the process_feedback() method.

The process_connect() function initialises the resources necessary for adaptive streaming, whereas the process_request() selects the requested stream and starts streaming by calling start_sending_mmdata(). The latter starts a timer mechanism implemented by SrvMmdataTimer class. This class, which implements the Transmission Shaper, is in charge with performing the timeout-driven streaming process at the rates associated with the different quality streams whose tracefiles were registered with the multimedia database. The actual sending of data packets is performed by the send_mmdata_pkt() function that uses the underlying transport agent's capabilities in order to do this. The process_shutdown() function releases all the resources used by the application.

5.1.3.3 Feedback Manager and Server Core

The Feedback Manager and the Server Core's Server Arbitration Scheme (SAS) work in conjunction in order to receive and process the incoming feedback and to take adaptive adjustments if necessary. As previously mentioned the feedback is received by the *recv_msg()* and is processed by the *process_feedback()* method. The latter implements the SAS mechanism which was described in detail in section 4.6 in conjunction with the *set_scale()* method. The actual state change is performed only at the beginning of a GOP after calling the *setTxState()* function.

In case that for some time no feedback is received, the QOAS server application estimates that there is a delivery-related problem and decreases its quality state and therefore the multimedia transmission rate. This mechanism is implemented by the *SrvTimeoutTimer* class and the *AdSrvApp*'s method *reduce pkt txrate()*.

5.1.4 Implementation of the QOAS Client Application Model

The QOAS client application model is implemented by the AdSrvAppClass, an extension of the TclClass base class. This class makes the association between the given OTcl name of "Application/AdCliApp" and the C++ class AdCliApp that was implemented such as it inherits the NS-2 Application class. Our implementation of the latter follows the client application architecture presented in the fourth chapter and is described next based on its block-structure in terms of the most important class methods and their roles.

5.1.4.1 MPEG Decoder and Multimedia Player

Since the simulation model does not play out the multimedia data received, there is no need for the decoding process. However since the data received has to be consumed to prevent overflowing the receiver buffer, the *AdCliApp*'s method *play_data()* does this as it plays the multimedia data with the associated display frequency. This reading frequency is controlled by an object that instantiates the specially defined *CliPlayoutTimer* class.

5.1.4.2 Client Communication Manager

Similar to the AdSrvApp class, the AdCliApp defines in its command() method control functions that allow for setting up of the QOAS client application model via a OTcl script. The "attach-agent" and "attach-agent2" functions associate transport layer agents in charge with the data transmission to this application. In this implementation the RTP agent already presented is used. A third function provided is "attach-recv-buffer" that associates a certain receiver buffer implementation to the client Communication Manager. In the implemented solution this receiver buffer was defined by the RecvBuffQueueClass that inherits the base class TclClass and associates the OTcl name "RecvBuffer" with the C++ class RecvBuffQueue. The later extends the TclObject C++ class and provides means of storing the packets that have arrived at the client using a FIFO policy. Its main methods are enque() for storing and deque() for retrieving data.

The AdCliApp class implements the RTSP client-based session control mechanism using a set of methods presented next. The RTSP client-side SETUP is performed by the start() method that calls the send_connect_pkt() method. If the server's answer is positive, the client sends the PLAY command by calling the send_request_pkt() member function. When the client desires the end of the streaming process SHUTDOWN command is sent to the server by calling the send_shutdown_pkt() method. These commands are sent to the underlying transport agent which does the actual sending of data by a call to the send_control_pkt() method. The server application answers to control messages by sending corresponding ACKs. These messages are processed by the AdCliApp application in its recv_msg() function by calling recv_ackconnect_pkt(), recv_ackrequest_pkt() and respectively recv_ackshutdown_pkt(). The recv_msg(), method called automatically when the underlying transport level receives any packet for this application, receives also all the multimedia data packets that are stored by the recv_mmdata_pkt() in the client receiver buffer, an instance of the RecvBuffQueue class.

5.1.4.3 Feedback Indication Unit and Client Core

The Feedback Indication Unit and the Client Core co-operate in order to support the functionality of the Quality of Delivery Grading Scheme (QoDGS) whose principle was presented in detail in section 4.5. Presented briefly, the QoDGS's goal is to monitor and to grade the network-related parameters' values and variations as well as the estimated end-user perceived quality during multimedia streaming. In order to do this, QoDGS's implementation makes use of the specially built structures qo_transmission and qot_XXX, where XXX stands for the monitored parameter and could be delay, loss, jitter and percequal (i.e. end-user perceived quality). These structure initialisation is performed by the init_qot_XXX() methods, the QoDGS-related information update by the adjust_qot_XXX() functions and the partial parameters' grading by the grade_qot_XXX() methods. The final computation of the QoD_{Score}-s is performed by the grade_tx() function, member of the AdCliApp class.

The initialisation of these feedback-related structures is done in the <code>init_variables()</code> method, the adjustment of parameters' values and variations in the <code>recv_mmdata_pkt()</code>, immediately after a new data packet was received at the client and the QoDGS final grading every time a feedback message is sent to the server. The frequency of the feedback messages is controlled by the timer class <code>CliFeedbackTimer</code> whose timeout is set via the OTcl script, allowing for high testing flexibility. Every time when the timeout occurs, a feedback <code>packet</code> is sent using <code>send_feedback_pkt()</code> to the RTP transport agent that actually does the actual transmission.

5.2 Implementation of the Real Prototype System

5.2.1 Prototype System's Implementation Overview

The prototype system built in a Windows environment using Microsoft Visual C++ 6.0 follows the block-level architecture presented in the fourth chapter. The implemented system consists of two applications: a server and a client, which inter-communicate via a network. Both the design of the system and its implementation follow an object-oriented approach and make use of the Microsoft Foundation Class (MFC) as the base class structure for the creation of the majority of the implemented classes. The Windows event and messaging systems support the message and event handling in both the client and the server applications. The implementation also makes use of the threading support offered by the Win32 multi-tasking environment, of the sockets mechanism provided by the Windows Sockets 2 (WinSock2) architecture for applications inter-communication and of the Microsoft's Open Database Connectivity (ODBC) API for accessing the Microsoft Access database used.

Before implementation details related to each of the system components, the server application and respectively the client application are given, next information related to implementation issues common for both of them are presented.

5.2.1.1 Applications' Inter-communication

For the implementation of the two communication channels, one for bi-directionally exchanged control messages, including feedback, and the second for unidirectional transport of data packets from the server to the client application, WinSock2 sockets mechanism and MFC library were used. The MFC's *CSocket* class is the base for the implementation of all the communication-related classes built, which inherit from it the basic socket functionality.

Figure 5-1 presents the implemented classes involved in the process of establishment, control and disconnection of the double client-server communication link and the two container classes *CMySrvDoc* and *CMyCliDoc*, located at the server and respectively at the client. The figure also schematically describes the process of double-channel creation for a requesting client and involves two steps. The first step consists of the establishment of the control link, whereas the second step involves the creation of the data link. During disconnection first the destruction of the data link is performed and then the control link is terminated.

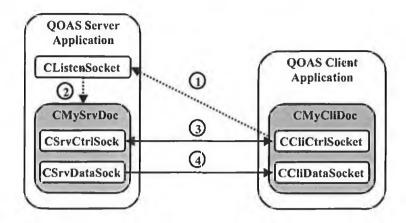


Figure 5-1 QOAS Client-server inter-application communication

At the server a listening socket, which instantiates the *CListenSocket* class, described by the publicly known pair server IP address - port number, allows for clients to request services. If a particular request is accepted, the exchanged control messages allow for the double communication link to be established. First the control link is established between instances of the client and server Control socket classes (i.e. CSrvCtrlSock and CCliCtrlSock) and then the data link is created between instances of their Data socket classes (i.e. CSrvDataSock and CCliDataSock). The Data classes always implement UDP which ensures fast, although unreliable packet transmission between the sender and the receiver. The Control classes can implement both TCP and UDP allowing for a choice when sending control messages.

Once the double-channel link has been successfully established, data can be sent across the network from one communication partner to the other. The Windows messaging and event systems permit a very simple implementation of the receiver-related functionality for all the socket-based classes. When a data packet is incoming, the application automatically calls the *OnReceive()* method that processes the data.

5.2.1.2 Data Buffering and Statistical Data Collection

Buffering has a very important role in this prototype system, not only because it is involved in the transmission, decoding and playing of multimedia content, but also because it provides important information to the client's Quality of Delivery Grading Scheme (QoDGS). QoDGS makes use of information related to network-related parameters' values and variations. In this context the Circular Buffer structure was especially designed for a double role: data buffering and statistical data collector.

The Circular Buffer is composed of a number of equal size buffers linked in a circular manner in order to allow for their usage in a way that simulates a pool of buffers. These buffers are allocated once during the initialisation phase and re-used during streaming for storing data. This provides a significant advantage since no CPU processing time or I/O activities are required during the streaming either for repeated memory allocations and de-allocations or for the structure management.

The Circular Buffer allows for the existence of a number of ordered lists of buffers within its structure. This allows both for re-ordering if necessary, thus restoring the original sequence of data packets and for storing data during different processing stages with little effort. For example the list of buffers with encoded data, the list of decoded data buffers, the list of buffers containing data being played out and the list of empty buffers are all stored in the same structure. This is possible because of the pipeline-like processing of the data buffers and of the circular structure of the buffering system. When the last processing stage (i.e. display) was completed the buffer will be returned for re-use as the last buffer in the list of free buffers. The buffer memberships to different lists are indicated using marker bits, without modifying the buffer position in the circular buffer system. These lists of buffers are accessed via a set of pointers that indicate the beginning of the lists and their management is in fact reduced at advancing these pointers on the circular link. In consequence the first buffer in a list becomes the last buffer in the next list in the order of data processing.

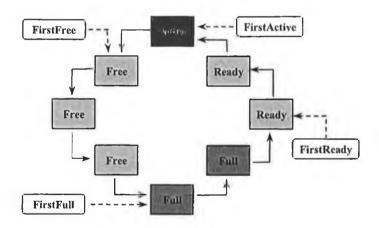


Figure 5-2 Basic structure of the Circular Buffer

Figure 5-2 presents the basic structure of the Circular Buffer system used at the client that includes four pre-defined lists of buffers. The list of free buffers is indicated by the *FirstFree* pointer, the list of buffers with encoded data received via data link by *FirstFull*, the list of buffers

with decoded data waiting to be played out by FirstReady and the list of buffers containing data being played out by FirstActive.

Not all these pre-defined lists have to be used. For example, in the pre-recorded streaming case, the server only uses three lists: one that links the empty buffers, one that stores encoded data read from the files and one that contains data being streamed.

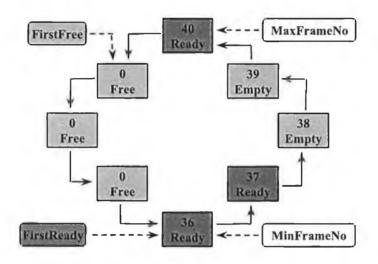


Figure 5-3 Enhanced structure of the Circular Buffer

Apart from storing data the Circular Buffer also helps in acquiring statistical information necessary for the QoDGS during its grading of the quality of streaming. Therefore, apart from the basic features already presented, the Circular Buffer was enhanced with supplementary capabilities and the structure shown in Figure 5-3 was obtained.

Apart from fields that store the corresponding data packet RTP sequence numbers received during transmissions, the enhanced Circular Buffer structure allows also for the continuous computation of the loss rate. It also provides a mechanism that allows the insertion of data packets in the order of their sequence numbers regardless of their order of arrival and a mechanism to communicate network parameter-related information (i.e. loss rate). This enhanced Circular Buffer structure takes into account packets that arrive ahead of their time or too late for their decoding and play out, by either leaving empty buffers, marked with the sequence numbers of the expected packets or by skipping them. *MinFrameNo* and *MaxFrameNo* indicate the minimum and the maximum sequence numbers for packets whose data are stored into the Circular Buffer.

The Circular Buffer is designed for operation as a shared resource between multiple threads. While one of them writes data into the buffer structure, the other reads it and uses it for

another purpose. In consequence a mechanism that protects the integrity of this information was deployed and will be described next.

The CCircBuff class implements the basic Circular Buffer structure and was used at the server, whereas the CNoCircBuff class implements the enhanced Circular Buffer and was used at the client.

5.2.1.3 Complex Producer-Consumer Problem

The producer-consumer problem aims at managing a number of parallel activities that share common resources some using the output of the others. The first aspect of the problem is to organise these activities in such a manner that they will perform continuously. The second is to protect the common resources from being interfered with while they are accessed.

During the prototype system's implementation a solution was found for the client's copier-decoder-player problem, a more complex version of the simple producer-consumer problem. The solution that involves three threads that share two common resources is presented in Figure 5-4, which also indicates some implementation-related objects. The data is read from a file or received through the sockets by a thread (Copier Thread) and placed in a buffer (CpyCircBuffer) from where it is retrieved by a second thread (Decoder Thread), decoded and stored in a second buffer (DecCircBuffer). From DecCircBuffer the decoded data is read and then it is sent to be played out by a third thread (Player Thread).

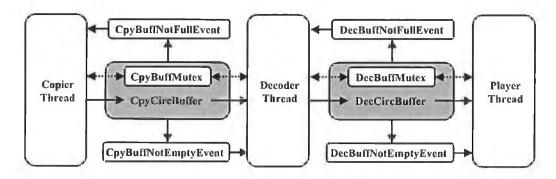


Figure 5-4 Solution for the copier-decoder-player problem

In this example implemented at the client, the shared resources are the two Circular Buffers and mutex objects (CpyBuffMutex and DecBuffMutex) control the access to each of them ensuring that all the operations on the shared structures are mutually exclusive. A pair of event objects associated with each of the Circular Buffers allows for the inter-thread synchronisation of the access

to data. For example if the *CpyCircBuffer* is full, the Copier Thread is sent to sleep and is awaken by the *CpyBuffNotFullEvent* only after the Decoder Thread retrieves some data, emptying a slot. If for example the *CpyCircBuffer* is empty, the Decoder Thread is sent to sleep and is awaken by the *CpyBuffNotEmptyEvent* only when the Copier Thread writes some data in the buffer structure.

The server uses a similar structure, although not as complex, in order to solve its copier-transmitter problem. The solution involves two threads (Copier and Transmitter) and one shared circular buffer structure.

5.2.2 Implementation of the QOAS Server Application

Before giving details about the server application implementation at block level, as presented in the architecture introduced in the fourth chapter, the server's application overall structure is presented next, in terms of defined classes and their roles.

As in the majority of MFC-based multiple document applications, MySrv consists of a main application class CMySrvApp that inherits the CWinApp class, a mainframe window class CMainFrame that extends the CMDIFrameWnd class and a main document-view structure that consists of three classes. The first class inherits the MFC's CDocument class is called CMySrvDoc and provides the container capabilities for the server application-related mechanisms. The second class inherits the CMDIChildWnd class is called CChildFrame and offers the multiple document child window-related functionality. Finally the third class, called CMySrvView, inherits MFC's class CEditView and provides the viewing facility for the server application's user. This structure is very flexible since it allows for a single server application instance to act as a pool of server applications, which share common resources, providing services to a high number of clients.

A second document-view structure was defined in order to implement the server-located Communication Manager. The components of this structure are the *CMpegRemoteTxDoc* class, which provides the container capabilities for most of the mechanisms, the *CChildFrame* class with the role of providing the child window functionality and the *CRemoteTxView* class that offers the viewing capabilities. Instances of the *CMpegRemoteTxDoc* class, that control the client-server communication at the session (control) level, are meant to work in conjunction with instances of the *CMpegSystemSrv* class, which was defined to work at the transport (data transmission) level. Therefore the latter encapsulates the server-side of the client-server communication mechanism-related objects, the server Copier and Transmitter threads, the server located circular buffer and the synchronisation mechanisms for the threads access to the shared resources which have been

previously discussed. It also includes the instances of the *CSrvStateMonitor* class that implements the Server Arbitration Scheme (SAS) and of the *CPRecDatabase* class that implements the database interrogation mechanism.

Next the implementation of each of the server's units is described, indicating the main classes and methods used.

5.2.2.1 Multimedia Acquirer and MPEG Encoder

Although in the block-level architecture of the server application, two units are indicated for audio-video capturing and MPEG encoding, this implementation uses the Canopus Amber MPEG hardware encoder-decoder card for both processes, since it uses directly analog signals as input. The implementation relies on the Canopus Amber's MPEG Development Kit MVR-D2000 version 3.20 [243] and on the associated API library.

For the multimedia acquiring and encoding another MFC-based document-view structure was defined and consists of a document-container class *CEncOverlayDoc*, a *CEncOverlayView* class that provides viewing facilities and a child window class *CEncRemoteOverlayChildFrame*. The latter also provides all the methods directly related to the functionality of the hardware encoding process. This seven-step process is detailed next.

Since multiple encoding-decoding cards can be used at a time, the first step consists of selecting of the next available card and marking it as used. Next, the initialisation of the selected encoder card is performed. This second step indicates some encoding-related parameters such as the type of the encoded stream (MPEG-1 or MPEG-2, System or Program, Audio or Video), the buffer size, the place where to temporarily store the processed data to (memory or disk), input TV system format (PAL or NTSC) and the callback function which is called each time a new set of data is encoded in order to give the application access to it. This is performed in the *OnCreate()* method of the *CEncRemoteOverlayChildFrame* class.

The third step is the creation of the control window that displays the video information being encoded and plays the associated sound. This is performed in the method CreateOverlayWindow(). The fourth step is the encoding, which is started when any of the following CEncRemoteOverlayChildFrame class's methods are called. OnEncodeAudio() was defined for local MPEG Audio encoding, OnEncodeVideo() - for local MPEG Video stream encoding, OnEncodePs() - for local MPEG System or Program stream encoding and

RemoteEncodePs() for live encoding used in conjunction with streaming, function supported only for MPEG System or Program. These functions initiate the encoding thread that acts as a copier in a producer-consumer-like situation, storing data in a local buffer from which a display or a transmitter thread retrieves it for local display and storage on disk or for streaming.

The encoding ends by calling the *OnClose()* method or *RemoteEncodeStop()*, according to the case and includes the destruction of the display control window and by the deactivation of the encoder card.

5.2.2.2 Server Communication Manager and Transmission Shaper

The principle and some implementation details in relation to the server's Communication Manager were described in section 5.2.1.1. Next the classes that implement the server-side communication mechanism are just mentioned.

CListenSocket is created to listen for incoming client link requests and if accepted, one of CSrvTCPSock or CSrvUDPSock classes is instantiated to create an object in charge with the server end of the control link. To this end the client is directed and consequently control messages can be then bi-directionally exchanged with the server. Next the unidirectional data link is established using the server side's CSrvUDPSock class and a similar class at the client and data can be sent to the clients.

The Transmission Shaper, in charge of controlling the data transmission, is implemented by the *CMpegSystemSrv* class that acts as a container for the copier thread (*CSrvCpThread*), the transmission thread (*CSrvTxThread*), the shared transmission circular buffer (*CCircBuff*) and the related mechanisms (i.e. events and mutex-based). These represent a solution for a producer-consumer-like problem that was already described in section 5.2.1.3 and it aims at transmitting data to the clients. It is interesting from the implementation point of view that the copier thread selects the source of data based on the QOAS – Server Arbitration Scheme's (SAS) decisions, taken on the basis of feedback received from the client. These decisions affect the server quality state and therefore the stream the data is read from. The exact position the switch between a source of data and the next one is determined after consulting the Multimedia Database whose implementation mechanism is detailed in section 5.2.2.4. Also the transmission thread adjusts the streaming rate according to the current quality state of the server.

The server Communication Manager is also controlling the multimedia session using RTSP VCR-like commands. The server receives the client commands and processes them in the CMpegRemoteTxDoc's ExecuteCommand() method that is invoked by the CSrvCtrlSock's OnReceive() method, automatically called by the framework when a new data packet is received. It then sends the acknowledgements by calling the CMpegRemoteTxDoc's SendMessage() method. ACK_SETUP, in form of PUT_CTRL_CONNECT and PUT_DATA_CONNECT, confirms the setup process completion, whereas ACK_PLAY (PUT_PREC_FILE) acknowledges starting streaming of the requested multimedia content. The server can initiate the destruction of the double link and the end of the session by sending TEARDOWN (SRV_SHUT_DOWN).

5.2.2.3 Feedback Manager and Server Application Core

One of the most important mechanisms operated in conjunction by the Feedback Manager and the Server Core is the Server Arbitration Scheme (SAS), which is implemented by the *CSrvStateMonitor* class. In the server application implementation the *CMpegSystemSrv* container class stores a reference to an object that instantiates *CSrvStateMonitor* class.

Among the most significant methods of the *CSrvStateMonitor* class are *AddFeedbackQoDGrade()* that processes a newly received QoD_{Grade} from the client and assesses the opportunity of a server quality state change and *ChangeSrvState()* that effectively changes this state.

Another important mechanism implemented in co-operation by the Feedback Manager and the Server Core is the quality timeout mechanism that determines a server quality state decrease if there is no information received from the client in form of feedback for a duration of time. This mechanism uses a timer from the MFC's timer pool facility built-in the CView class which is inherited by our *CRemoteTxView* class. The latter's *OnTimer()* function is called every time timeout occurs and measures have to be taken. The timer is restarted every time the server's Feedback Manager receives a feedback message from the client.

5.2.2.4 Database Support for Pre-recorded Streams

The goal of the database support is to provide a mean for storing, accessing and updating pre-recorded multimedia streams related information, necessary for the implementation of QOAS. This database implementation makes use of the Microsoft Open Database Connectivity (ODBC) a vendor-independent mechanism that allows access to data stored in a variety of data sources [244]

(including Microsoft Access used here) by executing SQL (Structured Query Language) statements against them.

In order to allow for any application to access a database with the ODBC mechanism, the database has to be first registered with the ODBC Administrator from the Windows Control Panel, indicating also the driver it can work with, as advised in [245]. The registration name for the database can be different than the real name of the database file and it is meant to be used by any application that wants to access the registered database.

The defined class *CPRecDatabase* which inherits the MFC's *CDatabase* class, provides methods (*OpenPRecDbConnection(*) and *ClosePRecDbConnection(*)) for connecting and disconnecting from the registered database. Once a connection has been made, operations on the data source are possible using either an object that instantiates the MFC's *CRecordset* class or executing the *CPRecDatabase ExecuteSQL(*) member function. Since the latter does not return any result, only creation and deletion of tables, insertion of data can be performed, whereas for the database queries the MFC's *CRecordset* class is used. Among other methods is *CreatePRecDbTable(*) which creates a new table in the database with the given name. The newly created table will have its fields named and of the types as indicated in the transmitted parameters. *ExistsPRecDbTable(*) verifies whether or not the indicated table exists and if it does not, it creates it in a similar manner with the previous function. *OpenPRecDbTable(*) and *ClosePRecDbTable(*) open and close the indicated table, while *DeletePRecDbTable(*) deletes it.

The MFC's CRecordset class provides the functionality of an ODBC SQL statement, including the returned row set by the statement. CreatePRecRecordset() DeletePRecRecordset() are creating and destroying the object used to operate on the database. OpenPRecRecordset() and ClosePRecRecordset() allow for selecting a set of records from the database that fulfil the indicated constraints for further processing and respectively renounce at the selected set once the task has been performed. Its usage is very simple and employs two query structures DbIntResult and DbStrResult. A call to OpenPRecRecordset() function requires a pointer to the CPRecDatabase object that manages the connection to the data source and two strings Filter and Sort that correspond to SQL's WHERE and ORDER BY clauses. The most important result is the SQL's selected set of records that is processed and essential information collected and retrieved via the query structures.

5.2.3 Implementation of the QOAS Client Application

As in the case of the server application, the implementation of the client application follows the block level architecture presented in the fourth chapter, section 4.2. Next the defined classes are presented and how they are used is discussed.

At the core of the *MyCli* application is the application class *CMyCliApp* that inherits MFC's *CWinApp* class, a mainframe window class *CMainFrame* that inherits MFC's *CMDIFrameWnd* class and a document-view structure. The three document-view classes are *CMyCliDoc* that provides the container capabilities for the client-related mechanisms, *CChildFrame* that offers the child window related functionality and the *CMyCliView* that provides the viewing facility for the client's user. The flexibility of this structure allows for the existence of multiple clients at the same time, which can connect to different server application instances and stream different multimedia content if wanted. In this way the single application acts as a pool of clients that run in parallel and share some resources.

A second document-view structure was defined in order to implement the client-located Communication Manager. The components of this structure are the *CMpegRemoteCliDoc* class, which provides the container capabilities for the most of the client-located streaming mechanisms, the *CChildFrame* class with the role of providing the child window functionality and the *CMpegRemoteCliView* class that offers the viewing capabilities. An extra view has been also added for real-time monitoring of the transmission-related parameters, class called *CRemoteRxView*.

Any instance of the *CMpegRemoteCliDoc* class, that controls the client-server communication at the session (control) level, is meant to work in conjunction with an instance of the *CHardSystemCli* class, which was defined to work at the transport (data transmission) level. The latter encapsulates the client-side of the client-server communication mechanism such as the client receiver (copier), decoder and player threads, two instances of the circular buffer and the associated synchronisation mechanisms, previously presented. It also includes the instances of the *CQoDGSMonitor* class that implements the Quality of Delivery Grading Scheme (QoDGS).

The implementation of each of the client's units is described next, indicating the main classes and methods used.

5.2.3.1 MPEG Decoder and Multimedia Player

Although in the architecture presented in the fourth chapter two separate units were indicated for MPEG decoding and respectively multimedia play-out, this implementation uses the Canopus Amber MPEG hardware decoder card for both processes. In a similar manner with the server application, the client relies on the Canopus MVR-D2000 Amber Development Kit version 3.20 [243] and its associated API.

Similar to the server-based acquiring and encoding processes, the MPEG decoding and multimedia playing has seven steps, which are described in detail next. They make use of an MFC-based document-view structure that was especially defined and consists of a document-container class CDecOverlayDoc, a CDecOverlayView class that provides viewing facilities and a child window class CDecRemoteOverlayChildFrame that contains all the methods used for decoding and playing as well as for the interaction with the rest of the client application units.

Since more than one MPEG encoding-decoding card can be used at a time, the first two steps consist of the selection of the next available card and its initialisation and are implemented in the *CDecRemoteOverlayChildFrame*'s method *OnCreate()*. The initialisation sets the type of the stream to be decoded (the card accepts only MPEG System or Program), the buffer size, the place the encoded data is read from (memory or disk), the TV systems (PAL or NTCS) and indicates the callback data-related function which is called each time a new piece of data is decoded, the error callback function and the status callback function.

The third step is the creation of the control window in which the streamed multimedia information is played out for the viewer and this is done in the *CreateOverlayWindow()* method of the *CDecRemoteOverlayChildFrame* class. The next step, the real decoding and playing of the multimedia data starts when the *OnDecodePlay()* method is called. This creates the copier-decoder-player thread and buffer structure that supports the play out. During the play out process, the card's driver expects to be fed with encoded data by the application and this is performed by the data callback function, which acts as the decoder thread.

The last stages involve the termination of the decoding process, freeing of the associated resources, including the destruction of the card's display control window and the deactivation of the encoder card. These are performed by the *CDecRemoteOverlayChildFrame's CloseFrame()* member function.

5.2.3.2 Client Communication Manager

The client-side Communication Manager supports in conjunction with the server's the double-link communication principle presented in section 5.2.1.1 First the client creates a control socket instantiating either of the classes *CCliTCPSocket* or *CCliUDPSocket* and tries to communicate with the server's listening socket described by an IP address-port number pair. If accepted by the server, a control link is established between the two of them. Using this link, next the unidirectional data link is created using the client side's *CCliUDPSocket* class and a similar class at the server for data to be received by the client.

The Communication Manager is also in charge of controlling the data reception and play out. This mechanism is implemented by the CHardSystemCli class that acts as a managing point for the receiver thread (CCliUDPSocket's method OnReceive()), the decoder thread (CDecRemoteOverlayChildFrame's callback function ReadRemoteSectorProc() that was associated to the decoder card), the player thread (started by a call to the CDecRemoteOverlayChildFrame's DecodeRemotePlay() method), the shared buffers (CCircBuff and the card driver's buffer) and the related synchronisation mechanisms (i.e. events and mutex-based). These represent a solution for the complex producer-consumer-like problem that was already described in section 5.2.1.3 and here aims at receiving data, decode it and play it out to the viewer.

The client Communication Manager also implements the client-side RTSP-related multimedia streaming session control. The client sends session control commands such as SETUP (GET_CTRL_CONNECT and GET_DATA_CONNECT) for establishing the double client-server communication link and PLAY (GET_PREC_FILE) for requesting multimedia streaming of certain content. It also uses STOP (STOP_PREC_FILE) for stopping the streaming and TEARDOWN (SRV_SHUT_DOWN) for initiating the destruction of the double link and the end of the session. Processing the server's answers is performed in the *CMpegRemoteCliDoc*'s member function *ExecuteCommand()* invoked from the *CCliCtrlSocket*'s *OnReceive()* method, which is automatically called by the framework when a new packet is received.

5.2.3.3 Feedback Indication Unit and Client Core

One of the most important mechanisms provided in conjunction by the Feedback Indication Unit and the Client Core is the Quality of Delivery Grading Scheme (QoDGS), which is implemented by the *CQoDGSMonitor* class according to the principles described in section 4.5. A

reference to the object that instantiates this class is a member of the CHardSystemCli container class.

Among the most significant methods of the *CQoDGSMonitor* class are *Update_QoD_After_ReceivePkt()*, *Grade_QoD()* and *InitQoDGSStruct()*. The first method processes the information related to a newly received packet and updates the monitored parameters' values and variations as well as the estimated quality of delivery. It bases its functionality on calls to *Adjust_QoD_XXX()* functions, where XXX stands for the name of the monitored parameter (e.g. Delay, Jitter, Loss, etc.). The second method computes the QoD_{Score}-s after computing partial scores for the monitored parameters by calling their associated methods with the following form: *Grade_QoD_XXX()*. The third method initialises the QoDGS structure and all its components by calling individual functions like *Init_QoD_XXX()*.

These computed QoD_{Scores} are sent to the server as part of feedback messages, using the control link provided by the client and server Communication Managers, informing it about the quality of delivery and allowing it to take adaptive measures.

5.3 Summary

This chapter has presented implementation details about both the simulation system model and the real prototype systems that were built in order to test the proposed QOAS for multimedia streaming using simulations and respectively emulations. The simulation environment used was Network Simulator version 2 and the implementation relies on some of the classes provided by it for lower level services such as transport-level communication. The programming environment used for building the prototype system was Microsoft's Visual C++ version 6.0 and the implementation made use of its MFC class structure. In both cases, this chapter has presented server and client-located mechanisms in terms of architectural blocks and details about their implementation in terms of classes and their main methods whose role was explicitly indicated.

Chapter VI Experimental Results

Abstract

The sixth chapter presents experimental results of tests that aim at both tuning the QOAS parameters in order to obtain best results in local broadband multi-service IP-networks and testing it in different delivery conditions to make sure that very good performances were achieved. The scheme is tested both objectively using a simulation environment and simulation models and subjectively using a prototype system and human subjects. The experimental test results are presented and commented on.

6.1 Overview

QOAS was proposed as an inexpensive solution to deliver high quality multimedia-based services to home residences via the local broadband multi-service IP-network. During its design many issues were taken into account in order to better achieve this goal, as presented in the previous chapters. However, testing is required both during and following the design phase, helping to propose a good solution and in the final stage for its verification and validation. In order to do this, a model was built and extensive **objective testing** was performed, involving simulations. Although these simulations allow for measuring diverse parameters and assessing the performances of the QOAS in various conditions, even in comparison to other solutions, they can only estimate the users' perceived quality with a certain degree of accuracy based on some metrics. Since currently there is not a total agreement that any of these metrics may reflect the viewers' opinions in a wide range of situations, **subjective tests** were also performed in order to determine the real users' perceptual assessment of the streamed multimedia clips' quality. Next both these sets of tests are presented and their results are commented.

6.2 Objective Testing

6.2.1 Simulation-based Testing

There is a growing recognition within the research communities of the importance of simulation tools in helping to design and test different proposed algorithms and schemes. New mechanisms, especially in networking research, present huge challenges for testing. These challenges are mainly related to the required large and complex environments. Both the designers and the testers recognise the significant advantages of simulations in terms of necessary computing resources, associated costs and convenience over duplication of a real word system in the lab. Simulations allow performing large-scale tests that are controlled, reproducible and do not involve significant costs. Therefore it is easy to explain why research in general and networking research in particular increasingly depends on simulation to investigate new schemes' behaviours, performances, and interactions.

There are different simulation tools or environments available for network simulations such as Network Simulator (NS-2) [246], OPtimized Network Engineering Tool (Opnet)⁶³, SIMSCRIPT II (former COMNET III)⁶⁴ and CNET⁶⁵. NS-2 was chosen because it is open-source and allows for easy extensibility, but mainly because it includes a large number of models for different layer protocols, traffic etc. such as UDP and TCP that can be used during testing.

The performed simulations include two different phases: tuning and testing. Since the QOAS design involves the existence of a number of parameters that have to be tuned in order to achieve best results in a given infrastructure, the tuning phase aims at determining these parameters' values. The parameters whose values have to be assigned are the weights associated to the QOAS's Quality of Delivery Grading Scheme (QoDGS), which was presented in the fourth chapter. The testing phase makes use of the tuning stage's results and aims at showing that the QOAS solution achieves expected performances, even in comparison to other multimedia streaming approaches in different network delivery conditions and subject to various cross traffic. In this section both tuning and testing phases are presented.

⁶³ OPtimized Network Engineering Tool (Opnet), OPNET Technologies Inc., http://www.opnet.com

 $^{^{64}\,}SIMSCRIPT,\,CACI\,International\,Inc.,\,http://www.caciasl.com/products/simscript.cfm$

⁶⁵ The CNET network simulator, http://www.csse.uwa.edu.au/cnet/

Next the simulation environment, the aims of the performance assessment, the simulation topology, the QOAS model and the multimedia clips used for objective testing are presented.

6.2.1.1 Network Simulator Version 2 (NS-2)

Network Simulator version 2 (NS-2) [246] is an open source, object-oriented, discrete event, network simulation environment that was built at University of California at Berkely⁶⁶ in order to test models proposed in the networking research area. It was developed and written in C++ and OTcl (an object-oriented extension to Tcl/Tk proposed at Massachusetts Institute of Technology) [247], but in order to be deployed, it also requires Tcl, Tk, OTcl and TclCL to be installed. NS-2 is primarily used for simulating local and wide area IP-networks. It implements network protocols such as TCP (different flavors) and UDP, traffic source behavior such as FTP, Telnet, WWW, constant bit-rate (CBR) and variable bit-rate (VBR), router queue management mechanisms like Drop Tail, Random Early Drop (RED) and Class-Based Queuing (CBQ), routing algorithms such as Dijkstra, and more. NS-2 also supports multicasting and some of the MAC layer protocols for LAN simulations. The NS project is now a part of the Virtual InterNetwork Testbed (VINT) project⁶⁷ and is supported by DARPA⁶⁸. More information about NS-2 can be found in the NS Manual [248] or in one of the NS tutorials such as [249, 250].

6.2.1.2 Simulation Topology

Regardless of the exact architecture of the local broadband multi-service IP-network chosen for the distribution of services (including multimedia-based ones) from the distribution hub to the residential users, there is in fact a single link that has to support all the traffic exchanged. Therefore, the problem is in fact the QOAS-based distribution of high quality multimedia to users behind a single bottleneck link situated between the distribution hub and the residential users. This bottleneck link may cause problems in term of increased or variable delays and/or loss that severely affect the quality of the provided services, especially of the multimedia-based ones. These problems not only originate in the multimedia traffic, but also in other types of traffic, with different size and variation patterns produced by other type of services provided through the same infrastructure. Therefore the problem of delivering multimedia-based services with little effort to residential users

⁶⁶ University of California at Berkeley, http://www.berkeley.edu

⁶⁷ Virtual InterNetwork Testbed (VINT) project, http://www.isi.edu/nsnam/vint/index.html

⁶⁸ Defense Advanced Research Projects Agency (DARPA), http://www.darpa.mil

or businesses becomes the problem of providing the same services to a group of customers behind a single bottleneck link traversed by traffic of different types, sizes and variation patterns. The "Dumbbell" topology presented in Figure 6-1 best describes this situation since assumes a single shared bottleneck link traversed by all the traffic.

The sources of traffic are located on one side of the bottleneck link, whereas the receivers are on the other side. Among these sources of traffic is the QOAS server that is modelled as a number of adaptive sources that are associated with corresponding QOAS receivers situated across the bottleneck link. The other sources produce the background traffic. Apart from the bottleneck link, the other links are provisioned in such a manner that the only drops and significant delays are caused by congestion that occurs at the bottleneck link. In this context the S_i-B₁ and B₂-C_i links have been assigned 200 Mb/s bandwidth and 0.005 s propagation delay. The buffering at the bottleneck link uses a drop-tail queue of length proportional with the product between the round trip time and the bandwidth of the bottleneck link. During simulations this bandwidth was set to 100 Mb/s, which ensures both an as real as possible situation that can also allow for easy generalisation to gigabit Ethernet and an average complexity of the simulations. The bottleneck link's delay was set to 0.1 s, allowing for testing the feedback-based QOAS for average-high latencies in local area networks.

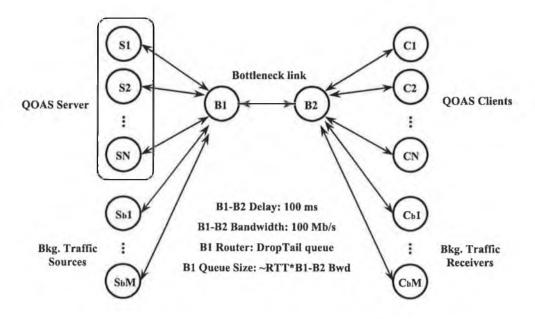


Figure 6-1 The "Dumbbell" topology includes a bottleneck link, a QOAS server and N QOAS receivers (clients), as well as a number of sources and receivers of background traffic

6.2.1.3 QOAS Model

The QOAS model used during simulations conforms to the general description made in the fourth chapter and is the result of the simulation implementation described in the fifth chapter. The QOAS server arbiter upgrade period was set to 6 s and the associated downgrade timeout was set to 1 s. As mentioned in the design, the QoDGS's short-term period and its long-term period were taken an order and two orders of magnitude greater than the feedback interval respectively. For 0.1 s inter-feedback transmission time, they are considered 1 s and 10 s respectively.

6.2.1.4 Multimedia Clips

In order to ensure a large range of types of multimedia clips for the simulations, five five-minute long video sequences were selected from movies with different degrees of motion content. The *diehard1* sequence includes a great deal of action, *jurassic3* and *dontsayaword* have average motion content, *familyman* has very little movement, whereas *roadtoeldorado* is a cartoon sequence, with average-high motion content. Since during testing five possible quality states for the QOAS application server were considered, these clips were MPEG-2 encoded at five different rates, equally distributed between 2 Mb/s and 4 Mb/s, using the same frame rate (25 frames/sec) and the same IBBP frame pattern (9 frames/GOP). Traces were collected, associated with the QOAS server states, and stored in the QOAS indexing database to be used during simulations. Table 6-1 presents statistics about all the quality versions of these multimedia clips used during testing. Table 6-2 presents a categorisation of these multimedia clips based on their 2.0 Mb/s versions.

Table 6-1 Peak/mean ratio for all the MPEG-2 encoded quality versions associated to the multimedia clips used during simulations

Quality Version (average rate) Clip Name	2.0 Mb/s	2.5 Mb/s	3.0 Mb/s	3.5 Mb/s	4.0 Mb/s
diehard1	7.48	7.43	6.31	5.65	4.06
roadtoeldorado	6.91	6.51	6.23	6.12	6.05
dontsayaword	5.56	4.51	4.36	4.08	3.56
jurassic3	4.83	4.38	4.04	3.71	3.41
familyman	3.99	3.67	3.42	3.09	2.93

The bitrates used for encoding ensure both a compromise between the quality of the streamed content and its corresponding bandwidth for necessary transmission and a balance between the degree of flexibility related to possible adaptations during streaming and the required storage space in the multimedia database.

Table 6-2 Categorisation of the multimedia clips used during simulations (based on their 2.0 Mbits/s MPEG-2 encoded quality versions)

Clip	Motion Content	Content	Peak Rate (bits/frame)	Peak/Mean Ratio
diehard1	High	movie	860648	7.48
roadtoeldorado	average-high	cartoons	693696	6.91
dontsayaword	average	movie	480840	5.56
jurassic3	average-low	movie	447528	4.83
familyman	Low	movie	322968	3.99

6.2.1.5 Performance Assessment

The performance of the QOAS-based adaptive solution is assessed in terms of loss, bottleneck link utilisation, estimated end-user perceived quality, and the number of clients simultaneously served from a fixed infrastructure. The loss rate refers to packet loss as measured at the receiver, whereas the link utilisation is computed in terms of current throughput versus existing bandwidth. The estimated end-user perceived quality is computed using the no-reference moving picture quality metric (Q) proposed in [133] and described in the second chapter, equation (2-5). The results are expressed on the five-point scale for grading subjective perceptual quality suggested in the ITU-T R. P.910 [63] and presented in Table 6-3. The number of simultaneous served clients is computed while maintaining a certain perceived quality for the served multimedia-based services. However, by differently setting the target limit related to the accepted end-user perceived quality degradation, the number of customers served simultaneously varies.

Rating **Impairment** Quality 5 Imperceptible Excellent 4 Perceptible, not annoying Good 3 Slightly annoying Fair 2 Annoying Poor 1 Very annoying Bad

Table 6-3 Quality scale for subjective testing

The results presented in the following sections rely on a sampling frequency of 0.1 s. The loss rates are expressed in percentage (%) with values from 0 to 100, whereas link utilization is computed and presented as a fraction, with values from 0 to 1.

6.2.2 Tuning QOAS

In order for the proposed QOAS to achieve the best performance in a given infrastructure, a tuning phase has to be performed. This tuning phase aims at determining the values of a number of weights used by the client-located QoDGS in its process of mapping network-related parameters variations into application level quality scores. These weights were presented in the fourth chapter where QoDGS was described. For simplicity it was considered $w_i = w'_i$, $1 \le i = 4$. Next, the following notations are made: $w_l = w'_l = W_{Delay}$, $w_l = w'_l = W_{Jitter}$, $w_l = w'_l = W_{Loss}$ and $w_l = w'_l = W_{Ql}$. These notations better reflect the association between the weights and the QoDGS's parameters.

The tuning process has two steps. First the tuning process aims at determining the range of values that these weights have to belong to in order to achieve best end-user perceived quality. Then values from within the suggested intervals are selected and tested in order to verify the expected results.

For each of the QoDGS parameters, the tuning process experimentally finds lower and upper limits for their contribution in the overall application-level QOAS scores that determines the highest end-user perceived quality for the streamed multimedia sequence. Therefore four issues were taken into account: streamed content, traffic conditions, quality variation patterns and background traffic type.

Three multimedia sequences with different motion content (average, high and low) were selected for streaming: jurassic3, diehard1, and familyman and the starting points were chosen at random within each sequence. These tests make use of the "Dumbbell" topology, which was presented in Figure 6-1. Since QOAS is designed to work in highly loaded network conditions, such a situation was created by generating a constant bit-rate (CBR) background traffic of 95.5 Mb/s using the NS-2 CBR traffic model. This CBR background traffic represents a well-multiplexed traffic composed of a high number of individual different types, shapes and variation patterns of data flows, as expected to happen in a local multi-service broadband IP-network. On top of this large background traffic a highly variable high-quality multimedia-like traffic presented in Figure 6-2 is transmitted across the delivery network. This traffic simulates all possible effects of user interactions to multimedia streams such as play, pause, re-play and stop. It even takes into account the effect of multiple consecutive play commands that increase the traffic in a staircase up manner, consecutive pause-play interactions with different frequency and applied on movies with different rate and consecutive stop requests by different viewers that decrease the traffic in a staircase-down fashion. This traffic was considered representative for this case since interactive controlled multimedia should account for the majority of the traffic carried by this local network.

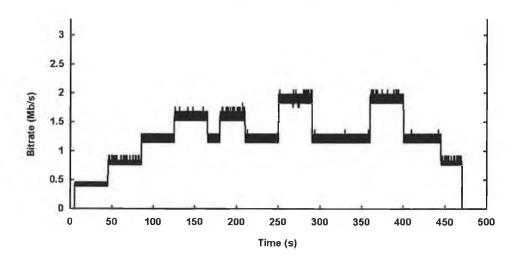


Figure 6-2 Background traffic variation on top of 95.5 Mb/s CBR traffic

This variable background traffic, on an already loaded delivery link, determines quality adaptations from the QOAS-based stream and consequent quality variations of the streamed multimedia sequence. The resultant end-user perceived quality is measured and averaged for the duration of the tests (500 sec), with the exception of two transitory periods of 50 sec at the beginning and at the end.

The first set of tests involved the usage of the sequence *jurassic3* with average motion content. The contribution of the Delay is first varied in steps of 20% between 0% and 100%, by changing the value of W_{Delay} between 0.0 and 1.0. The other parameters were equally sharing the remaining contribution. However, at all times, the equation (6-1) was respected.

$$W_{Delay} + W_{Jitter} + W_{Loss} + W_{Q} = 1 ag{6-1}$$

Table 6-4 presents the average end-user perceived quality for the duration of these tests, as estimated by the no-reference metric Q at the receiver. Analyzing the results, the best average perceived quality is obtained for a contribution of the Delay in the interval 80-100%. For achieving better granularity, a supplementary test was performed for $W_{Delay} = 0.9$, revealing that the best results were obtained in the 90-100% interval in the QoDGS process of QoD scores' computation. The interval and values are marked in the next table.

Table 6-4 Average end-user perceived quality when varying W_{Delay} in QoDGS

W _{Delay} (%)	0	20	40	60	80	90	100
Avg. Quality (1-5)	4.194	4.312	4.428	4.288	4.444	4.462	4.467

A similar set of tests involved the variation of Jitter's contribution in the computation of QoD scores at the QoDGS level and the results from Table 6-5 were obtained. The best results in terms of average estimated end-user perceived quality were obtained for a Jitter contribution in the QoDGS grading process between 50% and 60%.

Table 6-5 Average end-user perceived quality when varying W_{Jitter} in QoDGS

W _{Jitter} (%)	0	20	40	50	60	80	100
Avg. Quality (1-5)	4.335	4.252	4.346	4.406	4.367	4.136	1.927

When Loss rate's contribution was varied in the same manner, the best results were obtained for the interval 30-40% as shown in Table 6-6.

Table 6-6 Average end-user perceived quality when varying W_{Loss} in QoDGS

W _{Loss} (%)	0	20	30	40	60	80	100
Avg. Quality (1-5)	4.186	4.168	4.212	4.206	3.952	4.043	4.098

The last set of tests in this first stage involved variations in the contribution of the QoDGS's parameter Q. The results suggest for W_Q an interval between 0.0 and 0.1, as presented in Table 6-7.

Table 6-7 Average end-user perceived quality when varying Wo in QoDGS

W _Q (%)	0	10	20	40	60	80	100
Avg. Quality (1-5)	4.457	4.466	4.336	4.410	4.329	4.222	4.244

Table 6-8 concludes these experimental test results, indicating both minimum and maximum limits for these "best results" intervals. It also presents the remaining contributions for the QoDGS parameters whose values were not varied for the duration of these individual tests.

Table 6-8 Minimum and maximum limits for the QoDGS weights when the highest end-user perceived quality was achieved during QOAS-based streaming of the average motion content movie *jurassic3*

Variation	W _{Delay}	(%)	W _{Jitter}	(%)	W _{Loss}	(%)	W_Q	(%)
	Min	Max	Min	Max	Min	Max	Min	Max
Delay	90	100	0	3.33	0	3.33	0	3.33
Jitter	13.33	16.67	50	60	13.33	16.67	13.33	16.67
Loss	20.00	23.33	20.00	23.33	30	40	20.00	23.33
Q	30	33.33	30	33.33	30	33.33	0	10
Average	38.33	43.33	25.00	30.00	18.33	23.33	8.33	13.33

Based on these suggestions, by averaging the values of the limits that correspond to the same parameter, the following intervals were suggested for the weights that correspond to the QoDGS parameters: W_{Delay} between 0.383 and 0.433, W_{Jitter} between 0.25 and 0.30, W_{Loss} between 0.183 and 0.233 and W_Q between 0.083 and 0.133 (see Table 6-8).

The same tests were repeated using the *diehard1* sequence, with high motion content and the results are presented in Table 6-9. The suggested intervals are the following: W_{Delay} between 0.383 and 0.433, W_{Jitter} between 0.283 and 0.333, W_{Loss} between 0.15 and 0.20 and W_Q between 0.083 and 0.133.

Table 6-9 Intervals for QoDGS weights when QOAS has achieved the highest end-user perceived quality when streaming the high motion content movie *diehard1*

Variation	W _{Delay}	(%)	$\mathbf{W}_{\mathrm{Jitter}}$	(%)	W _{Loss}	(%)	W_Q	(%)
	Min	Max	Min	Max	Min	Max	Min	Max
Delay	90	100	0	3.33	0	3.33	0	3.33
Jitter	10	13.33	60	70	10	13.33	10	13.33
Loss	23.33	26.67	23.33	26.67	20	30	23.33	26.67
Q	30	33.33	30	33.33	30	33.33	0	10
Average	38.33	43.33	28.33	33.33	15	20	8.33	13.33

The third time the same tests were performed using the low motion content sequence familyman. The suggested limits, presented also in Table 6-10, are as follows: W_{Delay} between 0.367 and 0.417, W_{Jitter} between 0.267 and 0.317, W_{Loss} between 0.167 and 0.217 and W_Q between 0.10 and 0.15.

Table 6-10 Suggested limits for QoDGS weights in tests that have involved streaming using QOAS of the low motion content movie: familyman

Variation	$\mathbf{W}_{ ext{Delay}}$	(%)	W _{Jitter}	(%)	W _{Loss}	(%)	\mathbf{W}_{Q}	(%)
	Min	Max	Min	Max	Min	Max	Min	Max
Delay	80	90	3.33	6.66	3.33	6.66	3.33	6.66
Jitter	13.33	16.67	50	60	13.33	16.67	13.33	16.67
Loss	23.33	26.67	23.33	26.67	20	30	23.33	26.67
Q	30	33.33	30	33.33	30	33.33	0	10
Average	36.67	41.67	26.67	31.67	16.67	21.67	10.00	15.00

Averaging the limits obtained when using different motion content movies, the following suggestions for the intervals are made: W_{Delay} between 0.378 and 0.427, W_{Jitter} between 0.268 and 0.317, W_{Loss} between 0.167 and 0.217 and W_Q between 0.089 and 0.139 (see Table 6-11). Normally the equation (6-1) has to be respected when choosing the weights' values within these limits.

$\mathbf{W}_{ ext{Delay}}$	(%)	W _{Jitter}	(%)	W _{Loss}	(%)	W_Q	(%)
Min	Max	Min	Max	Min	Max	Min	Max
37.77	42.77	26.67	31.67	16.67	21.67	8.89	13.89

Table 6-11 Suggested contributions for the parameters in the QoDGS

Apart from the values of these weights, other weights have also to be assigned values. w_A and w_B determine the contribution of the short-term and respectively the long-term monitoring and grading of the parameters' values, variations and variation patterns in the total QoDGS-based scoring process.

The existing research such as [90, 103, 118] indicates that, for any network-related parameters' monitoring, the closer the monitored period to the present, the more accurate the estimation is. Since the variations' patterns can be considered only during long-term monitoring of parameters, some of these works also suggest taking long-term monitoring into account. However, an increased importance should be given to short term monitoring. Taking this advice into account, Table 6-12 presents broad limits for the contributions of short-term and long-term monitoring in the overall QoDGS scoring process. These limits impose the following constraints on the values of the weights: w_A between 0.6 and 0.9 and w_B between 0.1 and 0.4.

Table 6-12 Suggested contributions for short-term and long-term monitoring in the QoDGS

Short-term	Short-term (%)		(%)
Min	Max	Min	Max
60.00	90.00	10.00	40.00

The second step of tuning aims at determining more precise values for all the QoDGS-related weights.

In the intervals suggested by the first step of the tuning process the following set of weights were selected such as they respect the equation (6-1): $W_{Delay} = 0.4$, $W_{Jitter} = 0.3$, $W_{Loss} = 0.2$ and $W_Q = 0.1$. Testing the QOAS for the transmission of the average motion content clip *jurassic3*, the average estimated user-perceived quality was 4.479, the highest recorded during testing with this clip. When QOAS was used for streaming the *diehard1* and *familyman* clips with high and low motion content, respectively, the averages for the estimated quality were similarly very high (4.384 and 4.489 respectively).

In order to determine the best values for w_A and w_B , since the suggested interval is very large, further tests have been performed. These tests have varied the w_A and w_B values within the suggested interval and were performed for three multimedia sequences with different motion content. The resulting performance of the QOAS adaptation was assessed in terms of average estimated end-user perceived quality. Table 6-13 presents these results, highlighting also the best results obtained for each type of movie. The results indicate a significant improvement of the end-user perceived quality for the high-motion content clip when the contribution of the QoDGS's short-term grading increases, whereas for the low motion content clip the estimated end-user quality becomes better when the long-term grading has a more important contribution. However, in both cases, as well as for the QOAS streaming of an average motion content clip, good end-user perceived quality was obtained for w_A between 0.7 and 0.8 and w_B between 0.2 and 0.3. Therefore we have selected $w_A = 0.75$ and $w_B = 0.25$ within these intervals, respecting also the condition: $w_A + w_B = 1$.

Table 6-13 Average end-user perceived quality when varying QoDGS's w_A and w_B for different motion content movies: *jurassic3*, *diehard1* and *familyman*

Weights' values	Avg. motion clip	High motion clip	Low motion clip
$w_A = 0.6$ $w_B = 0.4$	4.306	4.304	4.479
$w_A = 0.7$ $w_B = 0.3$	4.478	4.460	4.479
$w_A = 0.8$ $w_B = 0.2$	4.476	4.455	4.346
$w_A = 0.9$ $w_B = 0.1$	4.352	4.455	4.298

6.2.3 Testing QOAS

After QOAS was tuned for this infrastructure based on experimental tests and values for the six QoDGS weights were determined, these results have to be validated by testing the QOAS in different delivery conditions. During the tests the QOAS's performance related to multimedia delivery alone and in comparison with other streaming solutions was assessed. These tests include a single QOAS-based delivery of a multimedia stream in increased traffic conditions and in the presence of background traffic of different types, shapes and variation patterns. They also involve assessing QOAS-related performance when streaming a single multimedia stream in parallel with multiple other interactive multimedia streams. The delivery of multiple QOAS-based adaptive streams over the same infrastructure is also tested and the consequent benefits are assessed. As previously mentioned these tests aim at analysing the performance of the QOAS in terms of loss, bottleneck link utilisation, estimated end-user perceived quality and the number of clients simultaneously served from a fixed infrastructure.

6.2.3.1 Single QOAS-based Streaming Against Different Types of Traffic

The first set of tests aims at assessing the performance of the delivery of a single multimedia stream using QOAS in increased traffic conditions and in the presence of background traffic commonly expected in IP-based networks of different types, shapes and variation patterns.

As for all the objective tests, NS-2 is the simulation environment and more details about it were presented in section 6.2.1.1 of this chapter. The "Dumbbell" topology used for these tests, as well as for the other simulation tests, was presented in section 6.2.1.2 and more details about the tested QOAS model were given in section 6.2.1.3. Section 6.2.1.5 states the principles followed for the scheme's performance assessment, whereas section 6.2.1.4 presents the multimedia clips used during testing.

Since NS-2 includes models for many protocols and can generate different types of traffic, as it would be outputted by real applications, these models were used in these tests to generate different types of background traffic commonly expected in IP-based networks with different shapes and sizes. UDP-based (constant bit-rate - CBR and variable bit-rate - VBR) and TCP traffic are two main classes of traffic taken into account on top of a CBR traffic of at least 95.5 Mb/s that corresponds to a well multiplexed high load that ensures highly increased traffic conditions for these tests. Each of these classes is presented next, as well as the different types of traffic taken into account for each of them and the effect on the QOAS streaming.

6.2.3.1.1 UDP - CBR as Background Traffic

Some multimedia streaming solutions use smoothing techniques in order to reduce the burstiness of the traffic generally associated with multimedia deliveries, whereas some others use CBR encoding to produce a flat rate output stream that would be easily transmitted over IP networks. Also if the traffic is very large, even if consisting of different types and shapes of individual flows, it is subject to a process of statistical multiplexing that produces an almost CBR output. The effect of these traffic types is studied in this section, taking into consideration different variation shapes such as periodic, staircase up and staircase down, with different frequency and variation step size, related to the size of the QOAS adaptation step which is 0.5 Mb/s.

6.2.3.1.1.1 CBR Periodic

The step size of the CBR periodic traffic variation is set to 0.5 Mb/s in the first set of tests and to 0.7 Mb/s in the second set while the frequency of the periodic variations is varied. These traffic variations are made on top of the 96 Mb/s CBR background traffic that represents a well-multiplexed traffic and aims at creating high loaded network delivery conditions. Next the results of the QOAS-based adaptive multimedia streaming subject to CBR periodic background traffic that varies with different frequencies such as 20 s on and 40 s off, 30 s on and 60 s off and 40 s on and 80 s off are presented and are commented on. These results aim at presenting the QOAS-driven adaptation in relation to the background traffic variation, the estimated end-user perceived quality using the no-reference moving pictures quality metric Q in comparison with the quality provided by an ideal adaptive scheme in this conditions, the loss rate and the achieved link utilisation.

CBR periodic - size: 0.5 Mb/s and frequency: 20 s on - 40 s off

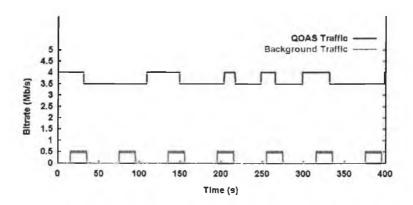


Figure 6-3 QOAS bitrate adaptation versus CBR periodic background traffic with size: 0.5 Mb/s and frequency: 20 s on - 40 s off

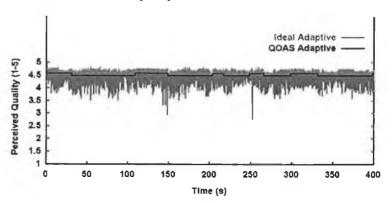


Figure 6-4 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR periodic background traffic with size: 0.5 Mb/s and frequency: 20 s on – 40 s off

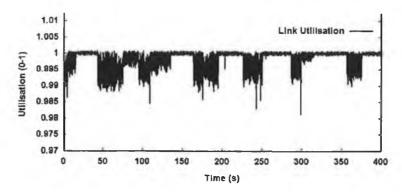


Figure 6-5 Link utilisation for QOAS-based multimedia streaming with CBR periodic background traffic size: 0.5 Mb/s and frequency: 20 s on – 40 s off

CBR periodic - size: 0.5 Mb/s and frequency: 30 s on - 60 s off

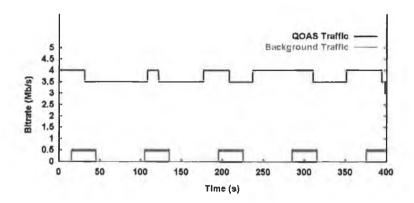


Figure 6-6 QOAS bitrate adaptation versus CBR periodic background traffic with size: 0.5 Mb/s and frequency: 30 s on - 60 s off

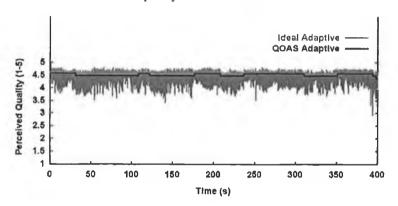


Figure 6-7 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR periodic background traffic with size: 0.5 Mb/s and frequency: 30 s on – 60 s off

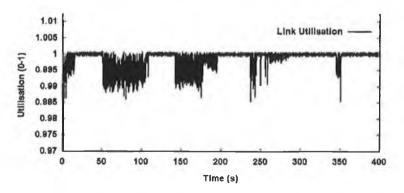


Figure 6-8 Link utilisation for QOAS-based multimedia streaming with CBR periodic background traffic size: 0.5 Mb/s and frequency: 30 s on - 60 s off

CBR periodic - size: 0.5 Mb/s and frequency: 40 s on - 80 s off

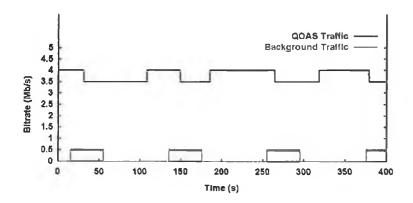


Figure 6-9 QOAS bitrate adaptation versus CBR periodic background traffic with size: 0.5 Mb/s and frequency: 40 s on - 80 s off

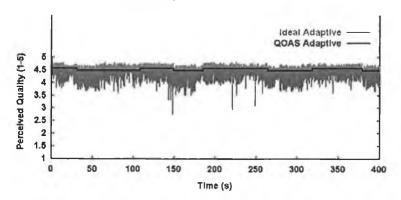


Figure 6-10 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR periodic background traffic with size: 0.5 Mb/s and frequency: 40 s on - 80 s off

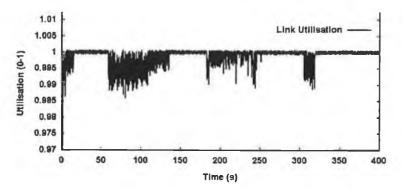


Figure 6-11 Link utilisation for QOAS-based multimedia streaming with CBR periodic background traffic size: 0.5 Mb/s and frequency: 40 s on -80 s off

CBR periodic - size: 0.7 Mb/s and frequency: 20 s on - 40 s off

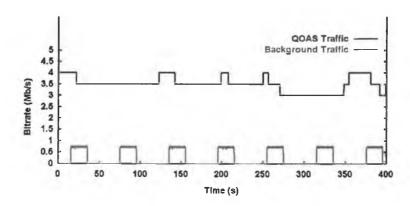


Figure 6-12 QOAS bitrate adaptation versus CBR periodic background traffic with size: 0.7 Mb/s and frequency: 20 s on -40 s off

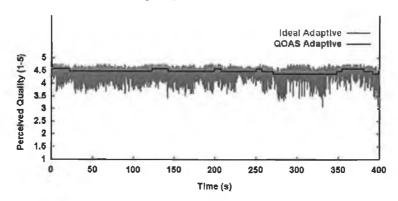


Figure 6-13 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR periodic background traffic with size: 0.7 Mb/s and frequency: 20 s on – 40 s off

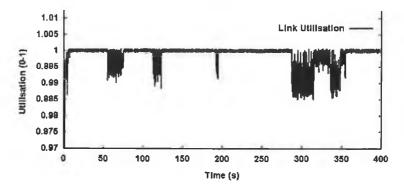


Figure 6-14 Link utilisation for QOAS-based multimedia streaming with CBR periodic background traffic size: 0.7 Mb/s and frequency: 20 s on -40 s off

CBR periodic - size: 0.7 Mb/s and frequency: 30 s on - 60 s off

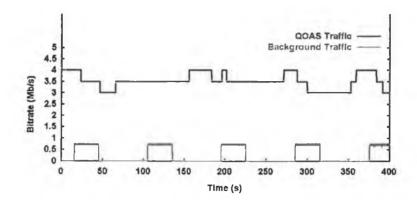


Figure 6-15 QOAS bitrate adaptation versus CBR periodic background traffic with size: 0.7 Mb/s and frequency: 30 s on -60 s off

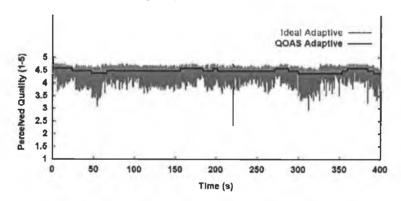


Figure 6-16 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR periodic background traffic with size: 0.7 Mb/s and frequency: 30 s on – 60 s off

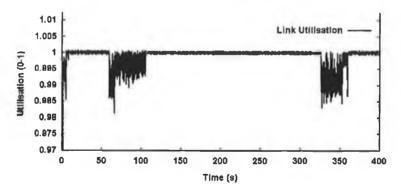


Figure 6-17 Link utilisation for QOAS-based multimedia streaming with CBR periodic background traffic size: 0.7 Mb/s and frequency: 30 s on -60 s off

CBR periodic - size: 0.7 Mb/s and frequency: 40 s on - 80 s off

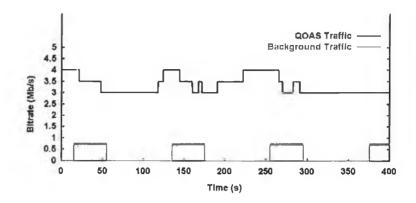


Figure 6-18 QOAS bitrate adaptation versus CBR periodic background traffic with size: 0.7 Mb/s and frequency: 40 s on - 80 s off

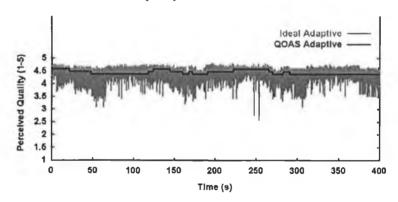


Figure 6-19 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR periodic background traffic with size: 0.7 Mb/s and frequency: 40 s on - 80 s off

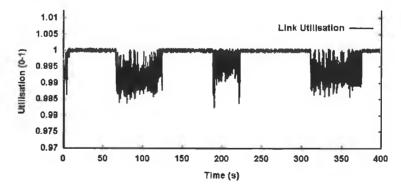


Figure 6-20 Link utilisation for QOAS-based multimedia streaming with CBR periodic background traffic size: 0.7 Mb/s and frequency: 40 s on -80 s off

Comments

In all cases when the CBR background traffic was periodically varied with steps of 0.5 Mb/s, which are comparable with the QOAS adaptation step, regardless of the variation frequency, QOAS has successfully followed the change in the traffic. Figure 6-3, Figure 6-6 and Figure 6-9, which show both the background traffic and the triggered QOAS adaptation, present how the OOAS bit-rate changes in the pre-defined interval of 2-4 Mb/s with almost the same frequency of the background traffic variation. This adaptation is beneficial, completely avoiding packet loss. Consequently the end-user perceived quality achieves very high values, in spite of very high and variable traffic, as shown in Figure 6-4, Figure 6-7 and Figure 6-10. These figures compare the enduser perceived quality achieved by QOAS with the one that may have been achieved by an ideal adaptive scheme in the same conditions. The ideal adaptive scheme is using all the available bandwidth for transmitting multimedia data, achieving therefore 100% link utilisation with no loss and yielding the best end-user quality possible in these conditions. These plots show that the enduser perceived quality when using QOAS tends to the highest values of the one estimated for the ideal adaptive scheme, without having its multiple variations that may disturb the viewers. Also its stand-alone values above 4, the "good" subjective quality level, are impressive, indicating very good QOAS performance from this point of view. At the same time the link utilisation is very close to 100% for the majority of time and even its temporary variations do not lower it below 99%, achieving very good results. Its variation for the duration of these tests is presented in Figure 6-5, Figure 6-8 and Figure 6-11.

The second set of tests involved CBR background traffic variations with the same periodicity, but with steps of 0.7 Mb/s, which are much higher than the QOAS's adaptation step. Similarly Figure 6-12, Figure 6-15 and Figure 6-18 show the QOAS bit-rate adaptive variations triggered by the CBR background traffic, which is also presented in the plots. Although more than one QOAS adaptive step has to be performed, the adaptation is successful and no loss is recorded. In consequence the resulting end-user perceived quality is very high (much above the "good" level) and tends to the highest levels of the one that may have been achieved by an ideal adaptive scheme in the same conditions as shown in Figure 6-13, Figure 6-16 and Figure 6-19. In the same time also the link utilisation values are very close to the 100% as presented in Figure 6-14, Figure 6-17 and Figure 6-20.

More detailed statistics about both the CBR background traffic and the results in these cases are presented in Table 6-14 and Table 6-15.

Table 6-14 Different shapes and variation patters for the tested UDP-CBR periodic background traffic

Traffic Code	Traffic Shape	Size (Mb/s)	Duration (s)	Frequency	Other Traffic Size (Mb/s)
1	Periodic	1 x 0.5	400	20 s on - 40 s off	96.0
2	Periodic	1 x 0.5	400	30 s on - 60 s off	96.0
3	Periodic	1 x 0.5	400	40 s on - 80 s off	96.0
4	Periodic	1 x 0.7	400	20 s on - 40 s off	96.0
5	Periodic	1 x 0.7	400	30 s on - 60 s off	96.0
6	Periodic	1 x 0.7	400	40 s on - 80 s off	96.0

Table 6-15 Statistical results for UDP-CBR periodic background traffic

Traffic Code	QOAS Avg. Bit-rate	Ideal Avg. Bit-rate	QOAS Avg. Perceived Quality (Q)	Ideal Avg. Perceived Quality (Q)	Bandwidth Utilisation (%)	Loss Rate (%)
1	3.670	3.833	4.521	4.548	99.804	0.0
2	3.737	3.848	4.532	4.550	99.840	0.0
3	3.764	3.838	4.537	4.548	99.873	0.0
4	3.496	3.580	4.490	4.505	99.858	0.0
5	3.520	3.581	4.495	4.505	99.876	0.0
6	3.331	3.555	4.458	4.501	99.721	0.0

6.2.3.1.1.2 CBR Staircase

Similar to the CBR periodic variation of the background traffic, the CBR staircase-up and CBR staircase-down variation patterns aim at verifying the QOAS adaptation in very difficult traffic conditions. The background traffic is increased in four steps of 0.4 Mb/s and 0.6 Mb/s respectively and is added to a 95.5 Mb/s CBR background traffic that creates highly loaded network delivery conditions. The QOAS's reaction is then tested when step-wise decreasing the traffic. The consequent QOAS bit-rate adaptations, resulting end-user perceived quality, loss rate variations and link utilisations in these tested cases are presented next.

CBR staircase up and down - steps of size 0.4 Mb/s

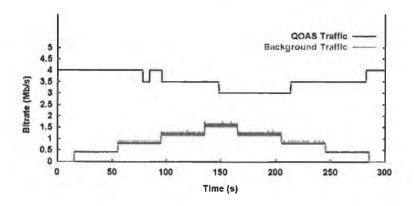


Figure 6-21 QOAS bitrate adaptation vs. CBR staircase background traffic with steps of 0.4 Mb/s

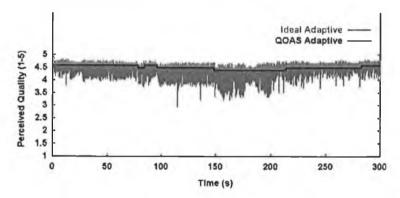


Figure 6-22 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR staircase background traffic with steps of 0.4 Mb/s

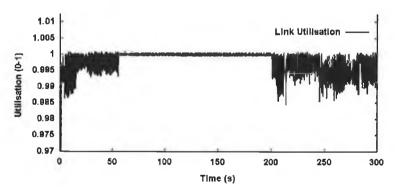


Figure 6-23 Link utilisation for QOAS-based multimedia streaming with CBR staircase background traffic with steps of 0.4 Mb/s

CBR Staircase up and down - steps of size 0.6 Mb/s

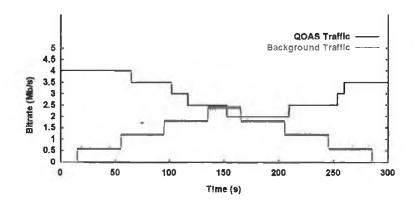


Figure 6-24 QOAS bitrate adaptation vs. CBR staircase background traffic with steps of 0.6 Mb/s

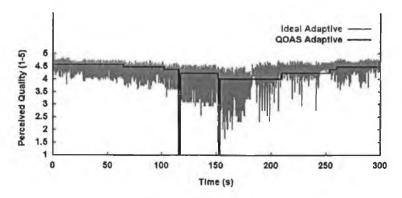


Figure 6-25 End-user perceived quality: QOAS versus ideal adaptive streaming subject to CBR staircase background traffic with steps of 0.6 Mb/s

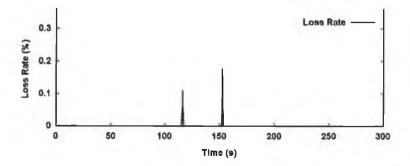


Figure 6-26 Loss rate variation for QOAS-based multimedia streaming with CBR staircase background traffic with steps of 0.6 Mb/s

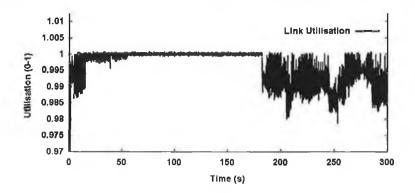


Figure 6-27 Link utilisation for QOAS-based multimedia streaming with CBR staircase background traffic with steps of 0.6 Mb/s

Comments

In these tests the CBR background traffic was varied in a staircase up and staircase down manner, with steps of 0.4 Mb/s and 0.6 Mb/s, lower and respectively higher than the QOAS adaptation step of 0.5 Mb/s. In these conditions the performance of the QOAS's consequent adaptations is assessed.

In the staircase up situations, it is significant to observe that the adaptiveness of the QOAS has successfully and immediately followed the change in the traffic. Figure 6-21 and Figure 6-24 show how the step-by-step increase in the background traffic has triggered staircase-like OOAS adaptations, regardless of the background traffic step size. These adaptations are loss free when the traffic step size is lower than the QOAS adaptation step, but loss occurs for short periods of time when such significant variations in background traffic occur. The duration of these periods of loss is minimised by the QOAS's adaptive reaction, which proves to be very successful. However, as Figure 6-22 and Figure 6-25 show, the end-user perceived quality that is maintained much above the "good" perceptual level for the whole duration of the streaming process in the first case, decreases to the lowest level for these, extremely short, lossy periods. But the real benefit of the QOAS is highlighted when its end-user perceived quality is compared with the end-user perceived quality of a potential ideal adaptive scheme. The latter decreases to the "fair" level for a duration of around 100 s and even dropping to "poor" for more than 20 s, whereas the QOAS's two lossy periods average 1.75 s in duration, as presented in Figure 6-26. Also it is significant to mention that in rest of the time the QOAS's end-user perceived quality is maintained at least at the "good" subjective level. In these tests the link utilisation is very close to 100% for all the time when the background traffic ensures increased traffic conditions, as shown in Figure 6-23 and Figure 6-27.

These results also indicate the QOAS's asymmetric adaptive reaction to network recovery after highly loaded situations. Figure 6-21 and Figure 6-24 show how the QOAS bit-rate adaptations take place a certain period of time after the CBR background traffic has varied. The significant difference between the case when the background traffic step is higher than the adaptation step and when it is lower is that sometimes more than one QOAS adaptive step has to be performed in order to achieve equilibrium. In all the tested cases the adaptation is successful and no loss is recorded due to eventual miss-estimation of the available bandwidth after the background traffic backs off. In consequence the resulting end-user perceived quality is very high (much above the "good" level) at all times and tends to match the levels that may have been achieved by an ideal adaptive scheme in the same conditions as shown in Figure 6-22 and Figure 6-25. The link utilisation, whose values are very close to the 100% when the link is fully loaded, decreases with the decrease in background traffic as presented in Figure 6-23 and Figure 6-27. Figure 6-26 presents the loss rate during streaming and is significant only for the assessment of the staircase-up background traffic variation.

More detailed statistics about the CBR background traffic variation in staircase-like manner are presented in Table 6-16 and Table 6-17. These tables present separately the results related to the period when the traffic varied in a staircase-up manner and put high pressure on QOAS and to the overall testing period.

Table 6-16 Different shapes and variation patters for the tested UDP-CBR staircase background traffic

Traffic Code	Traffic Shape	Size (Mb/s)	Duration (s)	Step Length (s)	Other Traffic Size (Mb/s)
7	Staircase up	4 x 0.4	200	40	95.5
8	Staircase up	4 x 0.6	200	40	95.5
9	Staircase up-down	4 x 0.4	300	40	95.5
10	Staircase up-down	4 x 0.6	300	40	95.5

QOAS Avg. Ideal Avg. QOAS Avg. Ideal Avg. Link Traffic **Loss Rate** Bit-rate **Bitrate** Percy. Quality Percy. Quality Utilisation Code (%) (Mb/s)(Mb/s)(%) (1-5)(1-5)7 3.592 3.617 4.508 4.512 99.905 0.000 8 3.085 99.945 0.091 3.031 4.300 4.391 9 3.568 3.696 4.503 4.525 99.771 0.000 10 3.019 3.296 4.309 4.451 99.628 0.059

Table 6-17 Statistical results for UDP-CBR staircase background traffic

6.2.3.1.2 UDP - VBR as Background Traffic

The majority of multimedia streaming solutions produce very bursty output traffic especially MPEG-encoded streams, due to the different compression achieved for their I, P and B frames. The effect of this kind of background traffic is studied next, taking into consideration different situations in terms of average bit-rate and degree of burstiness.

6.2.3.1.2.1 Constant Average Bit-rate and Variable Burstiness

This section examines the effect of different burstiness associated with the VBR traffic on the QOAS-based adaptation. This background traffic is sent across the bottleneck link along with a 95.5 Mb/s CBR background traffic that simulates a well-multiplexed traffic and creates high loaded network delivery conditions. The characteristics of the VBR background traffic, exponentially generated, are: 0.001 s on - 0.1 s off, 0.01 s on - 0.1 s off and 0.1 s on - 0.1 s off, whereas the bitrate is maintained constant at 1 Mb/s. The QOAS is assessed in terms of its adaptation in relation to the background traffic variation, the estimated end-user perceived quality using the no-reference moving pictures quality metric Q, the loss rate and the link utilisation.

VBR - size: 1.0 Mb/s and burstiness: 0.001 s on - 0.1 s off

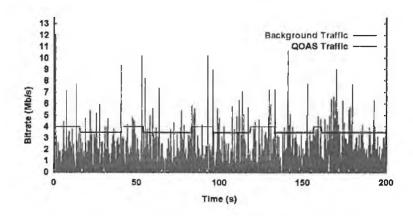


Figure 6-28 QOAS bitrate adaptation versus VBR background traffic with size: 1.0 Mb/s and burstiness: 0.001 s on - 0.1 s off

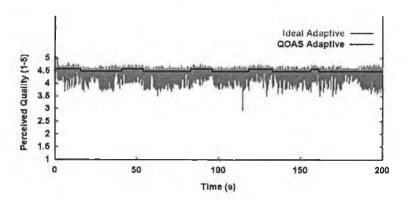


Figure 6-29 End-user perceived quality: QOAS versus ideal adaptive streaming subject to VBR background traffic with size: 1.0 Mb/s and burstiness: 0.001 s on -0.1 s off

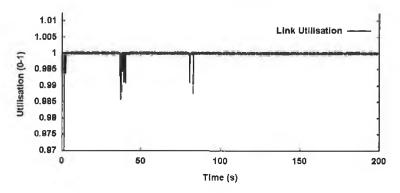


Figure 6-30 Link utilisation for QOAS-based multimedia streaming with VBR background traffic size: 1.0 Mb/s and burstiness: 0.001 s on -0.1 s off

VBR - size: 1.0 Mb/s and burstiness: 0.01 s on - 0.1 s off

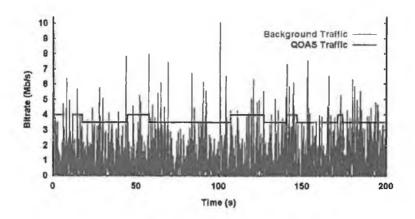


Figure 6-31 QOAS bitrate adaptation versus VBR background traffic with size: 1.0 Mb/s and burstiness: 0.01 s on -0.1 s off

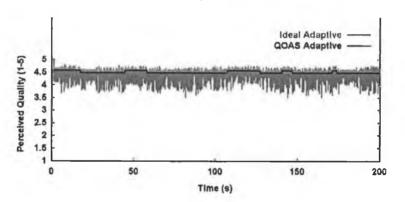


Figure 6-32 End-user perceived quality: QOAS versus ideal adaptive streaming subject to VBR background traffic with size: 1.0 Mb/s and burstiness: 0.01 s on -0.1 s off

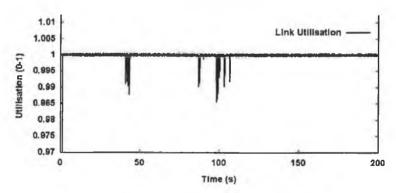


Figure 6-33 Link utilisation for QOAS-based multimedia streaming with VBR background traffic size: 1.0 Mb/s and burstiness: 0.01 s on - 0.1 s off

VBR - size: 1.0 Mb/s and burstiness: 0.1 s on - 0.1 s off

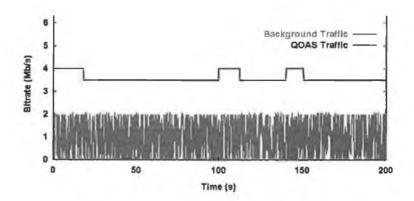


Figure 6-34 QOAS bitrate adaptation versus VBR background traffic with size: 1.0 Mb/s and burstiness: 0.1 s on -0.1 s off

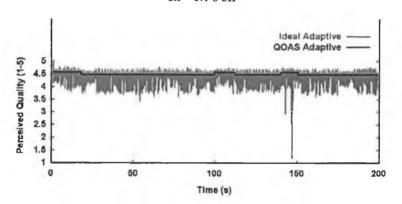


Figure 6-35 End-user perceived quality: QOAS versus ideal adaptive streaming subject to VBR background traffic with size: 1.0 Mb/s and burstiness: 0.1 s on -0.1 s off

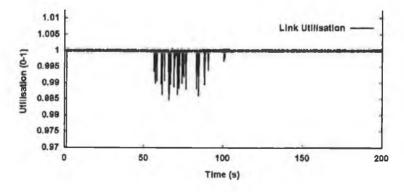


Figure 6-36 Link utilisation for QOAS-based multimedia streaming with VBR background traffic size: 1.0 Mb/s and burstiness: 0.1 s on - 0.1 s off

Comments

Although the average bit-rate of the VBR background traffic is maintained at the same level of 1.0 Mb/s, the burstiness is significantly varied in order to test QOAS's adaptiveness and to assess its performance. Figure 6-28, Figure 6-31 and Figure 6-34 show how QOAS adapts in response to the VBR traffic. This very important result shows that, even with extreme background traffic variation patterns, QOAS adapts successfully in these delivery conditions, achieving no loss. It is also significant to observe that the QOAS's variations are slow, not following the VBR traffic variations and therefore not decreasing much the end-user perceived quality as shown in Figure 6-29, Figure 6-32 and Figure 6-35. Moreover, these figures show that even in comparison to the end-user perceived quality achieved by a hypothetic ideal adaptive scheme that may use all the available bandwidth for multimedia streaming with no loss, QOAS determines very high end-user perceived quality above the "good" subjective level at all times. In these conditions also the link utilisation is very close to the 100 % limit at all times, as presented in Figure 6-30, Figure 6-33 and Figure 6-36. More statistical information is presented in Table 6-18 and Table 6-19.

Table 6-18 Constant average bit-rate and variable burstiness background traffic of type UDP - VBR exponential

Traffic Code	Traffic Shape	Size (Mb/s)	Duration (s)	Traffic characteristic	Other Traffic Size (Mb/s)
I	VBR Exponential	1.0	200	0.001 s on - 0.1 s off	95.5
II	VBR Exponential	1.0	200	0.01 s on - 0.1 s off	95.5
III	VBR Exponential	1.0	200	0.1 s on - 0.1 s off	95.5

Table 6-19 Statistical results for tests with constant average bit-rate and variable burstiness background traffic of type UDP - VBR exponential

Traffic Code	QOAS Avg. Bitrate	Ideal Avg. Bitrate	QOAS Avg. Percv. Quality (Q)	Ideal Avg. Percv. Quality (Q)	Bandwidth Utilisation (%)	Loss Rate (%)
I	3.651	3.658	4.518	4.519	99.942	0.000
II	3.649	3.659	4.517	4.519	99.935	0.000
III	3.601	3.639	4.509	4.516	99.926	0.000

6.2.3.1.2.2 Constant Burstiness and Variable Average Bit-rate

For constant burstiness related to the VBR traffic, chosen as the one that puts the most pressure on the infrastructure, the bit-rate is varied in order to study its effect on the QOAS-based adaptation. This background traffic is sent across the bottleneck link on top of the 95.5 Mb/s CBR background traffic that creates high loaded network delivery conditions. The characteristics of the VBR background traffic, exponentially generated, are: 0.001 s on -0.1 s off and the bit-rates in these tests are 0.8 Mb/s, 1.0 Mb/s and 1.2 Mb/s.

<u>VBR - size</u>: 0.8 Mb/s and burstiness: 0.001 s on -0.1 s off

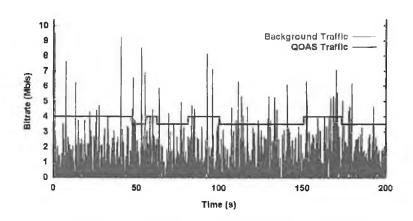


Figure 6-37 QOAS bitrate adaptation versus VBR background traffic with size: 0.8 Mb/s and burstiness: 0.001 s on -0.1 s off

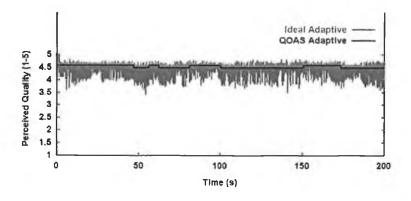


Figure 6-38 End-user perceived quality: QOAS versus ideal adaptive streaming subject to VBR background traffic with size: 0.8 Mb/s and burstiness: 0.001 s on - 0.1 s off

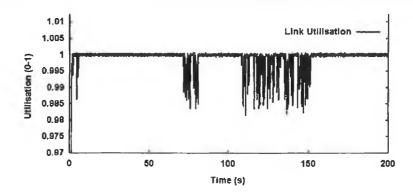


Figure 6-39 Link utilisation for QOAS-based multimedia streaming with VBR background traffic size: 0.8 Mb/s and burstiness: 0.001 s on -0.1 s off

VBR - size: 1.0 Mb/s and burstiness: 0.001 s on - 0.1 s off

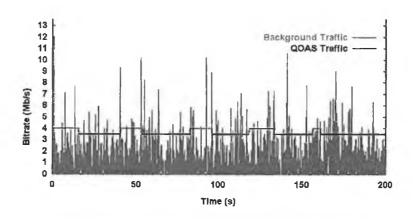


Figure 6-40 QOAS bitrate adaptation versus VBR background traffic with size: 1.0 Mb/s and burstiness: 0.001 s on - 0.1 s off

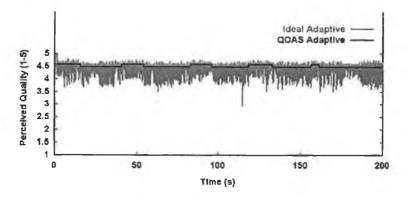


Figure 6-41 End-user perceived quality: QOAS versus ideal adaptive streaming subject to VBR background traffic with size: 1.0 Mb/s and burstiness: 0.001 s on - 0.1 s off

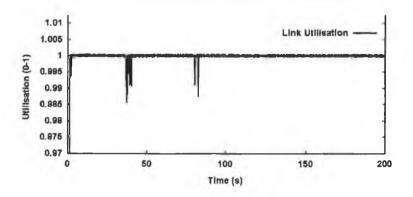


Figure 6-42 Link utilisation for QOAS-based multimedia streaming with VBR background traffic size: 1.0 Mb/s and burstiness: 0.001 s on - 0.1 s off

VBR - size: 1.2 Mb/s and burstiness: 0.001 s on - 0.1 s off

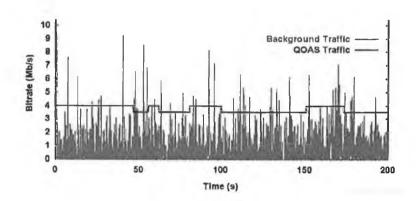


Figure 6-43 QOAS bitrate adaptation versus VBR background traffic with size: 1.2 Mb/s and burstiness: 0.001 s on - 0.1 s off

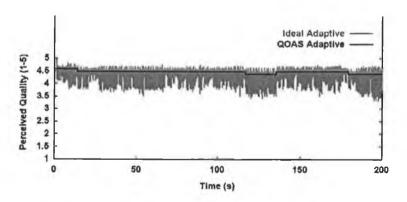


Figure 6-44 End-user perceived quality: QOAS versus ideal adaptive streaming subject to VBR background traffic with size: 1.2 Mb/s and burstiness: 0.001 s on -0.1 s off

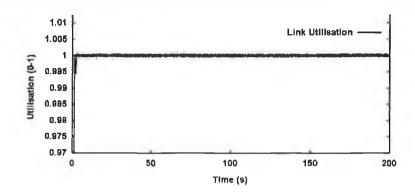


Figure 6-45 Link utilisation for QOAS-based multimedia streaming with VBR background traffic size: 1.2 Mb/s and burstiness: 0.001 s on - 0.1 s off

Comments

Similar to the previous set of tests, QOAS has successfully adapted, regardless of the pressure put on the bottleneck link by increasing the average bit-rate of the VBR background traffic as presented in Figure 6-37, Figure 6-40 and Figure 6-43. However this increase has triggered a decrease in the average bit-rate of the QOAS-transmitted multimedia stream, from 3.74 Mb/s in the first case to 3.65 Mb/s in the second and to 3.43 Mb/s in the third situation. As Figure 6-38, Figure 6-41 and Figure 6-44 show, the end-user perceived quality was at all times very high, much above the "good" perceptual level, although its average has also slightly decreased with the increase in the background traffic from 4.53, to 4.52 and respectively 4.48. Link utilisations are maintained very high for the duration of these tests, achieving averages of roughly 99.9 % (see Figure 6-45).

More statistics related to the UDP - VBR traffic and its effect on the QOAS streaming are presented in Table 6-20 and Table 6-21.

Table 6-20 Constant burstiness and variable average bit-rate background traffic of type UDP - VBR exponential

Traffic Code	Traffic Shape	Size (Mb/s)	Duration (s)	Traffic characteristic	Other Traffic Size (Mb/s)
IV	VBR Exponential	0.8	200	0.001 s on - 0.1 s off	95.5
V	VBR Exponential	1.0	200	0.001 s on - 0.1 s off	95.5
VI	VBR Exponential	1.2	200	0.001 s on - 0.1 s off	95.5

Table 6-21 Statistical results for tests with constant burstiness and variable average bit-rate background traffic of type UDP - VBR exponential

Traffic Code	QOAS Avg. Bitrate	Ideal Avg. Bitrate	QOAS Avg. Percv. Quality (Q)	Ideal Avg. Percv. Quality (Q)	Bandwidth Utilisation (%)	Loss Rate (%)
IV	3.736	3.824	4.532	4.546	99.849	0.000
V	3.651	3.658	4.518	4.519	99.942	0.000
VI	3.433	3.445	4.478	4.481	99.950	0.000

6.2.3.1.3 TCP as Background Traffic

The very large majority of Internet traffic today consists of file transfers that use TCP as the transport protocol and among the most popular applications that have based their functionality on TCP are FTP applications employed for file transfers and WWW applications used for immediate viewing of the content. These applications were chosen because they are representative for two types of TCP-based traffic: **long-lived** and **short-lived**. The former is characterised by long duration processes that produce in general slow-changing traffic, whereas the latter is responsible for highly variable traffic, of short durations and therefore very bursty. The effect of these traffic types is studied in this section, taking into consideration different sizes that affect differently the network delivery conditions and therefore the QOAS-based multimedia streaming.

6.2.3.1.3.1 Long-lived TCP

Two sets of tests are performed that aim at transmitting 50 and 54 FTP flows, generated using the NS-2 built-in models, that account for a significant long-lived TCP background traffic on top a 75 Mb/s CBR background traffic that represents a well-multiplexed traffic and aims at creating high loaded network delivery conditions. Next the results of the QOAS-based adaptive multimedia streaming sent along with this long-lived traffic through the bottleneck link are presented and are assessed in terms of adaptiveness in relation to the background traffic variation, estimated end-user perceived quality in comparison with an ideal adaptive scheme that would perform in these conditions, loss rate and link utilisation.

Long-lived TCP - 50 FTP flows

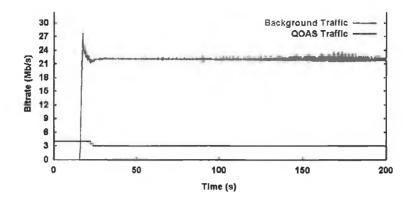


Figure 6-46 QOAS bit-rate adaptation versus 50 FTP flows as background traffic

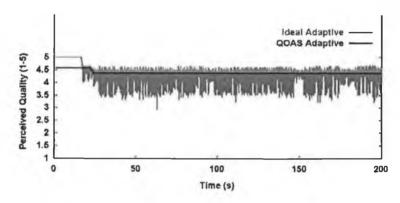


Figure 6-47 End-user perceived quality: QOAS versus ideal adaptive streaming subject to 50 FTP flows as background traffic

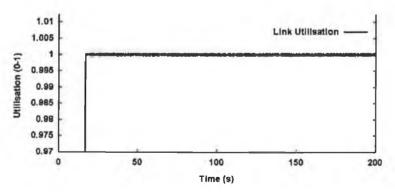


Figure 6-48 Link utilisation for QOAS-based multimedia streaming with 50 FTP flows as background traffic

Long-lived TCP - 54 FTP flows

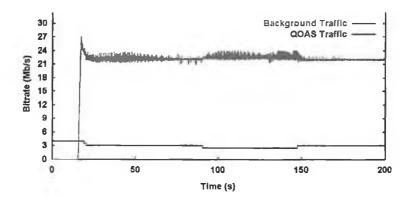


Figure 6-49 QOAS bitrate adaptation versus 54 FTP flows as background traffic

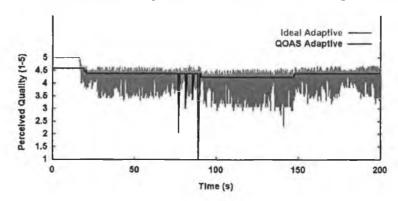


Figure 6-50 End-user perceived quality: QOAS versus ideal adaptive streaming subject to 54 FTP flows as background traffic

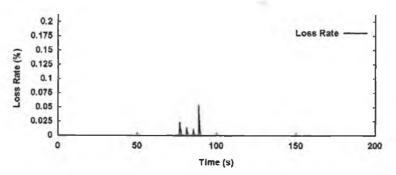


Figure 6-51 Loss rate variation when QOAS-based multimedia streaming with 54 FTP flows as background traffic

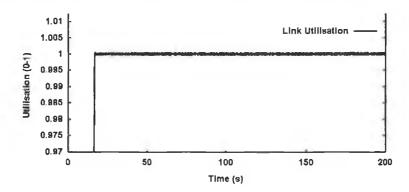


Figure 6-52 Link utilisation when streaming multimedia using QOAS with 54 FTP flows as background traffic

Comments

In these cases when long-lived TCP - itself based on an adaptive mechanism - was used as background traffic, the QOAS has adapted being both aggressive and TCP friendly. On one side it is significant to achieve the highest possible end-user quality by sending as much multimedia data in as timely manner as possible in the given conditions, regardless of the other sources of traffic. On the other hand it is also important not to undermine the other provided services by using all the available bandwidth. An approach that balances these directions is taken by QOAS that adapts to a certain level even in the presence of other services based on adaptive control schemes that also adapt. Figure 6-46 and Figure 6-49 include also the transitory period showing how both the longlived TCP and the QOAS gracefully adapt by sharing the available bandwidth. In relation to the TCP traffic, this behaviour is known as TCP-friendliness, as already mentioned in the second chapter. Due to this adaptation that completely avoids packet loss in the first case QOAS achieves very high end-user perceived quality as shown in Figure 6-47, reaching an average of 4.39, which is very close to the 4.42 computed for an ideal adaptive scheme in the same conditions (see Table 6-22 and Table 6-23). However, when 54 FTP flows are transmitted, increasing much the load on the bottleneck link, before it succeeds to adapt, QOAS experiences some loss that decreases temporarily its end-user perceived quality as shown in Figure 6-50 and Figure 6-51. But the average loss duration is 0.8 s for the QOAS streaming and these periods are the only ones when the end-user perceived quality decreases below the "good" subjective level, as compared to the ideal adaptive scheme that may have maintained the "fair" level for a duration of 100 s. The link utilisation is very close to 100 % at all times in both these tests as shown in Figure 6-48 and Figure 6-52. More statistics-related details are presented next in Table 6-22 and Table 6-23.

Table 6-22 Characteristics of the long-lived TCP background traffic

Traffic Code	Traffic Shape	Avg. Traffic Size (Mb/s)	Duration (s)	Other Traffic Size (Mb/s)	
a	50 x FTP	22.0	200	75.0	
b	54 x FTP	22.5	200	75.0	

Table 6-23 Statistical results for tests with long-lived TCP background traffic

Traffic Code	QOAS Avg. Bitrate (Mb/s)	Ideal Avg. Bitrate (Mb/s)	QOAS Avg. Perceived Quality (Q)	Ideal Avg. Perceived Quality (Q)	Bandwidth Utilisation (%)	Loss Rate (%)
a	3.042	3.140	4.394	4.417	98.423	0.000
b	2.781	2.729	4.291	4.309	98.425	0.036

6.2.3.1.3.2 Short-lived TCP

In this section during two sets of tests using 40 and 50 WWW sessions are generated using the NS-2 built-in models that account for short-lived TCP background traffic. This traffic is on top a 95.5 Mb/s CBR background traffic that stands for well-multiplexed different type traffic and aims at creating high loaded network delivery conditions. The traffic that corresponds to the WWW sessions is generated using the following characteristics, considered typical for a WWW session by the research in the WWW area [251, 252]:

- the inter-session time was exponentially distributed with an average of 2 s,
- the number of WWW pages retrieved during a session was constant and equal to 5,
- the retrieval time between consecutive pages was exponentially distributed with an average of 2 s,
- the number of objects within a page was considered constant and equal to 10,
- the time between two consecutive requests for objects belonging to the same page was considered exponential distributed with an average of 0.01 s,

• the size of the objects has followed a Pareto distribution with an average 10 KB and shape equal to 1.2.

Next the results of the QOAS-based adaptation, when streaming multimedia in these conditions, are presented and the achieved performance is highlighted.

Short-lived TCP - 40 WWW sessions

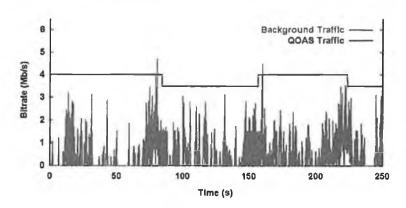


Figure 6-53 QOAS bit-rate adaptation versus 40 WWW sessions as background traffic

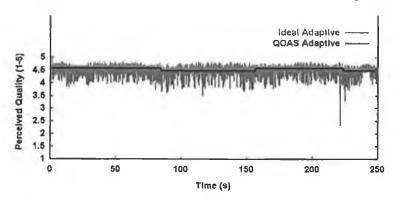


Figure 6-54 End-user perceived quality: QOAS versus ideal adaptive streaming subject to 40 WWW sessions as background traffic

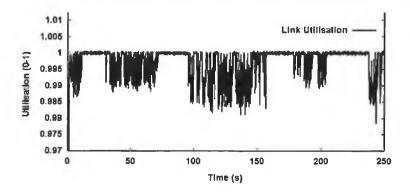


Figure 6-55 Link utilisation for QOAS-based multimedia streaming with 40 WWW sessions as background traffic

Short-lived TCP - 50 WWW sessions

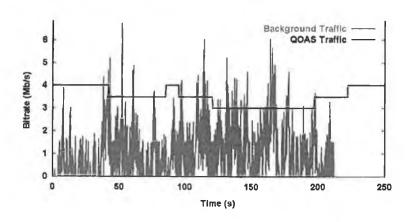


Figure 6-56 QOAS bitrate adaptation versus 50 WWW sessions as background traffic

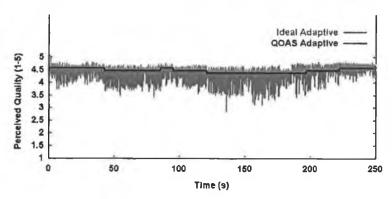


Figure 6-57 End-user perceived quality: QOAS versus ideal adaptive streaming subject to 50 WWW sessions as background traffic

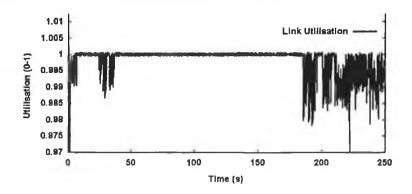


Figure 6-58 Link utilisation for streaming multimedia using QOAS with 50 WWW sessions as background traffic

Comments

In these tests the QOAS has achieved very good adaptation as shown in Figure 6-53 and Figure 6-56, not experiencing any loss in such bursty delivery conditions generated by the WWW background traffic. Consequently also the end-user perceived quality was very high, between the "good" and the "excellent" subjective levels for all the duration of these tests, achieving an average of 4.54 and respectively 5.49. The comparison to the quality achieved by a hypothetical ideal adaptive scheme is presented for each case, in Figure 6-54 and Figure 6-57 respectively. Link utilisation, although highly variable due to the burstiness of this background traffic was on average within 0.3 % from the maximum 100 %, which suggests very good performance for QOAS.

More detailed statistics about tests that have involved TCP-based background traffic are resented in Table 6-24 and Table 6-25.

Table 6-24 Characteristics of the TCP background traffic

Traffic Code	Traffic Shape	affic Shape Avg. Traffic Size (Mb/s)		Other Traffic Size (Mb/s)	
e	40 x WWW	0.48	250	95.5	
d	50 x WWW	0.91	250	95.5	

Traffic Code	QOAS Avg. Bitrate (Mb/s)	Ideal Avg. Bitrate (Mb/s)	QOAS Avg. Perceived Quality (Q)	Ideal Avg. Perceived Quality (Q)	Bandwidth Utilisation (%)	Loss Rate (%)
c	3.802	4.016	4.543	4.575	99.690	0.000
d	3.505	3.587	4.492	4.507	99.803	0.000

Table 6-25 Statistical results for tests with TCP background traffic

6.2.3.2 Comparison to an Ideal Adaptive Scheme

One of the most difficult challenges in building any adaptive control solution and especially when it targets multimedia streaming is to successfully adapt to changing network conditions and achieve an output rate that matches the available bandwidth for the transmission. The best assessment of the performance of any adaptive scheme should be the comparison with a hypothetic ideal adaptive scheme that uses all the available bandwidth for transmitting data achieves 0 % loss and 100 % link utilisation. This section presents this comparison involving the QOAS when streaming multimedia over the bottleneck link in the presence of background traffic with different types, shapes and variation patterns.

The tests presented in the previous section have tested the effect of different shaped UDP-CBR traffic on the QOAS such as *periodic*, with different periodicity, *staircase up* and respectively *staircase down*, each with different step sizes. UDP-VBR traffic was generated using an exponential distribution with on/off patterns that were varied from 0.001s/0.1s to 0.1s/0.1s and with different sizes from 0.8 Mb/s to 1.2 Mb/s. Different types of TCP connections were considered such as *long-lived TCP* (e.g. FTP flows) and *short-lived TCP* (e.g. WWW sessions) and in different number. The levels of the background traffic were such chosen in order to trigger adaptive variations of the streamed clips and to have different sizes relative to the adaptation step of 0.5 Mb/s. The simulations lasted on average 300 s and the first and the last transitory 50 s are not included in the results reported in the tables. During these simulations the potential behaviour of the ideal adaptive scheme in the same conditions is analysed and used as a base for comparison with the QOAS-related results.

Table 6-26 Background traffic of different types, shapes and sizes when testing QOAS

	Background Traffic							
Traffic Code	Traffic Type	Traffic Shape	Size (Mb/s)	Other Traffic Size (Mb/s)				
A	CBR UDP	Periodic - 40 s on 80 s off	1 x 0.5	96.0				
В	CBR UDP	Periodic - 40 s on 80 s off	1 x 0.7	96.0				
С	CBR UDP	Staircase up - 40 s step length	4 x 0.4	95.5				
D	CBR UDP	Staircase up - 40 s step length	4 x 0.6	95.5				
E	CBR UDP	Staircase down - 40 s step length	4 x 0.4	95.5				
F	CBR UDP	Staircase down - 40 s step length	4 x 0.6	95.5				
G	VBR UDP	0.001s on 0.1s off	1 x 0.8	95.5				
Н	VBR UDP	0.001s on 0.1s off	1 x 1.2	95.5				
I	TCP	FTP	50 x 0.44	75.0				
J	TCP	FTP	54 x 0.42	75.0				
K	TCP	www	40 x 0.012	95.5				
L	ТСР	www	50 x 0.018	95.5				

The comparative test results are presented in Table 6-26 and Table 6-27. The reported results represent computed average values for the duration of the tests. It is significant to observe that during these tests, regardless of the background traffic type, shape and size, the QOAS-based system scored highly in terms of perceived quality, loss rate and bottleneck link utilisation even in comparison with an ideal system, which it very unlikely to be ever built. The adaptation was so successful that the QOAS streaming has maintained loss rates of less than 0.1% in all cases, although the delivery network was fully loaded. The perceived quality scores are exceptional, not only that they are above the "good" perceptual level (4.00 on the 1-5 scale), but also in almost all cases they are within 1% from the ideal and in only one case is 3% adrift. The bottleneck link utilisation also reaches very high levels, QOAS making use of more than 99.5% of the bandwidth resources in the large majority of tests and even in the two remaining cases the available resources are less than 1.5% from being fully used. These results indicate a highly performant behaviour of the QOAS and shows good adaptations regardless of the background traffic type, size and shape.

Table 6-27 Comparison between QOAS and ideal streaming subject to concurrent traffic

Traff. Code	QOAS Avg. Rate (Mb/s)	Ideal Avg. Rate (Mb/s)	QOAS Quality (1-5)	Ideal Quality (1-5)	QOAS Loss Rate (%)	Ideal Loss Rate (%)	QOAS Utilis. (%)	Ideal Utilis. (%)
A	3.76	3.84	4.54	4.55	0.0	0.0	99.87	100.00
В	3.33	3.55	4.46	4.50	0.0	0.0	99.72	100.00
С	3.59	3.62	4.51	4.52	0.0	0.0	99.90	100.00
D	3.03	3.09	4.31	4.39	0.09	0.0	99.95	100.00
E	3.57	3.70	4.50	4.53	0.0	0.0	99.77	100.00
F	3.02	3.30	4.31	4.45	0.006	0.0	99.63	100.00
G	3.74	3.82	4.53	4.55	0.0	0.0	99.85	100.00
Н	3.43	3.45	4.48	4.49	0.0	0.0	99.85	100.00
I	3.04	3.14	4.39	4.42	0.0	0.0	98.42	100.00
J	2.73	2.78	4.29	4.31	0.04	0.0	98.43	100.00
K	3.80	4.02	4.54	4.58	0.0	0.0	99.69	100.00
L	3.50	3.59	4.49	4.51	0.0	0.0	99.80	100.00

6.2.3.3 Single QOAS-based Streaming Against Multimedia Traffic

6.2.3.3.1 Overview

This set of tests aims at assessing the performance of the delivery of a single multimedia stream using QOAS in increased traffic conditions and in the presence of other multimedia streams. Since the QOAS is designed for the local broadband multi-service IP networks in which the majority of traffic is expected to be multimedia-related, this section analyses in detail how different concurrent multimedia streams affect the QOAS-controlled adaptive stream. The effect on the QOAS-based adaptive streaming of repeated VCR-like operations such as play, pause and stop that involve these concurrent streaming processes is especially studied. The QOAS-related performances are then compared to other streaming solutions such as the adaptive schemes LDA+ [7] and TFRCP [6], already presented in the second chapter and the non-adaptive (NoAd) solution.

For these objective tests NS-2 is used and the "Dumbbell" topology, which was presented in section 6.2.1.2. In order to generate CBR-UDP background traffic, NS-2's CBR-traffic model is used. The QOAS model was presented in section 6.2.1.3 and the LDA+, TFRCP and NoAd models are described in the next section. From the multimedia clips presented in section 6.2.1.4 *diehard1* was chosen since it has the highest motion content and it may be affected the most by background traffic variations. Details about its five pre-recorded different quality versions are presented in Table 6-1. As mentioned in section 6.2.1.5 the performance is assessed in terms of schemes' adaptiveness to background traffic, resulted end-user perceived quality, loss rate and link utilisation.

6.2.3.3.2 NoAd, TFRCP and LDA+ Models

The **NoAd model** implements the non-adaptive multimedia streaming approach which transmits multimedia data using the highest available rate, regardless of the background traffic or eventual other problems that may affect the delivery process (e.g. loss, increased delays etc.). In order to allow for a fair comparison to the QOAS model, during testing the NoAd implementation streams the multimedia clips at their maximum rate of 4 Mb/s, which is the maximum available also for QOAS-based adaptations.

The TFRCP model relies on the TCP-Friendly Rate Control Protocol (TFRCP), an equation-based TCP-friendly adaptation scheme proposed in [6]. The adaptive scheme uses estimates of the round-trip delay and loss rates for the latest transmission round i to determine the adaptation policy for the next round i+1.

In the case of zero loss in the previous interval ($p_i=0$), the current transmission rate is doubled as shown in equation (6-2).

$$r_{i+1} = 2 * r_i {(6-2)}$$

In case of a non-zero loss rate, TFRCP restricts the transmission rate to the equivalent of a TCP flow transmitted in the same conditions, as computed by the TCP model proposed in [101]. Equation (6-3) and (6-4) present the formula according to which the transmission rate is computed in which W_{max} is the receiver's window size, R is the round trip time, p_i the loss rate in round i and B the TCP base timeout value.

$$r_{i+1} = f(W_{\text{max}}, p_i, R, B)$$
 (6-3)

$$f(W_{\text{max}}, p, R, B) = \begin{cases} \frac{\frac{1-p}{p} + W(p) + \frac{Q(p, W(p))}{1-p}}{R(W(p) + 1) + \frac{Q(p, W(p))G(p)B}{1-p}} & \text{for } W(p) < W_{\text{max}} \\ \frac{\frac{1-p}{p} + W_{\text{max}} + \frac{Q(p, W_{\text{max}})}{1-p}}{R(\frac{W_{\text{max}}}{4} + \frac{1-p}{pW_{\text{max}}} + 2) + \frac{Q(p, W_{\text{max}})G(p)B}{1-p}} & \text{otherwise} \end{cases}$$
(6-4)

where W(p), Q(p, w) and G(p) are computed as in equations (6-5), (6-6) and (6-7).

$$W(p) = \frac{2}{3} + \sqrt{\frac{4(1-p)}{3p} + \frac{4}{9}}$$
 (6-5)

$$Q(p,w) = \min\left(1, \frac{(1-(1-p)^3)(1+(1-p)^3(1-(1-p)^{w-3}))}{1-(1-p)^w}\right)$$
(6-6)

$$G(p) = 1 + p + 2p^{2} + 4p^{3} + 8p^{4} + 16p^{5} + 32p^{6}$$
(6-7)

The sender can update its rate in intervals of 2 to 5 s. The implemented TFRCP model uses 5 s long rate update intervals as suggested in [6] for link delays greater than 100 ms as in our topology. The transmission rates for the multimedia streams are maintained between 2 Mb/s and 4 Mb/s allowing for comparison to be performed with QOAS in similar conditions.

The LDA+ model is based on the Loss-Delay-based Adaptation Algorithm (LDA+), which was presented in [7]. LDA+ is an AIMD algorithm that changes its transmission rate after each receiver report according to the estimation of the network situation and of the share of the bandwidth already used. Therefore the scheme works in rounds between such receiver reports making the rate for round i to be based on the reports about the delivery performance in the previous round i-1.

In loss situations, the rate is decreased by the factor $1-p^{1/2}$, p being the loss rate, with final value not lower than the one suggested by the TCP model proposed in [101]. The transmission rate for round i: r_i is computed as in equation (6-8), based on the TCP rate formula presented in equation (6-9).

$$r_i = \max(r_{i-1} * (1 - \sqrt{p}), r_{TCP})$$
 (6-8)

$$r_{TCP} = \frac{S}{t_{RTT} \sqrt{\frac{2Dp}{3} + t_{out} \min\left(1, 3\sqrt{\frac{3Dp}{8}}\right) p(1 + 32 p^2)}}$$
(6-9)

where S is the packet size, t_{RTT} is the round trip delay, t_{out} is the TCP retransmission timeout, D is the number of acknowledged TCP packets by a single ACK packet and p is the loss fraction.

In cases with no loss, the additive value A_i for the rate is computed as the minimum between three values. The first value Aadd_i is computed in inverse relation with the share of the bandwidth that the current flow utilises. A second value $Aexp_i$ is meant to limit the increase to the bottleneck link bandwidth as it converges to 0 when this happens. The third value A_{TCPi} is determined in such a manner that, at no time, the rate should increase faster than a TCP flow sharing the same link. The formulas according to which the rate is computed are presented in the equations (6-10) - (6-14).

$$r_i = r_{i-1} + A_i \tag{6-10}$$

$$A_i = \min(A_{add_i}, A_{\exp_i}, A_{TCP_i})$$
(6-11)

$$A_{add_i} = \left(2 - \frac{r_{m-1}}{Bw}\right) \times A_{i-1} \tag{6-12}$$

$$A_{\exp_i} = \left(1 - \exp^{-\left(1 - \frac{r_{i-1}}{Bw}\right)}\right) \times r_{i-1}$$
(6-13)

$$A_{TCP_i} = \frac{N}{T} \rightarrow \frac{\frac{T}{R} + 1}{2 \times R}$$
, with $N = \sum_{n=0}^{T/R} n = \frac{\left(\frac{T}{R} + 1\right) \times \frac{T}{R}}{2}$ (6-14)

where Bw is the bottleneck bandwidth, N the number of packets TCP would increase its window with, T the interval between two receiver reports and R the round trip delay.

The LDA+ model's implementation uses a receiver report feedback interval of 5 s as suggested in [7] to minimise the quality variations and achieve better performances.

6.2.3.3.3 Background Traffic

Since QOAS was designed especially for highly loaded network conditions, CBR-UDP background traffic with a rate of 95.5 Mb/s is generated using the NS-2's model for the CBR traffic. This traffic represents a well-multiplexed real-life traffic composed of a high number of individual data flows of different types, shapes and variation patterns, as expected in a local multi-service broadband IP-network. On top of this traffic a complex multimedia-like traffic, presented in Figure 6-59, is transmitted. This traffic simulates all possible effects of user interactions to multimedia streams such as repeated *play*, *pause*, re-*play* and *stop*. It even takes into account the effect of multiple consecutive *play* commands that increase the traffic in a staircase up manner, consecutive *pause-play* interactions with different frequency and applied on movies with different rate and consecutive stop-s that decrease the traffic in a staircase down fashion. QOAS is tested with this traffic and its performances are compared with the ones obtained by using LDA+, TFRCP and NoAd.

6.2.3.3.4 Testing QOAS 2.5 Background Traffic 2 Bitrate (Mb/s) 1.5 1 0.5 0 0 50 100 150 200 250 300 350 400 450

Figure 6-59 Background traffic variation on top of 95.5 Mb/s CBR traffic

Time (s)

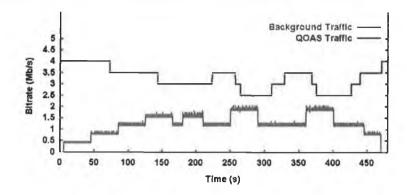


Figure 6-60 QOAS bit-rate adaptation versus complex multimedia traffic

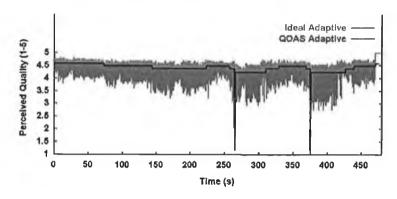


Figure 6-61 End-user perceived quality: QOAS versus ideal adaptive streaming subject to complex multimedia background traffic

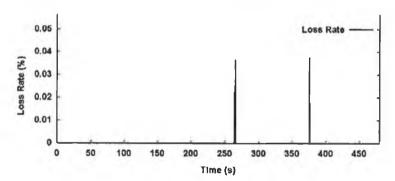


Figure 6-62 Loss rate variation when QOAS-based multimedia streaming with complex multimedia as background traffic

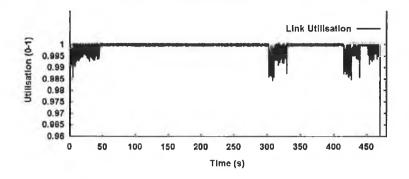


Figure 6-63 Link utilisation when QOAS-based multimedia streaming with complex multimedia as background traffic

Comments

QOAS was successfully tested with CBR background traffic of different shapes and variation patterns and, as expected, achieves very good performance also when used for streaming against a more complex multimedia-like traffic sequence. Figure 6-60 presents the adaptiveness of the QOAS with the background traffic variations. Apart from the staircase-up and staircase-down traffic variations that trigger good adaptations from the QOAS and were already discussed, it is significant to mention that QOAS's asymmetric behaviour related to upgrades and downgrades in the streaming rate pays off after 165 s. At this moment a 10 s pause in streaming of a multimedia sequence does not determine QOAS adaptation and the quality of the streamed multimedia does not change. This is unlike what happens after 320 s when the pause between the play commands is long enough to trigger quality adaptations. Unfortunately the significant size of the step with which the rate of the background traffic increases determines temporary losses presented in Figure 6-62. These losses affect the end-user perceived quality for short moments of time when they occur (average loss duration is 1.45 s), but the QOAS succeeds to adapt, restoring quickly the original quality as shown in Figure 6-61. In spite of this decrease in quality, the average end-user perceived quality has scored on average 4.38, between the "good" and the "excellent" subjective quality levels. Figure 6-63 shows the link utilisation variation during this test. It is important to note in relation to this link utilisation that its average was 99.93 %.

6.2.3.3.5 Testing TFRCP

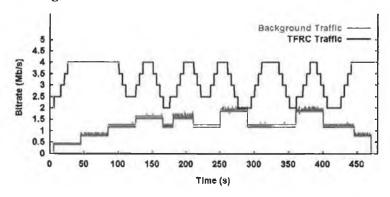


Figure 6-64 TFRCP bit-rate adaptation versus complex multimedia traffic

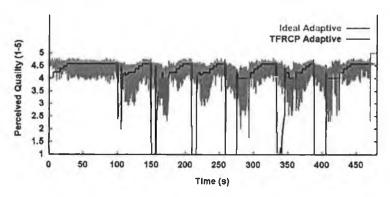


Figure 6-65 End-user perceived quality: TFRCP versus ideal adaptive streaming subject to complex multimedia background traffic

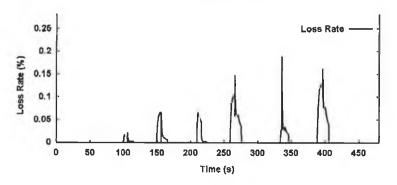


Figure 6-66 Loss rate variation when TFRCP-based multimedia streaming with complex multimedia as background traffic

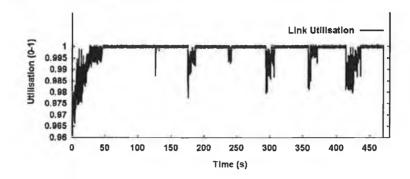
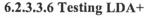


Figure 6-67 Link utilisation when TFRCP-based multimedia streaming with complex multimedia as background traffic

Comments

TFRCP mainly bases its adaptation on the loss rate in a similar manner with TCP, backing off when loss occurs and step-wise increasing its transmission rate in case of no loss. This very variable behaviour, acknowledged by the authors in [6], determines the multimedia clip's streaming rate to vary much between the minimum and maximum rate limits as shown in Figure 6-64. Since the scheme does not prevent the loss from happening, adapting only when loss occurs (see Figure 6-66), the end-user perceived quality is severely affected for periods that exceed 10 s in length, as shown in Figure 6-65. Although high, the link utilisation does not achieve the performance obtained when using QOAS for streaming and this is mainly due to the TFRCP behaviour that reduces its rate to a value below the available bandwidth as soon as loss is experienced. Figure 6-67 shows the link utilisation variation during TFRCP streaming.



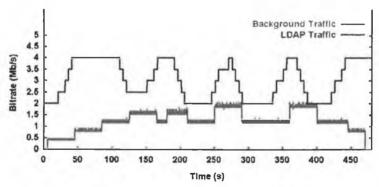


Figure 6-68 LDA+ bit-rate adaptation versus complex multimedia traffic

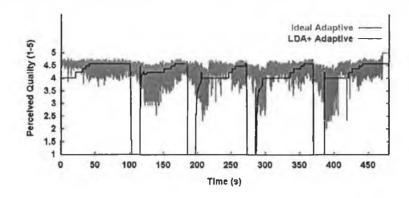


Figure 6-69 End-user perceived quality: LDA+ versus ideal adaptive streaming subject to complex multimedia background traffic

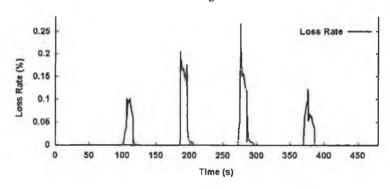


Figure 6-70 Loss rate variation when LDA+-based multimedia streaming with complex multimedia as background traffic

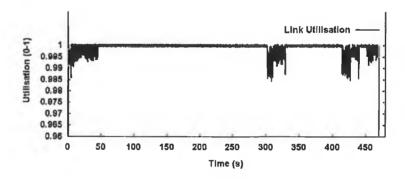


Figure 6-71 Link utilisation when LDA+-based multimedia streaming with complex multimedia as background traffic

Comments

Although based on a more complex algorithm, LDA+ also fails to adapt successfully to the available bandwidth when the background traffic significantly varies as in Figure 6-59. Adapting

the transmission rate in response to loss, the LDA+-related bit-rate bounces much between the minimum and the maximum limits as shown in Figure 6-68. The consequent end-user perceived quality not only that varies with these bit-rate changes, but it is also severely affected by loss for long periods of time, as shown in Figure 6-69. The loss rate variations that cause these effects on the end-user perceived quality during the LDA+ streaming process are presented in Figure 6-70. The achieved link utilisation is very high 99.67 %, but lower than the QOAS's, and varies as shown in Figure 6-71

6.2.3.3.7 Testing NoAd

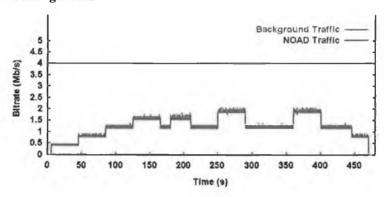


Figure 6-72 NoAd bit-rate versus complex multimedia traffic

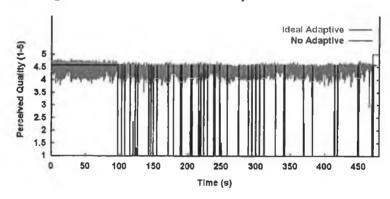


Figure 6-73 End-user perceived quality: NoAd versus ideal adaptive streaming subject to complex multimedia background traffic

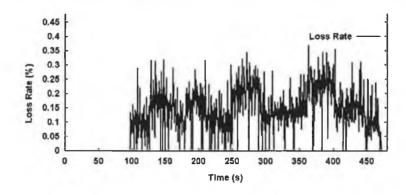


Figure 6-74 Loss rate variation when NoAd-based multimedia streaming with complex multimedia as background traffic

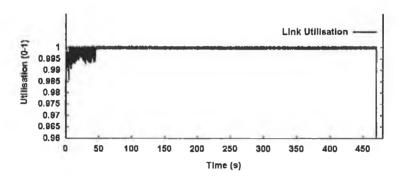


Figure 6-75 Link utilisation when NoAd-based multimedia streaming with complex multimedia as background traffic

Comments

NoAd can successfully stream multimedia data in normal traffic conditions at a flat average rate, regardless of the variation in delivery conditions, as shown in Figure 6-72 Unfortunately, as soon as the background traffic increases much, significant loss occurs and severely affects the enduser perceived quality that drops to the "bad" perceptual level for almost the whole duration of the streaming process. The loss rate variation is presented in Figure 6-74 and the consequent end-user perceived quality is shown in Figure 6-73. Although the link utilisation reaches 100%, the very low end-user perceived quality does not support NoAd as an acceptable solution for streaming multimedia.

6.2.3.4 Single QOAS - Comparison to Other Streaming Solutions

The QOAS solution was used for streaming the high motion content multimedia clip diehard1 in highly variable multimedia-like background traffic. In identical conditions, TFRCP and

LDA+-based adaptive solutions and a non-adaptive mechanism were used for streaming the same clip and the results were presented in terms of adaptiveness to background traffic, end-user perceived quality, loss rate and link utilisation in the previous section. The QOAS's performance is compared next with the performances of the other streaming solutions.

In normal traffic conditions, when the delivery network is not heavily loaded, all four schemes perform similarly by transmitting maximum quality data. After 85 s the staircase-like background traffic exceeds the available bandwidth determining adaptive reactions from QOAS, LDA+ and TFRCP-based solutions and causing losses in non-adaptive, LDA+ and TFRCP cases. The QOAS reacts faster than TFRCP and LDA+, reducing the quantity of the transmitted data and successfully avoids losses that occur in the other two cases, significantly degrading their end-user perceived quality. However, TFRCP reacts faster and minimises the lossy period in comparison to the LDA+. For the non-adaptive streaming, starting from this moment, the corresponding user-perceived quality is extremely poor for the whole duration when network conditions are highly loaded.

QOAS's conservative behaviour that maintains the current transmission state unless there is a significant change in the delivery conditions in comparison to both the TFRCP and LDA+ that tend to aim for a higher rate until loss occurs, pays off for example at 125 s. At this moment the background traffic further increases with 0.4 Mb/s and QOAS successfully adapts avoiding losses, whereas both TFRCP and the LDA+ experience significant losses, severely degrading the perceived quality. The duration of the period with low perceived quality is short in the TFRCP case since the stream finally adapts to the available bandwidth.

The asymmetric reaction to events prevents the QOAS adaptive system from immediately responding to the decrease in background traffic that occurs at 165 s during the 10 s-long brief pause. Therefore when the traffic increases again at 175 s, the stream neither experiences losses, nor has to adapt, maintaining a stable user-perceived quality. LDA+ also responds with certain latency to improvements in the delivery conditions and reacts fast to negative changes in the network traffic. This is the cause for its successful reaction to short breaks in streaming of concurrent multimedia streams, not experiencing losses. Unfortunately this was not the case for TFRCP whose associated end-user perceived quality decreases again to the "bad" level for certain period of time as a direct consequence of loss.

When the decrease in background traffic is prolonged as it is in the case of the longer pause starting at 290 s, although all adaptive schemes correctly determine that the congestion has passed,

QOAS obtains better results in terms of perceived quality in comparison to both other solutions due to its slow steps-based policy of increasing the transmission rate to a level determined according to long-term behaviour-related information it maintains. Both LDA+ and TFRCP use a more aggressive manner of recovery after network problems and increase their transmission rate faster. This policy may achieve high throughput in some occasions, but when the background traffic varies sharply like in this situation at 360 s, it may lead to packet loss.

The effect of a potential high and steep increase in the background traffic when the system is already heavily loaded is tested at 250 s and 360 s. QOAS performs significantly better that both LDA+ and TFRCP-based adaptations reacting much faster to the sharp change in traffic, minimising the losses and therefore much reducing the period when the perceived quality is degraded. The TFRCP's average loss period is 20 s, the LDA+'s is 17 s, whereas the QOAS's is only 1.2 s.

At the end of the simulation, the effect on the tested streams of successive ends of individual streaming processes was also analysed. All the adaptive schemes have increased their rates to compensate for the decrease in background traffic, but LDA+ has done it faster than TFRCP and both much faster than QOAS. Nevertheless, the difference in the perceived quality between the results of these adaptive solutions was less than 2% during this period, which is not highly significant.

More detailed statistics related to the behaviour of these streaming schemes in the tested conditions are presented in Table 6-28. The statistical values from the table are computed for the duration of these tests (480 s) and do not include two 50 s transitory periods at the beginning and at the end. These performance-related values show how much improvement the QOAS brings in comparison to the other tested schemes.

Table 6-28 Statistical comparison between QOAS, TFRCP, LDA+ and NoAd when streaming *diehard1* in multimedia-like background traffic conditions

Streaming Scheme	Avg. Tx. Rate (Mb/s)	Avg. Loss Rate (%)	Avg. Perceived Quality (1-5)	Avg. Link Utilisation (%)
QOAS	3.12	0.015	4.384	99.93
TFRCP	3.16	1.057	3.789	99.88
LDA+	2.95	1.465	3.766	99.67
NoAd	4.00	13.667	1.490	100.00

6.2.3.5 Multiple QOAS-based Streaming in Highly Loaded Conditions

6.2.3.5.1 Overview

The set of tests reported in this section focuses on assessing the performance of the delivery of multiple multimedia streams using QOAS. The number of these streams is incrementally increased so that increased traffic delivery conditions are determined. The "Dumbbell" topology, already presented in section 6.2.1.2, is used for testing. Since the QOAS is designed for local broadband multi-service IP networks in which the majority of traffic is expected to be multimedia-based, this section analyses in detail the benefits brought by using QOAS for streaming a high number of concurrent multimedia clips of different types. These benefits are related at all times to other streaming solutions' such as LDA+, TFRCP and NoAd.

The tests presented in this section use the QOAS model that was presented in section 6.2.1.3 and LDA+, TFRCP and NoAd models, described in section 6.2.3.3.2. The multimedia clips used during testing are: diehard1 with high motion content, jurassic3 and dontsayaword with average motion content and familyman with low degree of action, as well as the roadtoeldorado, a cartoons movie. The clips were encoded at five different rates between 2 Mb/s and 4 Mb/s and traces were collected, associated to corresponding quality states and used during simulation. Statistics about these sequences and more information about the collected traces are presented in section 6.2.1.4. As mentioned in section 6.2.1.5 these streaming solutions' related performances are assessed in terms of resulted end-user perceived quality, loss rate and link utilisation. The estimated end-user perceived quality is computed using the no-reference moving pictures quality metric (Q) presented in section 2.4.3.2.10 and described in detail in section 4.4 and expressed using the ITU-T R P.910 five-point scale for grading subjective perceptual quality [63].

6.2.3.5.2 QOAS, TFRCP, LDA+ and NoAd Testing

The simulations involve a number of clients that randomly select both the movie clip and the starting sequence from within the chosen clip. They do not take into account other factors such as for example the popularity of the movies. The length of the simulations was 250 s, but when statistics were gathered the first and last transitory 50 s were not taken into account.

The QOAS, TFRCP, LDA+ and NoAd approaches were used in turn as the video streaming method, and the number of clients was gradually increased above a base line of 23 in each case. This number of clients was chosen because it allowed for lossless streaming and maximum end-user perceived quality in each of the four cases. Figure 6-76 shows the loss rate as a function of the

increase in the number of simultaneously served clients, Figure 6-77 presents the end-user quality as a function of the increase in the number of simultaneously served clients, and Figure 6-78 plots the bottleneck link utilisation values when the number of clients similarly increases.

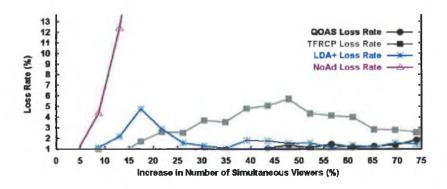


Figure 6-76 Loss rate vs. increase in the number of served clients above a base line of 23

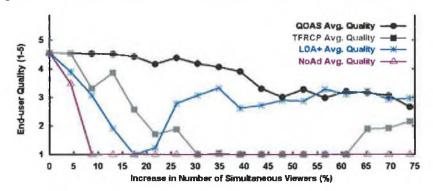


Figure 6-77 End-user average quality versus increase in the number of clients simultaneously served above a base line of 23

The results presented in Figure 6-76 show that in the NoAd case, an increase of only 4% in the number of clients caused a loss rate of just below 1%. When the number of clients was increased by more than 15%, the loss exceeded 10%, severely affecting the perceived quality, which drops quickly to the minimum level 1 ("bad") on the ITU-T R. P.910 five-point scale.

Under identical conditions, when QOAS was used, an increase of up to 40% in the number of clients (32 viewers) had very little effect on the loss rate, which remained below 0.5%. Figure 6-77 shows how for QOAS the resulting end-user quality remained above the "good" level of 4. Increases of up to 70% in the number of clients (39 viewers) resulted in loss rates of around 1%, which did not significantly affect the stream quality, which remained above the "fair" level of 3. Further increases in the number of clients caused both an increase in the loss rate and a fall in the

perceived quality below the "fair" level, which is considered here as the minimum acceptable quality level.

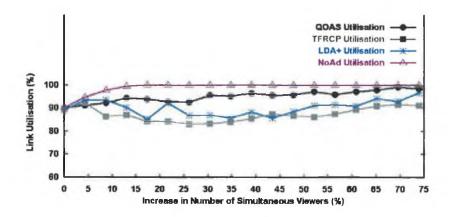


Figure 6-78 Bottleneck link utilization using different approaches, while increasing the number of simultaneous viewers

In comparison, tests using TFRCP streaming achieved only a 13% increase in the number of clients (26 viewers) when maintaining a loss rate below 1% and a corresponding perceived quality around the "good" level. For increases in the number of clients above 17%, the loss rate exceeded 1% and the end-user quality fell below the "fair" level. Given similar increases in the number of clients, LDA+ maintains an average loss rate below 1% and a perceived quality above the "good" level only for 24 clients (4% increase). However it maintained a "fair" end-user quality level for 30 simultaneous clients (30% increase) and loss rates around 1% for all tests performed in highly increased traffic conditions.

In terms of efficient usage of available bandwidth, QOAS was superior at all times to TFRCP and LDA+-based streaming. Using QOAS, the bottleneck link utilization exceeded 95% for 30 simultaneous clients and reached 99% for 40 clients. The values obtained for TFRCP and LDA+ are more modest: around 84% and respectively 87% for 30 simultaneous clients, and 92% and respectively 96% for 40 clients. Under the same conditions, the 100% figures obtained by NoAd came with severe costs in terms of loss and significantly reduced end-users quality.

6.2.3.6 Multiple QOAS - Comparison to Other Streaming Solutions

QOAS was used for streaming multiple multimedia clips to an increasing number of simultaneous clients so that the delivery network becomes increasingly high loaded. In identical conditions, TFRCP and LDA+-based adaptive solutions and a non-adaptive mechanism were used for similar multiple clips' streaming and the performance-related data was collected and compared to the QOAS's. The performance is assessed in terms of average end-user perceived quality, average loss rate and average link utilisation in all the cases presented in the previous section.

Table 6-29 shows comparative performance-related statistics for all the tested streaming approaches when choosing "fair" and "good" subjective quality levels as targets. In the table the increases in the number of clients are computed relative to the NoAd case. Since no post-processing techniques were applied, the "fair" level was considered here as the minimum quality level of interest. However further increases in the number of clients could be achieved by using for example different error concealment solutions, in order to mask the resulting losses that would otherwise severely affect the end-users' perceived quality.

Table 6-29 Statistical comparison between QOAS, TFRCP, LDA+ and NoAd when streaming multiple multimedia clips

Streaming Scheme	QOAS		TFRCP		LDA+		NoAd	
Quality	"fair"	"good"	"fair"	"good"	"fair"	"good"	"fair"	"good"
Loss rate (%)	1.39	0.47	1.73	0.53	1.31	0.50	0.81	0.01
Link utilisation (%)	95.7	96.4	84.1	87.1	86.9	93.4	94.7	90.0
Number of clients	34	32	27	26	30	24	24	23
Increase in no. of clients (%)	41.7	39.1	12.5	13.0	25.0	4.4	-	-

The results presented in Table 6-29 show that for the same average end-user quality, "fair" or "good" in these examples, QOAS can accommodate a significantly higher number of simultaneous clients while achieving higher bandwidth utilisation. For example, to maintain a "good" perceptual quality level, by using QOAS 23% more clients could be served than by using TFRCP, 33% more clients than by using LDA+, and 39% more users than by using the NoAd solution. If the goal is to maintain a "fair" average quality level for the clients, the benefit of using

QOAS is 26% greater than TFRCP, 13% greater than LDA+, and 42% greater than NoAd. The results are even more impressive if compared to the NoAd scheme as in the table. In terms of efficient usage of available bandwidth, QOAS was superior at all times to TFRCP and LDA+-based streaming, but inferior to NoAd, which pays for this with a significant decrease in its associated end-user perceived quality.

Comparing the schemes' performances for the same number of clients, the average end-user quality is always higher for QOAS than for the other solutions tested. Table 6-30 presents comparative performance results for these tested schemes obtained during some of the performed tests when streaming multimedia to certain numbers of simultaneous clients.

Table 6-30 Performance comparison between QOAS, TFRCP, LDA+ and NoAd when streaming multiple multimedia clips to the same number of clients

Streaming Scheme		QOAS			TFRCI			LDA+			NoAd	
No. of clients	Loss Rate (%)	Link Util. (%)	Perc Qual (1-5)									
23	0.00	90.04	4.56	0.00	89.54	4.56	0.00	89.12	4.56	0.00	90.04	4.56
26	0.00	94.34	4.51	0.53	87.06	3.86	2.19	90.28	1.91	12.34	99.43	1.00
27	0.05	93.68	4.42	1.73	84.13	2.58	4.77	85.18	1.00	23.57	100.0	1.00
32	0.47	96.38	4.01	4.82	85.42	2.62	1.82	88.28	1.00	>50.0	100.0	1.00
35	1.11	97.06	3.28	4.35	86.18	1.00	1.59	91.04	2.87	>50.0	100.0	1.00
39	1.38	99.07	3.06	2.83	91.59	1.93	1.57	92.88	2.93	>50.0	100.0	1.00

Both TFRCP and LDA+ seem to perform better for very high loads (when their loss situation behavior is applied) than for an average number of clients when loss and zero-loss periods alternate. In comparison, QOAS has a linear and more predictable response to an increase in the number of clients, which is a significant advantage of the QOAS scheme. In this way QOAS facilitates the choice of network load level according to economic, technical, and quality goals. However, QOAS was designed for local broadband multi-service IP-networks and therefore it seems likely that it will be used by service providers and network operators in order to maximise their revenues from offering VoD services to an increased number of clients while delivering a target quality level. For example, by scaling these simulation results with the "good" target quality

level to a one gigabit Ethernet connection, QOAS could service 320 simultaneous users compared to only 260 using TFRCP, 240 using LDA+, and 230 using NoAd streaming.

6.2.3.7 Effect of Feedback Frequency on End-user Perceived Quality

The goal of the set of tests whose results are presented in this section is to determine what is the effect of the variation in the frequency of feedback sent by QOAS on the multimedia stream quality, as it is perceived by the end-users.

The tests involve a five quality state QOAS server streaming *diehard1*, the multimedia sequence with high motion content (see Table 6-1), a QOAS client over the "Dumbbell" topology, which was described in detail in section 6.2.1.2. Background traffic that simulates real-life multimedia-like traffic with the variation presented in Figure 6-79 is generated on top of a 95.5 Mb/s CBR traffic outputted by the NS-2 CBR traffic model, which simulates a well-multiplexed natural traffic. This traffic determines loaded delivery conditions on which the QOAS with different feedback frequency are tested. For this the QOAS model, described in section 6.2.1.3, is used. As previously mentioned, the model consists of a QOAS server component, located at the sender and a QOAS client component, located at the receiver.

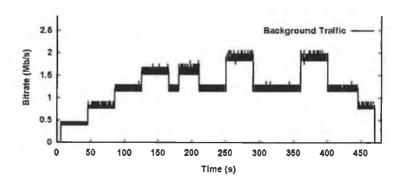


Figure 6-79 Multimedia-like background traffic variation on top of 95.5 Mb/s CBR traffic

The time between two consecutive feedback reports sent by the QOAS client to the QOAS server is varied from 0.01 s to 10 s and the QOAS's performance related results, expressed in terms of average transmission rate, average loss rate, average perceived quality and average link utilisation, are shown in Table 6-31.

Table 6-31 Effect of feedback frequency on the QOAS performance when streaming *diehard1* in multimedialike background traffic conditions

Feedback Interval (s)	Avg. Tx. Rate (Mb/s)	Avg. Loss Rate (%)	Avg. Perceived Quality (1-5)	Avg. Link Utilisation (%)
0.01	3.22	0.242	4.332	99.97
0.05	3.21	0.071	4.394	99.99
0.1	3.12	0.015	4.384	99.93
0.5	3.20	0.048	4.374	99.99
1.0	2.99	0.327	4.189	99.79
2.0	2.98	0.089	4.277	99.79
5.0	2.86	0.056	4.264	99.70
10.0	3.26	1.315	3.379	99.98

Analysing the results it is significant to mention that in general the end-user perceived quality decreases with the increase in the inter-feedback transmission time as expected since the control of the scheme becomes less tight. For very low feedback frequencies the QOAS's server component may not receive fast enough information about changes in the delivery conditions affecting the whole scheme's reaction to traffic variations and therefore not being able to avoid losses in loaded network situations. For example for an inter-feedback transmission time of 10 s the average end-user perceived quality has decreased to 3.38, around the "fair" subjective level from 4.39 much above the "good" perceptual level achieved when the feedback interval was set to 0.05 s.

In this context it seems that feedback has to be sent as often as possible. However sending feedback at high rates has at least two major disadvantages. First feedback takes bandwidth that is expensive and scarce in the environment the QOAS was designed for. Then processing feedback takes CPU computation time at both client machine and most important at the server. The latter can be easily overwhelmed by a very high number of feedback messages received from its clients. In consequence a compromise must be found for the inter-feedback transmission time, balancing the need for high quality with the low usage of shared resources, while taking into consideration the recommendations made in the RTCP standard [100] that specifies that feedback has to account for less than 5% of the bandwidth.

At the beginning the bandwidth used for the feedback transmission (BW $_{\text{feedback}}$) is computed for a single customer as in equation (6-15) and (6-16), where Time $_{\text{feedback}}$ is the inter-feedback transmission time.

$$BW_{feedback} = \frac{1}{Time_{feedback}} * Size_{feedback}$$
 (6-15)

$$Size_{feedback} = Size_{IPheader} + Size_{UDPheader} + Size_{RTCPheader} + Size_{Payload}$$
 (6-16)

For standard values for the headers' sizes (i.e. 20 bytes – IP header, 8 bytes – UDP header and 8 bytes – RTCP QOAS receiver report packet header) and for the size of the payload of 4 bytes, the feedback packet size becomes 40 bytes. At a very low average inter-feedback transmission time of 0.01 s the bandwidth used by feedback for a single client becomes BW_{feedback} = 4,000 bytes/s. Taking into consideration that QOAS solution was designed for delivering multimedia in increased traffic over local broadband multi-service IP-networks, for a one gigabit Ethernet on which 320 customers are being served with "good" perceived quality as shown in section 6.2.3.6, the total bandwidth used by feedback sent with this frequency is: 320*4,000=1,280,000 bytes/s. This figure, in fact 9.77 Mb/s, is less than 1 % of the total available bandwidth and does not add too much to the existing traffic. Unfortunately the number of feedback messages (No_{feedback}) that the QOAS server application must deal with in the presence of an increased number of customers (No_{customers}) becomes very high as computed with the formula from equation (6-17) and reaches (1/0.01)*320=32,000 every second for 0.001 s feedback interval.

No feedback =
$$\frac{1}{Time\ feedback} * No\ customers$$
 (6-17)

In order to lower the load from the server, decreasing at least ten times the feedback frequency is recommended. The expected benefit in reducing ten times the number of feedback messages that have to be processed by the QOAS server is followed by reducing ten times the used bandwidth that decreases to 0.1 % of the existing capacity. However, the end-user perceived quality is also reduced, but with not a significant value. Further decreases in the feedback transmission frequency may lead to more significant effects on the end-user perceived quality, which is decreasing to pay for further lowering the pressure on the QOAS server and the used bandwidth for feedback.

Therefore we recommend using a 0.1 s interval between consecutive feedback messages since this value balances the QOAS server component's need for fast and accurate information regarding the delivery conditions as reported by the client component with the use of shared resources.

6.2.3.8 Effect of Delivery Latency on End-user Perceived Quality

The set of tests whose results are presented in this section aims at determining whether QOAS is affected by the variation in the latency of the link delivery in the corresponding end-user perceived quality of the QOAS-streamed multimedia clips.

The tests involve the "Dumbbell" topology, which was described in detail in section 6.2.1.2. Its bottleneck link delivery latency is varied and the effects on the QOAS-related performance results when streaming *diehard1*, a multimedia sequence with high motion content presented in Table 6-1, are analysed. Background traffic that simulates real-life multimedia-like traffic with the variation presented in Figure 6-80 is generated on top of a 95.5 Mb/s CBR traffic outputted by the NS-2 CBR traffic model, which simulates a well multiplexed natural traffic. This traffic determines loaded delivery conditions on which the QOAS is tested. The QOAS model, described in section 6.2.1.3, is used and involves a five quality state QOAS server component streaming to a QOAS client component.

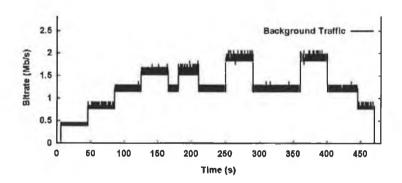


Figure 6-80 Multimedia-like background traffic variation on top of 95.5 Mb/s CBR traffic

The bottleneck link delay is varied from 0.01 s to 0.5 s while maintaining constant the interfeedback transmission interval of 0.1 s and the QOAS's performance related results, expressed in terms of average transmission rate, average loss rate, average perceived quality and average link utilisation, are shown in Table 6-32.

Table 6-32 Effect of delivery latency on the QOAS performance when streaming diehard1 in multimedia-like background traffic conditions

Delivery Latency (s)	Avg. Tx. Rate (Mb/s)	Avg. Loss Rate	Avg. Perceived Quality (1-5)	Avg. Link Utilisation (%)
0.01	3.17	0.031	4.391	99.94
0.05	3.07	0.026	4.353	99.84
0.1	3.12	0.015	4.384	99.93
0.2	3.11	0.280	4.279	99.89
0.5	3.16	0.777	4.086	99.90

It is significant to mention after analyzing these results that in general the QOAS's related end-user perceived quality decreases with the increase in the delivery link latency. This conclusion may seem natural since the longer the time the client has to wait for its reports about the quality of the delivery to be received and processed by the server and for the consequent adjustments to be felt back at the receiver, the greater the chance these adjustments not to match the new existing delivery conditions. For very long delays the QOAS's server component may not receive fast enough information about changes in the delivery conditions affecting the whole scheme's reaction to traffic variations and therefore not being able to avoid losses in loaded network situations. For example for a link delay of 0.5 s, the average end-user perceived quality has decreased to 4.09, at the "good" subjective level from 4.39, much above the "good" perceptual level achieved when the delivery latency was 0.01 s.

In this context it seems that feedback has to arrive at the server as fast as possible. However the link latencies depend very much of the architecture of the local broadband IP networks and in general the shortest the link delay, the more expensive the solution is. In consequence a compromise must be found for the link delay, balancing the need for high quality with the infrastructure-related costs of the solution.

Therefore it is recommended using 0.1 s as target for the maximum delivery latency since this value balances the QOAS server component's need for fast feedback with the resource-related costs.

6.2.4 Comments

Extensive objective simulation tests, based on NS-2 and both on its built-in and our specially built models, were performed in order to both tune the QOAS and test it. Tuning aimed at determining the design parameters for QOAS that lead to obtaining the best results in terms of estimated average end-user perceived quality of the QOAS-streamed multimedia clips over local broadband multi-service IP-networks. Once the values for these parameters were set, the goal of testing was to determine the QOAS's performances in terms of end-user perceived quality, loss rate, link utilisation and the number of simultaneous viewers served from a finite infrastructure. In consequence these tests have involved single QOAS-based multimedia streaming in loaded delivery conditions and subject to different background traffic. This traffic has included traffic of different types, shapes and variation patterns commonly encountered in IP-networks as well as multimedialike background traffic. The results were both analysed as they are and compared to those obtained by an ideal adaptive scheme and by existing other streaming solutions such as TFRCP, LDA+ and non-adaptive. Multiple QOAS-based multimedia streaming processes were simulated next and the results were compared to those obtained when using other streaming solutions. The effects of feedback transmission interval and of delivery link delay were also analysed.

In all tested situations QOAS has achieved very good results related to performance, even compared to the ideal adaptive scheme from whose performances QOAS's were very close. The results were significantly better than the ones obtained when using other streaming schemes in all tested conditions. Therefore these objective tests have shown that QOAS achieves link utilisations very close to 100 %, very low loss rates, a significant increase in the number of customers served from the same infrastructure and high estimated end-user perceived quality. However since there is not a generally accepted metric for measuring the latter, subjective tests are necessary to verify these objective results obtained using the no-reference moving picture quality metric (Q). The results of the subjective testing are presented in the following section.

6.3 Subjective Testing

6.3.1 Motivations

The objective testing results related to the performances of multimedia streaming when using QOAS were very significant, showing important benefits brought in terms of high end-user perceived quality, low loss rates, increased link utilisation and high number of simultaneous

viewers served from a given infrastructure. Since simulations were used to perform these tests, the end-user perceived quality could only be estimated using the no-reference moving pictures quality metric – Q, presented in section 4.4. Since there is not a standardised metric for measuring the end-user perceived quality when streaming multimedia clips and neither a general accepted metric that would very accurately estimate the end-user's subjective assessment of the quality of the remotely played stream, it was decided to use perceptual tests that involve real subjects in conjunction with the simulation tests in order to verify the results obtained by the latter.

6.3.2 Setup Conditions

6.3.2.1 Test Setup

In order to perform the real streaming tests, the test bed presented in Figure 6-81 was assembled. It consists of a local **Server** machine and a local **Client** computer, each part of a different network interconnected by a **Router**. An emulator installed on the Router captures all the packets and forwards them to the other network after introducing bandwidth and delay constraints.

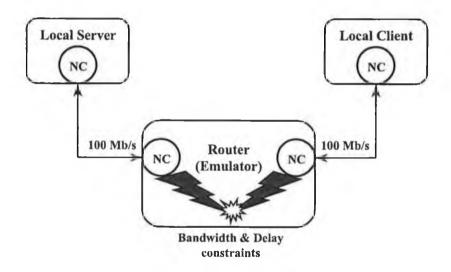


Figure 6-81 Test bed setup consisting of a local Server and a local Client part of different networks interconnected by a Router on which an emulator allows for bandwidth and delay variation

The Server is an IBM Netfinity 6000 R computer with two 700 MHz processors and 1 GB RAM on which Microsoft Server 2000 Advanced Edition is installed. The Client is a Fujitsu-Siemens Scenic machine with one 800 MHz processor and 512 MB RAM, with Microsoft Windows 2000 Professional as the operating system, whereas the Router is another Fujitsu-Siemens Scenic computer with one 800 MHz processor and 512 MB RAM, on which Linux was installed in order to facilitate the deployment of the NistNet Emulator [254]. The network cards (NC) are 3Com EtherLink XL PCI Combo NIC 3C900B at 100 Mb/s and UTP connections are used. The client has a Desktop PC with a 19 inch monitor.

6.3.2.2 Applications' Setup

QOAS server application, part of the QOAS prototype system whose implementation details were given in the fifth chapter, was installed on the server computer. In order to allow for the MPEG encoding of multimedia clips, a Canopus Amber Encoder/Decoder card was also installed on the server machine. Then the multimedia database that includes the multiple versions of pre-recorded multimedia clips was registered with the ODBC Data Source Administrator allowing it to be accessed and communicated with. The QOAS client application was deployed on the client machine. It makes use of a Canopus Amber MPEG Decoder card which was installed on the same machine.

On the server computer the latest Microsoft version of Media Producer⁶⁹ (series 9) [253] was installed and Windows Media Services was enabled. In this way the server side streaming application was deployed. This allows for Windows Media (WM) file streaming, including the Multiple bit-rate (MBR) ones that QOAS will be tested against. On the client machine, Windows Media Player series 9 was installed constituting the client application for WM streaming.

6.3.2.3 Tested Approaches

During subjective testing three different approaches are assessed by the test participants with different motion content clips and background traffic variation: the streaming of multimedia clips based on QOAS-the adaptive scheme proposed in this thesis, the commercially available adaptive Windows Media (WM) Multiple bit-rate (MBR) solution, launched as part of the WM series 9 products and a non-adaptive streaming solution.

⁶⁹ Windows Media, Web Site, Microsoft, http://www.microsoft.com/windows/windowsmedia/default.asp

6.3.2.4 Test Environment

The test environment was determined according to the ITU-T R P.910 recommendations [63]. The streamed multimedia clips are displayed on a Desktop PC 19 inches monitor situated in room with no natural light. The only source of light barely allows for the answer sheets to be filled in, it is localised and it does not reflect in the monitors nor disturb the subjects. The parameters for the monitor (brightness, luminance, hue etc.) have been set at average values. The viewing distance was set at 5 times the height of the picture, within the limits suggested by ITU-T R P.910 and should remain fixed for the duration of the testing. The audio component of the multimedia is played out by two 10 W Creative Cambridge SoundWorks SBS52 speakers which are the only source of sound in the testing room.

6.3.2.5 Multimedia Clips

In order to perform the tests, multimedia clips were encoded from high quality DVD sources. The WM Producer was used to encode a multiple bitrate (MBR) stream that can adapt to five audiences at 2.0 Mb/s, 2.5 Mb/s, 3.0 Mb/s, 3.5 Mb/s and 4.0 Mb/s. For the QOAS-based system, five streams were MPEG-2 encoded at 2.0 Mb/s, 2.5 Mb/s, 3.0 Mb/s, 3.5 Mb/s and 4.0 Mb/s. This process was repeated for different movies with various motion content or types, maintaining constant the IBBP-pattern, the number of frames per GOP at 9 and the resolution at 320 x 240. 15 minutes long multimedia sequences were encoded from the following movies: *Die Hard 1* — with very high motion content, *Jurassic Park 3* with an average — high motion content, *Don't Say A Word*, with average — low motion content and *Family Man* with very little action in it. A cartoons movie was also encoded — *Road To El Dorado*. From these clips shorter sequences were used as source files for multimedia streaming.

Due to the fact the testing time for each subject has not to exceed 30 minutes according to the suggestion made by ITU-T R. P.910, the multimedia clips use for testing was limited to four, leaving aside the sequence from *Jurassic Park 3*. Since the subjects' attention has a time limit, only 1 minute-long clips were used from each movie, for each test. This is unlike the ITU-T R P.910 recommendations that suggest using 10 s long sequences, but we support the opinion presented in [132] which states that such a short sequence is not enough to allow the subject both to accommodate with the particular movie content and to notice quality differences, especially if the quality varies in time for each sequence, not only between sequences.

6.3.2.6 Test Method

The chosen test method is a combination between the Absolute Category Rating (ACR) or Single Stimulus (SS) and the Degradation Category Rating (DCR) or Double Stimulus (DS), presented in detail in ITU-T R P.910 [63]. ACR involves the presentation of the multimedia sequences one at a time and the subject is asked to grade each of them separately on a given category scale. In this case the sequences with identical content would be streamed using different solutions one after the other, with short breaks for grading after each of them. DCR implies the fact that the test sequences are presented in pairs. The first stimulus presented is always the reference while the second sequence is the tested one. In our case this second sequence has the same content, but it is delivered using a different approach. The subject is asked to grade only the quality of the second multimedia clip.

Unfortunately in order to apply only the ACR test method, the implicit reference must be well known by all the assessors, and this cannot be expected in this case. If only the DCR method had been applied, the reference clips would have to be displayed too many times, the test would take too long, the subjects would become bored and the accuracy of the results would suffer. Therefore the ACR and DCR methods were combined and therefore a reference clip is shown first and then the multimedia sequences that have to be assessed. After each of them the subject is asked to grade its subjective quality on the given quality scale. The grading process should be very short in order to minimise the time passed since the viewer has seen the reference clip and also to minimise the total duration of the testing procedure.

6.3.2.7 Grading Scale

The multimedia quality could be graded on different scales such as, for example, a binary one (e.g. good/bad), a continuous graphical scale with no explicit labels (e.g. ranging between bad and excellent), or the quality scales for subjective testing suggested by ITU-T R P.910 with 5, 9 or 11 points. Since lately many systems, including commercial ones, have been rated on the 1-5 scale (see Table 6-33), which offers enough information in order to significantly assess the results being also simple to work with, this grading scale was selected for the perceptual tests, too. Apart from this, the quality metric Q used to objectively assess the quality of the streamed multimedia during the simulations uses the same 1-5 scale, allowing for a simple comparison between the simulation objective test results and these subjective test results.

Rating **Impairment** Quality 5 Imperceptible **Excellent** 4 Perceptible, not annoying Good 3 Slightly annoying Fair 2 Annoying Poor 1 Very annoying Bad

Table 6-33 Quality scale for subjective testing

6.3.3 Tests Description and Goals

6.3.3.1 Test Goals

The subjective tests performed have two main goals:

- Quantification of the perceived quality of the multimedia clips streamed using QOAS adaptive approach in highly loaded delivery conditions and subject to multimedia-like background traffic which should account for the majority of the traffic in the local broadband multi-service IP-networks QOAS was designed for. These conditions force the QOAS multimedia system to adjust the transmitted quantity of data by modifying the clips' quality. The intention is to test whether these adaptive variations are noticed by the viewers, are acceptable or disturbing for them and with what degree. It is important also to study if the movies' motion content affects differently the perceived quality result as subjectively graded by viewers by using clips with high, average and low motion content, as well as a cartoons clip.
- Comparison of the QOAS-based adaptive streaming with non-adaptive and multiple bitrate (MBR) Windows Media remote multimedia delivery. We would also like to
 determine what were the most appreciated features and the least liked characteristics of
 the streaming performed with each of these schemes.

These tests aim at complementing the simulation test results, verifying their findings: confirming or contradicting them.

6.3.3.2 Tests' Description

The tests performed using the simulation model and presented in the sections 6.2.3.3 and 6.2.3.4 of this chapter consist of delivering the multimedia clip with the highest motion content *diehard1* using QOAS and other streaming approaches over the "Dumbbell" topology (presented in section 6.1.2.5) that raises similar problems as those in a local broadband IP-network. The deliveries were subject to loaded delivery conditions and multimedia-like background traffic that simulated viewers' VCR interactivity such as play, pause and stop. Among the analysed results (the estimated end-user perceived quality, the loss rate and the link utilisation), the viewers' subjective quality assessment will be verified by the perceptual tests that involve real subjects presented next.

Since the duration of each test has to be minimised and since different aspects of the compared streaming schemes have to be tested, two separate tests were devised that differ in terms of the background traffic and consequently of the degree of expected reaction from the streaming schemes. In order to test the schemes in most difficult conditions, from within the background traffic variation presented in Figure 6-59 the sequences when it varies in a staircase-up manner with step size of 0.4 Mb/s and when the variation is periodic with step size of 0.7 Mb/s, above the adaptation step of 0.5 Mb/s were selected. In consequence the first test aims at determining the schemes' reactions and their effect on the end-users' perceived performance in loaded and variable delivery conditions, which have not caused loss during QOAS-based simulations. The second test, which involves a sequence of background traffic that has caused short periods of loss during simulations when streaming using QOAS, intends to determine how these expected lossy periods affect the end-users' grading of the clips' overall quality.

Figure 6-82 and Figure 6-83 show the multimedia-like background traffic variation during the first and the second test. This traffic variation is on top of a CBR traffic that generates loaded delivery conditions and represents well-multiplexed different types, shapes and sizes individual traffic flows.

6.3.3.2.1 Test 1 - Staircase-up Multimedia-like Background Traffic

Background traffic is increasing in a staircase-like manner every 20 s, with three steps of 0.4 Mb/s and starting from the level set by a first step of 0.4 Mb/s, not part of these tests. This traffic is on top of a 95.5 Mb/s CBR traffic that represents various individual traffic flows that are well-multiplexed determining loaded delivery conditions. Figure 6-82 shows the consequent background traffic variation, which is replicated using the NistNet emulator [254] and aims to cause

a reaction from the QOAS prototype system, from the Windows Media system and to affect the non-adaptive streaming. Streaming each clip takes 1 minute and it is performed for each of the four selected multimedia clips with different motion contents and types, involving each of the three tested streaming approaches. Before using these approaches, the testing methodology suggests reference streaming of the same clip at maximum achievable quality in the testing conditions (4 Mb/s in this case).

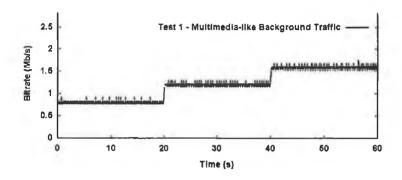


Figure 6-82 Staircase-up background traffic on top of 95.5 Mb/s CBR traffic during Test 1

6.3.3.2.2 Test 2 - Periodic Multimedia-like Background Traffic

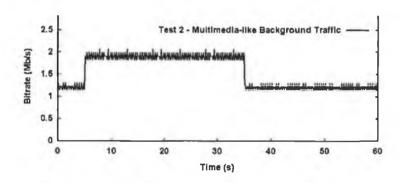


Figure 6-83 Periodic background traffic on top of 95.5 Mb/s CBR traffic during Test 2

During the second test the background traffic is part of a multimedia-like traffic that periodic varies with a pattern having an on period of 30 s and an off period of 60 s and an amplitude of 0.7 Mb/s, much higher than the QOAS scheme's adaptation step of 0.5 Mb/s. This traffic is on top of a 95.5 Mb/s CBR traffic that determines loaded network delivery conditions. Since the duration of the test is 60 s, it includes the on period and half of the off period as shown in Figure 6-83. This background traffic variation is replicated using the NistNet emulator affecting the streaming processes and consequently the end-user perceived quality in a higher or a lower degree,

depending on the scheme's capability to adjust to these variations. All three tested approaches (QOAS, WM MBR and non-adaptive) are used for streaming and all four selected movie clips are streamed. As in the first test type, reference streaming is required for each clip at maximum achievable quality (4 Mb/s in this case) in existing setup conditions.

6.3.3.2.3 Test Phases

Each of the two test types consists of four phases that each involves a different clip from the four movies with different motion content taken into consideration and presented in section 6.3.2.4. In order not to bias the viewers' decision regarding a movie type or another, the order in which these clips were shown was randomised. For each multimedia clip, first the reference streaming is performed at the highest quality taken into consideration (4.0 Mb/s) and the delivery is not subject to any background traffic that might interfere with its quality as seen by the remote viewer. Then each of the three streaming approaches is employed for delivering the clip to the remote viewer and after each of them, the subject is asked to grade its quality and to highlight the quality-related feature he/she liked the most and the one that he/she disliked the most. The order in which these approaches were used with the same clip was also randomised not to affect the obtained results.

6.3.3.2.4 Test Considerations

In order to ensure good testing results, it is very significant to include a **training phase** prior to starting the test sequence. In this training phase the test operators have to explain what is the goal of these tests and what it is required from each participant. More detailed information about this phase is given in Appendix C.

Since the **subjects' visual acuity** greatly differ and some of the participants may suffer from visual impairments that may affect their assessment of the streamed multimedia clips' quality, prior to testing the viewers should be screened for normal visual acuity or corrected-to-normal acuity and for normal colour vision. The results of these findings can be used during the results' analysis.

Participants' boredom or fatigue have an important impact on the multimedia clips' quality assessment as well as on the accuracy of the answers. Therefore the participation to testing should be voluntary, so the subjects could leave at any time.

Although the QOAS adaptation does not interfere with the audio component (we have considered that it takes only a small fraction of the bandwidth in comparison to the video

component), since the quality of any multimedia sequence is influenced also by the associated sound (e.g. audio and video must be appropriately synchronized), during testing the clips are streamed with their soundtracks.

During testing the participants are asked to indicate whether they have **liked** some **characteristics** related to the quality of the multimedia streaming such as continuity, audio/video synchronisation, clarity etc. They are also asked to mark any defects they noticed and they have **disliked** during streaming such as tiling, jerkiness, de-synchronisation etc.

A sample of a **test questionnaire** is presented in Appendix C.

6.3.4 Tests Results

6.3.4.1 Test 1 - Staircase-up Multimedia-like Background Traffic

The subjective *Test 1* has involved 42 subjects with ages between 18 and 48, with various experience related to multimedia streaming (i.e. 22 - familiar, 19 - not familiar and 1 - expert), 19 of which wearing glasses or contact lenses and none with other visual impairments that may affect their perception of the multimedia quality. The *Test 1* results are presented in the next tables as follows. Table 6-34 presents statistics related to the average subjects' perceived quality when the four multimedia sequences named *Die Hard 1*, *Don't Say a Word, Family Man* and *Road to El Dorado* were streamed in the background traffic conditions mentioned in section 6.3.3.2.1. Table 6-35 presents the test results related to the participants' most appreciated of the clips' quality characteristics such as continuity, quality stability, image clarity and media synchronisation in all the tested situations during *Test 1*. The figures in the table represent the percentage of the subjects that have appreciated the most the associated multimedia stream characteristic. Table 6-36 presents the percentage of the participants that have mostly disliked certain multimedia streaming related features indicated in the table such as jerkiness, quality variation, blurring, tiling and media desynchronisation. The results are presented for all the tests performed, including all the streaming schemes and involving all the multimedia clips taken into account.

Table 6-34 Statistical results related to subjective quality assessment on the 1-5 grading scale obtained for Test 1 for all the *Die Hard 1*, *Don't Say a Word*, *Family Man* and *Road to El Dorado* multimedia clips

Movie Clip	Die I	Hard 1	Don't Say A Word		Famil	y Man	Road To El Dorado		
Streaming Scheme	Avg.	Std.Dev.	Avg.	Std.Dev.	Avg.	Std.Dev.	Avg.	Std.Dev.	
QOAS	4.00	0.71	4.18	0.75	4.21	0.83	3.74	0.71	
WM MBR	2.02	0.78	2.12	0.71	2.00	0.74	2.38	0.78	
Non-Adaptive	2.02	0.72	2.44	1.00	1.85	0.79	1.93	0.72	

The results from Table 6-34 show how the QOAS streaming was very appreciated by the test subjects, scoring above 4, the "good" quality level on the 1-5 ITU-T grading scale, for all the movies and close to 4 for the cartoons sequence. The low standard deviation values that are also presented in the table show that the results obtained are consistent, although the granularity of the grading process was quite coarse, since the difference between the acceptable grades was 1. These positive results become more significant if compared with WM MBR commercial solution that achieves grades above "poor" quality level or with non-adaptive streaming solution whose subjective quality scores are below the "poor" level.

The results obtained for QOAS seem to suggest that the higher the motion complexity of a sequence the lower the subjective appreciation in loaded delivery conditions. However, more tests are needed to verify such an assumption. Nevertheless there is significant difference between the subjective scores obtained for the clips that contain movie scenes and the cartoons clip. A potential problem might be the different MPEG-2 encoding output for the cartoons sequences as shown in Table 6-1. Unlike for the movie content, for cartoons content the peak/mean ratio computed in relation to the size of the encoded frames does not significantly increase with the decrease in the average encoding bit-rate. Also the content with many colors and edges might be more affected in terms of the end-user subjective quality corrupted during streaming.

Figure 6-84 presents the QOAS bit-rate adaptation with the variation of the background traffic during streaming of the *Die Hard 1* clip, multimedia sequence with the highest motion content. This adaptation is very similar to the QOAS bit-rate variation while streaming the *diehard1* sequence during the simulation tests whose results are presented in Figure 6-60.

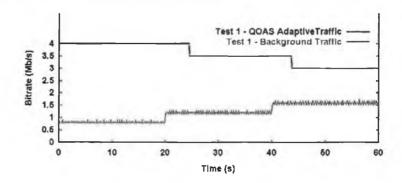


Figure 6-84 QOAS bit-rate adaptation with background traffic variation when streaming *Die Hard 1* clip during Test 1

Table 6-35 Statistical results related to what the subjects have appreciated the most when streaming *Die Hard*1 (A), *Don't Say a Word* (B), *Family Man* (C) and *Road to El Dorado* (D) multimedia clips during Test 1

(%)	QOAS				WM MBR			Non-Adaptive				
Clip	A	В	С	D	A	В	С	D	A	В	С	D
Continuity	52.4	66.7	69.0	47.6	2.4	4.8	7.1	11.9	4.8	16.7	9.5	4.8
Q. Stability	45.2	64.3	52.4	35.7	19.0	23.8	14.3	33.3	4.8	11.9	4.8	4.8
Clarity	78.6	78.6	71.4	69.0	42.9	61.9	50.0	52.4	35.7	38.1	31.0	28.6
Media Synch.	54.8	52.4	59.5	28.6	14.3	14.3	7.1	9.5	14.3	31.0	9.5	9.5

Analysing the results from Table 6-35, one could conclude that, regardless of the streamed content which influences only the degree of the opinion, the subjects have appreciated the same characteristics of the multimedia streamed clips related to their perceived quality. For example during streaming using QOAS the most appreciated was the clarity of the video content, followed by media synchronisation and continuity with results in generally much above 50 %. Although quality stability has scored less that the other features, it achieved on average high values (around 50 %), in spite of deliberately introduced variations in quality by the QOAS adaptation process. However streaming of the cartoons sequence was lower rated than the remote delivery and playing of the other clips.

WM MBR solution has scored high (around 50 %) only at image clarity and failed to impress the viewers in relation to its media synchronisation and continuity, which have got the least

number votes (around 10 % and respectively less than 10 %). Quality stability was appreciated by around 20 % of subjects.

The non-adaptive solution has obtained more than 30 % of votes only related to the clarity of the image, but did not attract appreciation related to any of the other features.

Table 6-36 Statistical results related to what the subjects have disliked the most when streaming *Die Hard 1*(A), *Don't Say a Word* (B), *Family Man* (C) and *Road to El Dorado* (D) multimedia clips during Test 1

(%)		QO	AS			WM	MBR			Non-A	daptive	e
Clip	A	В	С	D	A	В	С	D	A	В	С	D
Jerkiness	28.6	33.3	31.0	45.2	90.5	90.5	88.1	69.0	69.0	59.5	88.1	78.6
Q. Variation	26.2	14.3	14.3	28.6	19.0	9.5	23.8	16.7	71.4	64.3	73.8	71.4
Blurring	2.4	0.00	2.4	9.5	28.6	16.7	14.3	26.2	16.7	14.3	16.7	21.4
Tiling	9.5	11.9	23.8	21.4	4.8	4.8	9.5	2.4	69.0	61.9	78.6	69.0
Media Desyn.	14.3	19.0	9.5	45.2	64.3	47.6	64.3	64.3	59.5	47.6	66.7	59.5

Looking at the results contained in the Table 6-36 that relates to the QOAS performance, a surprisingly high percentage of subjects (30 %) have found jerkiness the most annoying aspect of the QOAS and only around 20 % the quality variations. Maybe some problems related to the implementation for the prototype system may have triggered such a result and less the QOAS-related aspects. A very low number of participants have indicated tiling (around 15 %), media desynchronisation (around 14 %) and blurring (almost none) as causes for dissatisfaction.

WM MBR solution has been mostly blamed for the jerkiness (almost 90 % of participants) and media de-synchronisation (around 60 %), although blurring and quality variation also have disliked to around 20 % of viewers.

The non-adaptive approach has negatively impressed the subjects from many points of view. More than 70 % of them have indicated that they have mostly disliked jerkiness, tiling and quality variation, whereas around 60 % were annoyed by media de-synchronisation and only 15 % were disturbed by blurring.

On aggregate, the results of Test 1 indicate that the QOAS solution achieves good subjective quality performance being appreciated by subjects both as stand-alone and in comparison to WM MBR and non-adaptive streaming solutions, confirming the objective testing results.

6.3.4.2 Test 2 - Periodic Multimedia-like Background Traffic

The subjective *Test 2* involved 42 subjects with ages between 21 and 45, with various experience related to multimedia streaming (i.e. 19 - familiar, 21 - not familiar and 2 - experts), 16 of which wearing glasses or contact lenses and none with other visual impairments that may affect their perception of the multimedia quality. The *Test 2* results are shown in the next tables in a similar manner with the *Test 1* results. Table 6-37 presents statistics related to the average subjects' perceived quality when the four multimedia sequences named *Die Hard 1*, *Don't Say a Word, Family Man* and *Road to El Dorado* were streamed in the background traffic conditions mentioned in section 6.3.3.2.2. Table 6-38 presents the test results related to the participants' most appreciated of the clips' quality characteristics such as continuity, quality stability, image clarity and media synchronisation, expressed as percentage of the total number of subjects, in all the tested situations during *Test 2*. Table 6-39 indicates the percentage of the participants that have disliked some multimedia streaming related features that are indicated in the table. The results are presented for all the tests performed, including all the streaming schemes and involving all the multimedia clips taken into account.

Table 6-37 Statistical results related to subjective quality assessment on the 1-5 grading scale obtained for Test 2 for all the Die Hard 1, Don't Say a Word, Family Man and Road to El Dorado multimedia clips

Movie Clip	Die l	Hard 1	Don't Say A Word		Famil	ly Man	Road To El Dorado	
Streaming Scheme	Avg.	Std.Dev.	Avg.	Std.Dev.	Avg.	Std.Dev.	Avg.	Std.Dev.
QOAS	4.22	0.69	3.98	0.64	4.24	0.66	3.85	0.69
WM MBR	2.32	0.69	2.62	0.70	2.36	0.73	2.33	0.66
Non-Adaptive	1.33	0.67	1.45	0.67	1.31	0.56	1.37	0.62

As with the results for *Test 1*, the results from Table 6-37 show the test subjects' appreciation of the QOAS-based streaming. QOAS has scored around and above 4, the "good" quality level on the 1-5 ITU-T grading scale, for all the movies and below 4, but close to it, for the cartoons sequence. The low standard deviation values presented in the table show that the results

obtained are consistent, and such low values were obtained in spite of the fact that the grading process has not accepted fractional quality grades. These positive results become more significant if compared with WM MBR commercial solution that achieves grades above "poor" quality level or with non-adaptive streaming solution whose subjective quality scores are close to the "bad" level. It is very important to notice that the short lossy periods that have occurred during *Test 2* have not significantly influenced the perceived quality of the results which are comparable with the ones obtained for *Test 1*.

Analysing these results it seems that there is not a quantifiable relationship between the motion complexity of a sequence and the subjects' quality appreciation in loaded delivery conditions. Yet, the significant difference between the subjective scores obtained for the clips that contain movie scenes and the cartoons clip has been maintained in highly increased delivery conditions that has also triggered loss. These delivery conditions made the difference between the WM MBR approach and the non-adaptive to become more significant in the favour of the former, which succeeds to adapt to the traffic conditions.

Figure 6-85 presents the QOAS bit-rate adaptation triggered by the background traffic variation when streaming *Die Hard 1* multimedia clip during Test 2. These results are similar to those obtained during simulations with the *diehard1* sequence and were presented in Figure 6-60. However, since the duration of the test was only 60 s, QOAS did not complete its adaptation period after the drop in background traffic and consequently the transmission rate does not reach the starting value.

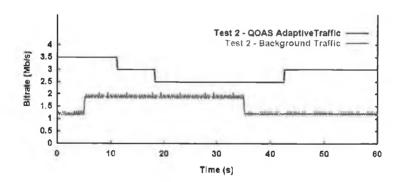


Figure 6-85 QOAS bit-rate adaptation with background traffic variation when streaming Die Hard 1 clip during Test 2

Table 6-38 Statistical results related to what the subjects have appreciated the most when streaming *Die Hard* 1 (A), *Don't Say a Word* (B), *Family Man* (C) and *Road to El Dorado* (D) multimedia clips during Test 2

(%)	QOAS				WM MBR			Non-Adaptive				
Clip	A	В	С	D	Α	В	С	D	A	В	С	D
Continuity	76.2	61.9	71.4	45.2	9.5	14.3	9.5	4.8	7.1	0.0	0.0	0.0
Q. Stability	57.1	47.6	54.8	45.2	31.0	35.7	35.7	35.7	7.1	0.0	0.0	2.4
Clarity	66.7	66.7	71.4	66.7	45.2	50.0	42.9	45.2	9.5	11.9	7.1	9.5
Media Synch.	66.7	52.4	64.3	38.1	16.7	14.3	14.3	14.3	9.5	0.0	0.0	0.0

Analysing the results from Table 6-38 obtained during streaming using QOAS in conditions imposed by *Test 2*, like in the case of *Test 1*, the most appreciated was the clarity of the video content and continuity with results close to 70 %, as well as media synchronisation that scored roughly 60 %. Although quality stability scores again less that the other features, it also reaches a high value around 55 %, in spite of the quality variations introduced by the QOAS and losses that have occurred and may have slightly decreased the subjective clips' quality.

WM MBR solution has scored high not only at image clarity (again around 50 %), but also at quality stability (around 35 %) unlike in the first test. It failed again to impress the viewers in relation to media synchronisation and streaming continuity.

The non-adaptive solution has obtained almost no appreciation, the solution collapsing from the quality point of view in these highly increased traffic conditions.

Table 6-39 Statistical results related to what the subjects have disliked the most when streaming *Die Hard 1* (A), *Don't Say a Word* (B), *Family Man* (C) and *Road to El Dorado* (D) multimedia clips during Test 2

(%)	QOAS				WM MBR			Non-Adaptive				
Clip	A	В	С	D	Α	В	С	D	A	В	С	D
Jerkiness	14.3	31.0	9.5	28.6	92.9	85.7	81.0	81.0	83.3	90.5	78.6	83.3
Q. Variation	16.7	16.7	23.8	21.4	19.0	16.7	11.9	16.7	66.7	64.3	57.1	64.3
Blurring	0.0	7.1	2.4	2.4	19.0	4.8	11.9	7.1	21.4	19.0	16.7	19.0
Tiling	26.2	28.6	23.8	38.1	2.4	4.8	0.0	2.4	88.1	85.7	90.5	78.6
Media Desyn.	9.5	16.7	4.8	31.0	66.7	52.4	66.7	66.7	76.2	81.0	83.3	76.2

Table 6-39 indicates in relation to the QOAS performance that around 25 % of subjects have tiling the most disturbing, followed by quality variation (roughly 20 % of them) and jerkiness (~17 %). However in such difficult delivery condition the number of dissatisfied viewers is very low and this is in favour of the QOAS-based solution.

WM MBR solution has been again mostly blamed for the jerkiness (almost 85 % of participants) and media de-synchronisation (around 60 %), and less for blurring, quality variation and tiling, in this order.

The non-adaptive approach determined more than 80 % subjects to indicate that they have mostly disliked jerkiness, tiling and media de-synchronisation, although quality variation also scored more than 60 %.

These results show that QOAS solution has successfully adapted even to very difficult delivery conditions achieving good subjective quality performance appreciation. *Test 2* confirms the results obtained by the first test relative to QOAS and in comparison to both WM MBR and non-adaptive streaming solutions, verifying also the objective testing results.

6.3.5 Comments

During the simulation tests the QOAS-based system has adapted switching the source of transmission from the 4.0 Mb/s stream to the 3.5 Mb/s and then to the 3.0 Mb/s one in the first test without experiencing any loss when transmitting *diehard1* sequence. It has also changed the transmission rate from 3.5 Mb/s to 3.0 Mb/s and 2.5 Mb/s and then back to 3.5 Mb/s during the background traffic variations as in the second test, experiencing short lossy periods. However it has maintained the average estimated end-user perceived quality above 4.0 (4.42 and 4.24 respectively). When similar conditions as in the simulation tests were emulated and the QOAS prototype system was tested it has also achieved subjective quality results above 4.0 for the same movie (4.00 and 4.22 respectively). These results confirm the objective testing results, in spite of some slight differences between the resulting values. However the test participants' subjective quality assessment related to all the movie sequences streamed in different delivery conditions using QOAS was around "good" ITU-T quality level, a fact that confirms the good performance of QOAS and recommends it as a viable solution.

These conclusions about QOAS are also confirmed by the comparison with the non-adaptive streaming and with the Windows Media MBR adaptive solution. The former, as expected,

does not succeed to provide even the least acceptable streaming quality in such loaded delivery conditions. The latter, although succeeds to adapt to the background traffic variation and maintains very high quality for each of the individual pictures that the video sequence is composed of, lacks a gracefulness of the display continuity between frames. Since detailed information about the adaptation was not made public, it can be only assumed that this is because the single-file-based Microsoft adaptive solution, initially designed for very low bit-rates, fails to achieve good performance for high bit-rate clips.

6.4 Conclusions

Both objective and subjective tests have shown good QOAS performance when streaming multimedia clips with different motion content in highly loaded delivery conditions and with different types, sizes and shapes of background traffic. These tests were performed over a topology that raises the same problems as a local broadband multi-service IP network. The QOAS performance was assessed in terms of end-user perceived quality, loss rate, link utilisation and number of simultaneous served customers and the results obtained highly recommend QOAS as a very efficient inexpensive solution that ensures the delivery of good quality multimedia-based services along other services via a local broadband IP-network to residential customers.

6.5 Summary

This chapter presents experimental results related to QOAS testing and includes presentation of both objective and subjective test results that complement each other. The objective tests involve tests that have aimed at tuning QOAS and the determination of some design-related parameters necessary to more accurately map the network-related parameters' variation into an application level quality of delivery score, based on which QOAS adapts. Simulations have then tested a QOAS model in loaded delivery conditions and subject to different background traffic commonly encountered in IP networks such as UDP (CBR and VBR) and TCP (long-lived and short-lived), with different shapes and sizes. Multimedia-like background traffic was also generated and QOAS was tested against it and successfully compared to other streaming solutions such as TFRCP, LDA+ and non-adaptive. Multiple simultaneous QOAS streaming processes were considered in order to determine the number of simultaneous viewers that can be served at "good" quality from a limited infrastructure and QOAS has again achieved better performances than other solutions. The effects of the variations in feedback frequency and delivery latency on QOAS were studied next. The chapter ends with a presentation of the results of a set of subjective tests that have

verified the results obtained by simulations. These have confirmed QOAS as a viable solution for streaming of high quality clips to local viewers, achieving very significant performance improvements over other solutions.

Chapter VII Conclusions

Abstract

This chapter summarises the research reported in this thesis, highlights the significant achievements of the proposed Quality Oriented Adaptation Scheme (QOAS), presents the contributions of the research and underscores the benefits of the proposed solution. Some future work directions are suggested at the end.

7.1 Main Achievements

This research aimed to find a solution for delivering high quality rich content multimedia-based services that would both be attractive to customers and beneficial for the service providers. Since existing solutions involve high complexity, increased deployment costs, and/or a lack of concern for the end-user perceived quality, a Quality-Oriented Adaptation Scheme (QOAS) was proposed. Its aim is to provide good end-user perceived quality for very high rate multimedia-based services in highly loaded and variable delivery conditions. If used in local broadband multi-service IP networks, QOAS allows for serving a larger number of customers from the same network infrastructure while maintaining good end-user perceived quality, bringing significant benefits to service providers and network operators and helping to ensure the success of these services.

This thesis proposed and presented in detail the principles and the mechanisms behind QOAS and analysed results of various tests. The effects on QOAS performance of increased traffic of different types, with various sizes and variation patterns as might be outputted by the other services delivered through the same infrastructure were assessed. The effects of multimedialike background traffic with different variations commonly expected due to the users' VCR interactivity with multimedia services were also tested. Similarly the effects of multiple simultaneous QOAS-based streaming processes were analysed. All these effects on QOAS performance were assessed in terms of end-user perceived quality, loss rate, link utilisation and number of customers served from a limited infrastructure.

The results were compared with those of an ideal adaptive scheme that uses all the available bandwidth not used by the background traffic at anytime in order to stream multimedia data, achieving 100 % utilisation and no-loss and with other proposed streaming solutions such as the adaptive TFRCP and LDA+ schemes and a non-adaptive mechanism.

These tests involve instantiations of QOAS in both a simulation model and a prototype system and the test results show very significant performances of the QOAS adaptive multimedia streaming stand-alone and in comparison to the other streaming solutions, regardless of the implementation used.

In comparison to the ideal adaptive scheme the results are very impressive in relation to the end-user perceived quality when using QOAS in simulated heavy traffic conditions for streaming multimedia clips subject to background traffic of different types, shapes and variation patterns. This subjective quality is not only above the "good" perceptual level (4 on the ITU-T 1-5 scale), but also in almost all cases it is within 1% from the corresponding value estimated for the ideal adaptive scheme. Moreover QOAS does not experience the latter's multiple variations with the bandwidth made available by the cross traffic that may disturb the viewers. QOAS streaming maintained loss rates of less than 0.1% in all cases, despite the fact that the delivery network was fully loaded. The link utilisation also reaches very high levels, QOAS making use of more than 99.5% of the bandwidth resources in the large majority of tests and even in the remaining cases the available resources are less than 1.5% from being fully used.

QOAS has achieved very good results also in comparison to other proposed approaches for streaming multimedia, such as adaptive TFRCP and LDA+ and a non-adaptive solution. For streaming in very heavy traffic conditions and subject to highly variable multimedia-like background traffic, QOAS has maintained the average end-user perceived quality above the "good" perceptual level, whereas for both adaptive schemes it was between the "fair" and the "good" subjective level and for the non-adaptive solution it was close to the "bad" level. The loss rate experienced by QOAS was very close to the ideal (0.015 %), whereas for the other tested adaptive schemes it exceeded 1 % and for the non-adaptive solution has even reached 13 %, severely affecting the quality of the delivery. The link utilisation was high for all the schemes, but QOAS has obtained the highest (99.93 %) after the non-adaptive solution whose 100 % link utilisation comes with a high price paid in end-user perceived quality. In terms of the number of customers served from a limited infrastructure, QOAS has also scored highly. For example, in order to maintain an average "good" perceptual quality level for all the streamed multimedia clips, 23% more clients could be served by using QOAS than by using TFRCP, 33% more clients than by using

LDA+, and 39% more users than by using the non-adaptive solution. If the goal is to maintain a "fair" average quality level for the clients, the benefit of using QOAS is 26% greater than TFRCP, 13% greater than LDA+, and 42% greater than the non-adaptive solution in terms on the number of simultaneous served customers.

In order to verify these very good results in terms of end-user perceived quality, extensive subjective tests that have involved the QOAS prototype system were performed. These tests have used multimedia clips with different motion content and apart from QOAS also other streaming approaches such as Microsoft's Windows Media (WM) Multiple bit-rate (MBR) adaptive streaming solution and a non-adaptive scheme. When the background traffic was varied in a similar fashion to multimedia clips subject to VCR-like interactivity with customers, QOAS has achieved subjective quality results above 4, the "good" ITU-T perceptual quality level, for all the movies streamed and in spite of the severely loaded and highly variable delivery conditions. In similar conditions WM MBR has scored on average only between "poor" and "fair", whereas the non-adaptive scheme has achieved on average between "bad" and "poor" on the same scale.

7.2 Novel Contributions

In this section the contributions made by the QOAS-related research are highlighted, making the solution original.

I) QOAS uses a novel client-located grading scheme that maps some network-related parameters' values, variation and variation patterns onto application-level QoS scores that describe the quality of the delivery. The Quality of Delivery Grading Scheme (QoDGS) monitors the packet loss, the packet delay and the delay jitter, which most seriously influence end-user perceived quality, as well as the end-user perceived quality as measured by a no-reference moving picture quality metric (Q). The three-stage QoDGS is based on both short-term and long-term evaluation of these monitored parameters' variations. Short-term variations are important for learning quickly about transient effects, such as sudden traffic changes and for reacting as fast as possible to them. Long-term variations are monitored in order to track slow changes in the delivery environment (e.g. new users). Taking into account the relative differences in the importance of the monitored parameters in relation to the characteristics of the delivery architecture (by weighting their contributions), short-term (QoD_{ST}) and long-term (QoD_{LT}) grades are computed. These partial grades are then used to determine the application-level quality of delivery score (QoD_{Score}).

Unlike the other sender-driven adaptive approaches for streaming multimedia clips that collect delivery-related information such as loss rate, delay etc. at the clients and send it to the server for processing, QOAS bases its adaptation on a different concept. Its idea is that apart from the monitoring of all these delivery-related parameters QOAS also distributes the quality of delivery grading process among its clients. As consequence, the *only feedback that is transmitted to the server consists of the client-computed QoD scores* that estimate the current delivery conditions and suggest quality adjustment decisions to be made by the server. This has an effect in lowering the complexity of the computations to be performed at the server, in reducing the quantity of information sent across the network and in increasing the accuracy of the grading since the client - as the receiver - is in better position to assess the quality of the delivery than the sender.

- II) The end-user perceived quality as estimated by a no-reference objective metric for multimedia streaming is actively considered during the adaptation, as part of the QoDGS's grading process. Since the goal of this adaptive scheme for streaming multimedia clips is to maximise the end-user perceived quality, it seems logical to conclude that by monitoring it during streaming and by taking it into account "in-service", the effectiveness of the adaptation is increased and better results can be achieved in terms of the remote viewers' perceived quality. This was shown by the results of the performed experimental tests, both objective and subjective.
- III) QOAS's tuning on an infrastructure that raises the same problems as a local broadband IP-network has produced very good results during the extensive testing sessions. Significant results are obtained in comparison to existing streaming solutions, adaptive or not, commercial or research-proposed in different delivery conditions. However, it is more significant that the QOAS's behaviour is very close to that of an ideal adaptive scheme, which is unlikely to be ever built, in terms of estimated end-user perceived quality, loss rate and link utilisation when used for multimedia streaming in the presence of traffic of different types, sizes and variation patterns.
- IV) QOAS allows for a significant increase in the number of customers that can be simultaneous served from an existing infrastructure while maintaining a good end-user perceived quality for the multimedia-based services offered, even in comparison with other existing solutions for delivering multimedia, adaptive or not.

7.3 QOAS Benefits

Delivering multimedia streaming-related services by using QOAS has some significant benefits that are mentioned next.

Simultaneous Access to Diverse Services

Both the service providers and the network operators on one hand and the customers on the other look forward to providing and having access to diverse services such as VoD, VoIP (IP telephony), high rate data transfers, etc. Unfortunately these services have different types and therefore various requirements that have to be accommodated by the same multi-service broadband IP-based infrastructure without interfering with each other. In this context the tests have shown that QOAS delivers multimedia-based services that gracefully adapt to traffic produced by other types of services and positively influences this traffic by reducing its share of bandwidth.

Increased Network Infrastructure Utilisation

In order to offer the best possible service quality at the lowest cost, service providers and network operators have to take full advantage of the existing network infrastructure. However, increasing the number of simultaneously served customers and the network utilisation decreases the quality of service in general. QOAS serves an increased number of customers from the same network infrastructure while maintaining a good quality level for the services provided.

Easy Scalability and Upgrade

The tremendous growth of the Internet and the fast evolution of the current cable TV services show that scalability is a significant problem for the designers and has to be taken into account. Another important problem, common to any engineering solution, is aging with the time and therefore updates are required from time to time. In relation to these problems, QOAS for delivering multimedia-based services to the residential users scores very well. The solution allows for more users to be added to the system at short notice and with no other investments apart from those related to their cable connection. Being a software solution, it also permits for upgrades to be made easily without any difficult problems to be overcome. In this context the QOAS's client-located Quality of Delivery Grading Scheme (QoDGS) and Server Arbitration Scheme (SAS) can be replaced by new, improved versions, if they are developed.

Providing Personalised Services

The scalability issue may have another dimension apart from number: heterogeneity of customers. In order to be considered acceptable, any multimedia-based solution has to be able to satisfy customers with different expectations. Therefore QOAS implements a "one-to-one" relationship with the customer by providing personalised, interactive, on-demand services.

Independence from Distributed Ownership

Providing content for services, providing connectivity and transporting the selected service to the receivers are three activities that could be completely separated from the ownership and administration point of view. Three different companies could be respectively the program provider, the service provider and the network operator, each with different policies and security issues that may not overlap, making the co-operation difficult if necessary for providing good quality services (e.g. deployment of QoS enhancements into the network may not be acceptable for the network operator or the service provider). Therefore solutions such as QOAS that offer independence from the manner the distribution of services to the customers is managed are highly desirable.

7.4 Future Work

Although the performances of QOAS as a solution for streaming high quality multimedia clips over local broadband multi-service IP-networks are already very close to an ideal adaptive scheme, there are some aspects in relation to the applicability of the scheme or to its potential extension that could be further explored. Next this section presents some of them.

Use of Error Control Solutions in Conjunction with QOAS

The QOAS-based solution for streaming high quality multimedia-based services in highly loaded delivery conditions in the current form which was extensively designed and tested did not take into consideration any error control mechanism to work in conjunction with apart from certain error resilient encoding provided by the MPEG compression scheme. However some of these error control mechanisms such as those based on retransmissions and on forward error control (FEC) have been assessed and considered not suitable to be used in conjunction with QOAS since they require supplementary bandwidth, which is not available in the expected highly loaded delivery conditions. Some other error control solutions, including error concealment techniques, seem suitable to reduce the effects of eventual packet losses during multimedia streaming on the end-user perceived quality. As a direct consequence different target loss rates have to be set for QOAS (i.e. higher) that take into account the application of these error concealment methods. These target limits have to be determined after extensive testing, including subjective ones, which should aim at assessing the results of these error control mechanisms on the remote multimedia viewers. Also these tests have to determine which of these mechanisms are best suited for applicability and in what delivery conditions they are recommended.

QOAS Independence from the Encoding Scheme

QOAS uses for in-service estimation of end-user perceived quality the no-reference moving picture quality metric (Q), as part of its client-located Quality of Delivery Grading Scheme (QoDGS). Q, as a no-reference metric, makes use of some a-priory knowledge about the encoding scheme – MPEG and the effect packet loss has on the MPEG-encoded stream. Since QOAS uses Q, the current version of QOAS is highly dependent on the encoding scheme that may limit its applicability. In consequence further work, aiming at increasing the QOAS generality, may explore possibilities to either use a no-reference metric that is independent from the encoding scheme or a set of different no-reference metrics for a number of popular encoding schemes.

Scalability and Real-Life Testing

For any proposed solution it is significant to allow for scalability. QOAS was such designed that permits new users to be added to the multimedia delivery system with ease. Also the simulation tests performed have shown very good results in terms of consequent end-user perceived quality when new users are added, the QOAS-based system achieving much better performances than other tested streaming solution. However real-life testing may be necessary in order to fully assess how the viewers are affected if their number increases and this is a direction future work may take.

Live Content Streaming Testing

The tests that have already been performed have mainly focused on streaming of prerecorded multimedia clips, which are only a part of the multimedia-based services. Since QOAS allows also for real-time adaptation of live transmissions, further work may include real-life testing of live multimedia deliveries. They can be performed either using an encoding card capable of adaptively modifying the encoding process or multiple encoding cards that encode the same content in different quality versions at the same time. The first case uses QOAS only to command the adaptive measures to be taken, whereas in the second QOAS controls the adaptive streaming in similar fashion it does with the pre-recorded streams, switching the source of transmission between the existing ones.

QOAS Extension for Multicasting

QOAS was designed to allow for a "one-to-one" relationship with the customers to which it is meant to provide on-demand, personalised multimedia-based services as part of a multi-service set offered via local broadband IP-networks. Among these services an important position due to its

current popularity has broadcast TV. This popularity is explained mainly by the high attraction of the main news programs, the important live transmissions, the major talk shows or the best-known series and by the viewers' routine-like daily and weekly schedule, which provides very much convenience for the customers. Although broadcast events can be replaced in an on-demand driven environment by multiple unicast deliveries of content to the viewers, it is a waste of resources to stream the same content over the same infrastructure multiple times. Therefore future work may take advantage of this common schedule for a number of viewers and propose solutions that would make use better of the shared resources. Using multicasting for such deliveries seems a good direction for research since although it introduces some overhead, if the number of simultaneous viewers is above a certain threshold, may achieve better performances. However this threshold, the architecture, the localisation of the customers and the complexity of the solution are very important and have to be taken into account in order to achieve good performances.

QOAS Extension for Low Bandwidth and/or Wireless Environments

QOAS was designed such as it currently targets very high bit-rate multimedia streaming over broadband wireline IP networks. However, many services are currently being delivered through much narrow bandwidth links and it is expected that some of them, including multimedia-based services, to complement the broadband related ones, offering to the users a rich set of heterogeneous services. Therefore QOAS may be extended to target lower bandwidth environments, including wireless ones that introduce supplementary challenges such as higher and less predictable loss rates, different encoding schemes such as MPEG-4, for instance and user mobility.

QOAS Extension with MPEG-4

A significant extension to QOAS could take into account the object-based structure of the MPEG-4 encoding solution for multimedia streams. For instance the different quality versions defined for the same multimedia content may not be totally exclusive as in the current solution, but more like complementing each other. The fine granularity scalability (FGS) or the progressive fine granularity scalability (PFGS) that were proposed for MPEG-4 can be used in order to transmit first a base layer ("must have") and, if the delivery conditions permit, other enhancement layers that would increase the overall quality of the multimedia streams.

7.5 Summary

This last chapter presents first the conclusions drawn after the Quality-Oriented Adaptation Scheme (QOAS) has been designed and tested. The results of these tests, both objective and subjective, are briefly summarised and QOAS's significant performances, stand-alone and in comparison to other solutions, are listed indicating very important achievements. Next the contributions of the research performed and presented in this thesis are listed and briefly commented. The most important benefits of the QOAS-based solution for streaming multimedia are also presented in this chapter, which ends with some suggestions for future work.

Appendix A

Definitions for Technical Terms

Connectivity = a host A has "Type-P-Instantaneous-Unidirectional-Connectivity" to a host B at time T if a type-P packet transmitted from A to B at time T will arrive at B. Bidirectional connectivity refers to unidirectional connectivity from A to B and from B to A. [203]

One-way delay = the "Type-P-One-way-Delay" from host A to host B is dT at moment T if the host A sent the first bit of a type-P packet to B at moment T and host B received it at moment T+dT. The one-way delay is undefined (in fact, infinite) if the packet does not arrive at host B. [204]

One-way loss = the "Type-P-One-way-Packet-Loss" from host A to host B is 0 at moment T if the host A sent the first bit of a type-P packet to B at moment T and host B received that packet. If host B did not receive that packet "Type-P-One-way-Packet-Loss" at moment T is 1. [210] Note: In practice the one-way loss is measured over a period of time and is expressed as a percentage of the total number of packets sent.

Error propagation = the process of spreading of an error effect to a larger area, involving parts that were not affected directly by the original error. [65] Note: The term is used in multimedia streaming in relation with MPEG encoding and refers to the fact that an error that affects the compressed data that corresponds to a reference frame will affect not only this frame, but also other frames that use the reference frame data for decoding. [133]

Round-trip delay = the "Type-P-Round-trip-Delay" from host A to host B at moment T is dT if host A sent the first bit of a type-P packet to B at time T, B received it and immediately sent another type-P packet back to A that has received the last bit of that packet at time T+dT. The "Round-trip-Delay" from A to B at T is undefined (informally, infinite) if A sent the first bit of a type-P packet to B at time T but either B did not receive the packet, B did not send a type-P packet in response or A did not receive that response packet. [212]

Round-trip loss = the "Round-trip-Loss" between host A and host B is equal to the percentage of the packets sent by host A to host B that were followed by received answer packets by the host A from the total number of packets sent. Host B is supposed to receive the packets sent by host A and answer by sending other packets back to host A^{52} .

One-way delay variation = the "One-way-IP-Packet-Delay-Variation" for two packets sent from host A to host B, as the difference between the value of the One-way-delay for the second packet at T2 and the value of the One-way-Delay for the first packet at T1. T1 is the time at which A sent the first bit of the first packet, and T2 is the time at which A sent the first bit of the second packet. [213] Note: An alternate, but related, way of computing an estimate of delay variation (jitter) is given in RFC 1889 [100]. By taking the absolute values of the delay variation sequence (as defined in [213]) and applying an exponential filter with parameter 1/16 the estimate is generated: j_new = 15/16* j_old + 1/16*j_new.

Network congestion = a network situation when the traffic increases above a certain limit, the routers are not able to cope with the number of packets to be routed and they begin loosing them. If the traffic further increases, almost no packets are delivered. [215]

Loss pattern = refers to the manner the loss occurs. IETF IPPM Working Group [61] has defined two loss pattern metrics: the "loss period" metric captures the frequency and length (burstiness) of loss once it starts, and the "loss distance" metric captures the spacing between the loss periods. [216]

Packet reordering = refers to the process necessary to be performed in order to restore at the destination the order in which the packets were sent. In general packet order is not expected to change during transmission from a host to another one, but there are cases when it does change. For example when a single packet stream is sent from a host to another one between which there are two paths, one with slightly longer transfer time, the packets traversing the longer path may arrive out-of-order. [220]

Bulk transport capacity = measures the network's ability to transfer significant quantities of data with a single congestion-aware transport connection (e.g., TCP). The formal is: data_sent / elapsed_time, where "data_sent" represents the unique "data" bits transferred (i.e., not including header bits or emulated header bits). Note: The amount of data sent should only include the unique number of bits transmitted (i.e. if a particular packet is retransmitted the data it contains should be counted only once). [223]

Available bandwidth = maximum end-to-end throughput given cross traffic load. It is a metric that varies with the time, background traffic type and variation pattern, used in general as an average over certain time interval. [255]

Wire time = historically the term was to loosely denote the time at which a packet appeared on a link, without exactly specifying whether this refers to the first bit, the last bit, or some other consideration. This informal definition makes a distinction between when the packet's propagation delays begin and cease to be due to the network rather than the endpoint hosts. [62]

Clock offset = the difference between the time reported by the clock and the "true" time as defined by the universal time clock at a particular moment. If the clock reports a time Tc and the true time is Tt, then the clock's offset is Tc - Tt. [62, 204]

Synchronized clocks = a pair of clocks that are "accurate" with respect to one another (their relative offset is zero). Note: Clocks can be highly synchronized yet arbitrarily inaccurate in terms of how well they tell true time. For many measurements, synchronization between two clocks is more important than the accuracy of the clocks. [62, 204]

Clock accuracy = indicates how close the absolute value of the clock's offset is to zero at a particular moment. Ideally this should be 0. [62, 204]

Clock resolution = the smallest unit by which the clock's time is updated. It gives a lower bound on the clock's uncertainty. Resolution is defined in terms of seconds. However, resolution is relative to the clock's reported time and not to true time, so for example a resolution of 10 ms only means that the clock updates its notion of time in 0.01 second increments, not that this is the true amount of time between updates. Note: Clocks can have very fine resolutions and yet be wildly inaccurate. [62, 204]

Clock skew = the frequency difference (first derivative of its offset with respect to true time) between the clock and true time at a particular moment. [62, 204]

Clock drift = the variation in skew exhibited by some real clocks (the second derivative of the clock's offset with respect to true time). [62]

Host = a computer capable (if all is working properly) of communicating using the IP protocols. Note: Includes **routers**. [154]

Link = the link-level abstraction of a "virtual direct connection" between two or more hosts. Often thought of in terms of a single underlying physical connection. [154]

Router = a host that facilitates communication between other hosts by forwarding packets from one link to another. [154]

Path = the network-level abstraction of a "virtual link" from host A to host B. The Internet Protocol (IP) makes it appear to higher levels as though the host A has a *direct* connection to B. This apparent direct connection is a "path." The notion of "path" is a unidirectional concept. [154]

Route = a sequence of links and routers comprising a path. [154]

Router buffer size = the number of bits the router has available for buffering queued packets (the router is seen as a queueing server). [154]

Link bandwidth = a link's data-carrying capacity, measured in bits per second, where "data" does not include those bits needed solely for link-layer headers. [154]

Link propagation time = the time difference in seconds between the moment when host A on the link A-B begins sending one bit to host B and the moment when host B has received the bit. [154]

Motion vector = A two-dimensional vector used for motion compensation that provides an offset from the coordinate position in the current picture or field to the coordinates in a reference frame or reference field. [75, 76]

Mutex object = A mutual exclusion object that allows multiple threads to synchronise access to a shared resource. A mutex has two states: locked and unlocked. Once a mutex has been locked by a thread, other threads attempting to lock it will block. When the locking thread unlocks (releases) the mutex, one of the blocked threads will acquire (lock) it and proceed.

Appendix B MPEG 1 and MPEG 2 Encoding Schemes

B.1 MPEG 1 and MPEG 2 Video

The MPEG video compression algorithm is based on the fact that the human eye is more sensitive to brightness changes than chromatic ones. Therefore, in order to achieve compression, the image data is divided into one luminance and two chrominance components, the latter of a smaller size. After this lossy step the compression method used is Discrete Cosine Transform (DCT) and then quantisation (Q). These reduce the high spatial frequency components from the image based on the observation that the human viewer is more sensitive to the reconstruction errors of low frequency components. The quantisation is meant to reduce the precision of the DCT-coefficients according to the required image quality. The higher the O factor, the lower quality of the image will be obtained after decompression. A zig-zag scan arranges the low frequency coefficients to the beginning of the stream. The upper left corner coefficient represents the mean value of the block and is encoded using the difference from the previous block (DPCM). Since most of the high frequency coefficients are zero after the quantisation, run length encoding (RLE) is used for further compression. The final step in the compression process is to minimize the entropy using Huffman (or arithmetic) coding. The frame encoded in this way is called I-frame (intra frame) because the encoding process uses no information about other frames. A description diagram of the encoding process is presented in Figure B-1.

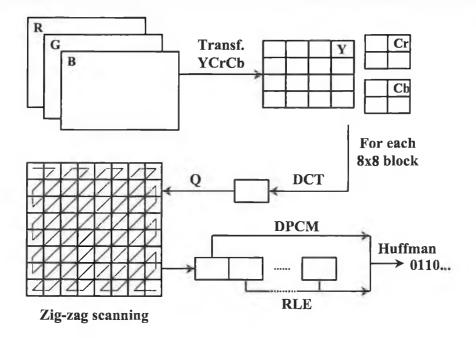


Figure B-1 The Spatial Compression Technique

Apart from the spatial compression technique, the temporal redundancy between frames can be reduced for further compression. The idea is to calculate a prediction error between corresponding blocks in the current and previous frames. Only the error values are then send to the compression process. The frames obtained by compressing prediction error values are usually called P-frames (prediction frame). If both previous and future frames are used as reference, the frame is called B-frame (bi-directional frame). Motion compensated prediction is an efficient tool to further reduce temporal redundancy between frames. The aim is to obtain the motion estimation between video frames. The motion is described by a small number of motion vectors, which gives the translation of a block of pixels between frames. The motion vectors and compressed prediction errors are then used. A short graphical description of the algorithm is presented in Figure B-2.

From structural point of view, the video stream consists of a number of video sequences, each of them having a sequence header and consisting of at least one group of picture. The latter has as components one or more pictures (frames). At this level, a typical encoded sequence is: I B B P B B I B B P B B..., but the actual pattern is up to the encoder. Each picture is made of slices, which accommodate macroblocks. Each macroblock has 6 blocks, 4 describing the luminance and 2 for the chrominance components.

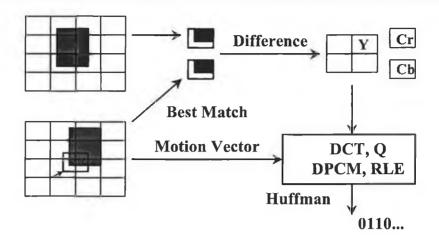


Figure B-2 The Temporal Compression Technique

B.2 MPEG 1 and MPEG 2 Audio

Because the uncompressed CD audio has 44,100 samples/sec, 16 bits/sample and 2 channels, which is difficult to transmit, the need for compression of audio is evident.

MPEG-1 Audio supports one (mono) or two (stereo) audio channels with a sampling rate of 32, 44.1, or 48kHz. The compressed bitstream varies with fixed bitrates in a range from 32 to 224kbits/sec per channel. This gives a compression grade ranging from 2.7 to 24 times, depending on the sampling rate. MPEG-1 Audio is divided into three parts, referred to as layers. Each layer describes a different method of encoding the audio, and higher-numbered layers involve higher complexity in the encoding and decoding processes, although all three methods are based on similar principles.

One of the key components of the encoding of MPEG Audio is the use of a psychoacoustic model and the use of the so-called masking effect. The masking effect relies on the observation that a human subject will not hear weaker sounds located near a strong sound, so they can be removed. First the frequency spectrum (20 Hz-20 kHz which is the range of human hearing) is broken up into 32 equal-width frequency bands, of 12 or 36 sub-bands each. This automatic filtering out of the inaudible frequencies achieves a direct saving. Then the filter examines each band, and identifies the key tones in each band. It calculates the masking effect of each tone, establishes a threshold for each band and removes all irrelevant (masked) tones from the band. Further on, different algorithms are used to achieve an efficient bit-stream formatting, according to the MPEG layer.

The Layer I algorithm is the simplest one and is best suited for bit rates above 128kbits/sec per channel. 384 audio samples are coded into every frame. Layer II is a bit more complex and improves the compression rate by coding data in larger groups. Layer II use 1152 samples/frame, which is the same as in Layer III and is targeted for bit rates around 128 kbits/s per channel. Layer III is the most complex but offers the best audio quality, particularly for bit rates around 64 kbits/s per channel.

MPEG-2 Audio extends the MPEG-1 standard with a set of additional features. The big difference is the support for multichannel and the multilingual support. It supports up to five high fidelity audio channels and one low frequency enhancement channel. This is perfectly suited for digital movies where you want surround sounds. The standard also has support for up to seven additional commentary channels. Another feature is the additional support for lower, compressed bitrates down to 8kbits/sec. MPEG-2 also introduces support for 16, 22.05 and 24kHz. The commentary channels are allowed to have a sampling rate that is half the high fidelity channel.

All MPEG Audio frames start with a 32-bit header. The header consists of an ID flag, a layer flag, an error protection flag, a bitrate index, a sampling frequency index, a mode flag, and other less important data (such as copyright, etc.). After the header the encoded data is placed in a format, which depends on the specific layer.

B.3 MPEG -1 Systems and MPEG 2 Program

The MPEG Systems stream combines one or more streams of video and audio as well as other data, into a single stream suitable for storage or transmission. The syntactical and semantic rules imposed by the standard enable correct synchronization and playback.

The basic principle of MPEG-1 Systems coding is the use of time stamps which specify the decoding and display time of audio and video and the time of reception of the multiplexed coded data at the decoder. This allows for a great degree of flexibility in decoder design, the number of streams, multiplex packet lengths, video picture rates, audio sample rates, coded data rates, digital storage medium or network performance. It also provides flexibility in selecting which entity is the master time base, while guaranteeing that synchronization and buffer management are maintained. Variable data rate operation is supported. A reference model of a decoder system is specified which provides limits for the ranges of parameters available to encoders and provides requirements for decoders.

The stream consists of a continuous sequence of elementary stream packets (known as "pack"-s). Each pack includes information regarding the clock reference and the stream rate. A System header follows and then one or more Packet blocks. Apart from packet data, each has a presentation time stamp, which helps the decoder in its playing process and a stream ID which indicates whose stream the packet belongs to.

The MPEG-2 Program stream is similar to the MPEG-1 Systems standard. It includes extensions to support new and future applications. Both are built on a common Packetized Elementary Stream packet structure, facilitating common video and audio decoder implementations and stream type conversions

Appendix C

Documents for Subjective Testing

Personal Information Page			
Ç		Record No:	
Gender:	a) male	b) female	
Age:			
Do you use glasses/contact lenses:	a) yes	b) no	
Are you long/short sighted:	a) long sighted	b) short sighted	c) no
Do you have other visual conditions that may affect your perception of movies (e.g. color blindness, glare):	a) yes	b) no	
How familiar are you with multimedia streaming:	a) I work in this domain	b) I am familiar	c) I am not familiar
Dou you rent DVDs/tapes:	a) often	b) sometimes	c) never
Do you go to cinema/theatre:	a) often	b) sometimes	c) never
Would you like to watch movies via Video on Demand streaming to your home (e.g. via cable TV):	a) yes	b) no	
Name (optional*):			
E-mail/phone no. (optional*);			

Disclaimer

The information collected will be kept separately from the perceptual test results and it will not be made public under any form. The name and e-mail address are collected only to allow us to deliver the prizes after the draw.

^{*} Fill the optional fields if you want to take part in the draw for prizes. This allows us to contact you.

Questionnaire

Record No:	
Test Type:	

Directions

Could you kindly answer the following questions about the last sequence shown?

- A) Grade the perceived quality of the streamed multimedia clip on the 1 (the worst quality) to 5 (the best) subjective scale presented in the following table.
 - B) State what you liked about the clip shown (e.g clarity, continuity etc.).
 - C) State what you disliked about the clip shown (e.g blurriness, discontinuity etc.).

Quality scale for subjective testing (ITU-T R P.910)

Rating	Quality	Impairment		
5	Excellent	Imperceptible		
4	Good	Perceptible, not annoying		
3	Fair	Slightly annoying		
2	Poor	Annoying		
1	Bad	Very annoying		

Example of Answer Sheet

Phase No:		
A) Grade the perceived	quality of the streamed multimedia clip	:
Clip Code:		Grade:
B) State what you have	appreciated at the clip shown (please inc	dicate others if any):
Continuity:	Quality Stability:	
Clarity:	Media Synchronisation:	
C) State what you have	disliked at the clip shown (please indica	ate others if any):
Jerkiness:	Quality Variation:	Bluring:
Tiling:	Media Desynchronisation:	
	259	

Answer Sheet

	Phase No:	
Clip Code:		
A) Grade the perceived qual	ity of the streamed multimedia clip:	
Grade:		
B) State what you have app	reciated at the clip shown (please indicate	e others if any):
Continuity:	Quality Stability:	
Clarity:	Media Synchronisation:	
C) State what you have disli	ked at the clip shown (please indicate oth	ers if any):
Jerkiness:	Quality Variation:	Bluring:
Tiling:	Media Desynchronisation:	
Clip Code:]	
A) Grade the perceived qua	ality of the streamed multimedia clip:	
Grade:		
B) State what you have app	reciated at the clip shown (please indic	cate others if any):
Continuity:	Quality Stability:	
Clarity:	Media Synchronisation:	
C) State what you have disl	iked at the clip shown (please indicate	others if any):
Jerkiness:	Quality Variation:	Bluring:
Tiling:	Media Desynchronisation:	

Clip Code:		
A) Grade the perceiv	ed quality of the streamed multimedia	clip:
Grade:		
B) State what you ha	ve appreciated at the clip shown (please	e indicate others if any):
Continuity:	Quality Stability:	
Clarity:	Media Synchronisation:	
C) State what you ha	ve disliked at the clip shown (please ind	licate others if any):
Jerkiness:	Quality Variation:	Bluring:
Tilings	Madia Dagymahyanisation	

Test Instructions

Welcome Message

Welcome to the perceptual testing session organised by the Performance Engineering Laboratory, Dublin City University.

Test Objectives

We have proposed a novel approach for multimedia streaming and we want to test it and compare it to other approaches. These subjective tests you take part in aim at quantifying the perceived quality of different multimedia clips, streamed using various approaches.

Disclaimer

Please fill in the personal information page. The information collected will be kept separately from the perceptual test results and will never be made public in any form. Your name and e-mail address are collected only to allow us to deliver the prizes after the draw.

Test Directions

The test consists of four phases. In each phase you will be first shown a high quality clip that gives you a reference for your judgement. Next you will be shown a series of multimedia clips and you will be asked to grade their quality on the indicated 1-5 scale. The grading is done immediately after the clip has ended. You are not allowed to change the screen position, the distance from the screen or to turn the speakers louder since they are fixed for all the test subjects. Once the test has started you are not allowed to pause it or to stop it or to ask questions. However if you feel bored or tired, you can leave the testing room anytime.

Example

Phase [X]

This is the reference clip for the [X]-th phase:

[Clip streaming]

The [Y]-th multimedia clip is shown next. Immediately after finishing it please answer the questions.

[Clip streaming]

Please grade its quality and answer the questions.

[Grading & Answering].

...

Our test has ended. Could we have your forms, please?

[Collection of all the forms]

Thank you for your kind participation.

Questionnaire

Could you kindly answer the following questions about the last sequence shown?

- A) Grade the perceived quality of the streamed multimedia clip on the 1 (the worst quality) to 5 (the best) subjective scale presented in the given table.
- B) State what you liked about the clip shown (e.g clarity, continuity etc.)
- C) State what you disliked about the clip shown (e.g blurriness, discontinuity etc).

List of Publications

Gabriel-Miro Muntean, Philip Perry, Liam Murphy, "A New Adaptive Multimedia Streaming System for All-IP Multi-Service Networks", accepted, IEEE Transactions on Broadcasting, 2003

Gabriel-Miro Muntean, Liam Murphy, "Adaptive Pre-recorded Multimedia Streaming", Proceedings of IEEE GLOBECOM 2002, Taipei, Taiwan, November 2002

Gabriel-Miro Muntean, Liam Murphy, "Adaptive Traffic-Based Techniques For Live Multimedia Streaming", Proceedings of the 9th IEEE International Conference on Telecommunication ICT'2002, Beijing, China, June 2002

Gabriel-Miro Muntean, Liam Murphy, "An Adaptive Mechanism For Pre-recorded Multimedia Streaming Based On Traffic Conditions", Proceedings of the 11th W3C World Wide Web Conference, Honolulu, Hawaii, USA, May 2002

Gabriel-Miro Muntean, Liam Murphy, "Feedback Controlled Traffic Shaping for Multimedia Transmissions in a Real-Time Client-Server System", Lecture Notes in Computer Science 2093, Springer-Verlag, 2001, Vol. II, pp. 540-548, ISSN 0302-9743

Gabriel-Miro Muntean, Liam Murphy, "A Novel Feedback Controlled Multimedia Transmission Scheme", Proceedings of the 8th IEEE International Conference on Telecommunication, Bucharest, Romania, June 2001, Vol. III, pp. 123-128

Gabriel-Miro Muntean, Liam Murphy, "Experimental Results for a Feedback-Controlled Multimedia Transmission System", Proceedings of the 17th IEE UK Teletraffic Sysmposium 2001, Dublin, Ireland, May 2001, pp. 15/1-15/6

Gabriel-Miro Muntean, Liam Murphy, "Some Software Issues of a Real-Time Multimedia Networking System", Transactions on Automatic Control and Control Science, Vol. 45, No. 59/III, pp. 35-40, Romania, 2000, ISSN 1224-600X

Gabriel-Miro Muntean, Liam Murphy, "An Object Oriented Prototype System for Feedback Controlled Multimedia Networking", The Irish Signals and Systems Conference 2000, University College of Dublin, Ireland, June 2000, pp. 173-180

Awards

Best Student Poster Paper Award - The 11th W3C World Wide Web Conference, Honolulu, Hawaii, USA, May 2002, for "An Adaptive Mechanism For Pre-recorded Multimedia Streaming Based On Traffic Conditions"

Best Paper Award - The 8th IEEE International Conference on Telecommunication, Bucharest, Romania, June 2001, for "A Novel Feedback Controlled Multimedia Transmission Scheme"

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